TeleVideo TS 806, TS 806C, TS 806H, TS 806/20 and TS 806H/20 Maintenance Manual

TeleVideo Systems, Inc.

TeleVideo® TS 806, TS 806C, TS 806H, TS 806/20 and TS 806H/20 Maintenance Manual

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PREFACE

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TS-806 THEORY OF OPERATION

This theory of operation describes the hardware layout, functions and operations.

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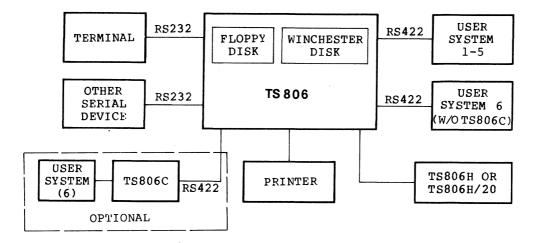
1. INTRODUCTION

The TS-806 is a multi-user system that supports up to 6 user stations. The user stations may be configured from TS-800, TS-801, or TS802-series systems.

The TS-806 is a single board microcomputer based on the Z80A family of microprocessor components. The system functions a central resource manager in a multi-user environment through RS-422 SDLC control. The terminal used with the TS-806 is a standard TeleVideo terminal used as a service terminal.

The TS806 contains a 10-megabyte Winchester technology hard disk drive. The TS806/20 uses a 20-megabyte Winchester drive in the same style cabinet as the TS806. Both systems use the WD-1000 disk interface board. The TS806 can be expanded with the TS-806H, and the TS806/20 is expanded with the TS806H/20.

The TS806C is a magnetic cartridge option to the TS806-series systems. The TS806C is interfaced to the TS806-series system through one of the RS-422 high speed serial data links used for a user station. A user station can be connected to the TS806C as shown in the block diagram to allow the full compliment of 6 stations.



2. GENERAL DESCRIPTION

a. Z80A-CPU

The main processing unit in the system (4.0 MHz operation).

b. Z80A-CTC

Counter/timer chip, generating the baud rate for RS232C serial channel.

c. Z80A-PIO

Parallel I/O chip, provides centronics-type printer interface.

d. Z80A-SIO

Serial I/O chips, providing RS232C interface (data rate to 19.2 KB) and/or RS422 high-speed serial interface with data rate of 800K bit/sec.

e. Z80A-DMA

Direct memory access controller chip is used for direct transfer of data between memory and peripheral I/O like floppy disk, Winchester hard disk, etc.

f. Memory

Main memory available to user is 64K bytes, using 64K X 1 DRAM 4K bytes of ROM and 1K byte of SRAM are used for system firmware (used on power up or reset only).

g. Floppy disk controller

Western Digital's FD 1793 chip along with its support chips WD 2143-01 and WD1691 provides control and interface to the 5 $\frac{1}{4}$ " floppy drive.

h. Winchester hard disk interface

Interface is provided to the disk controller board for the 5 $\frac{1}{4}$ " hard disk.

3. FUNCTION

The facilities available to the users are 64K bytes of main memory, floppy disk controller which can drive potentially up to 4 drives (0.5M byte each), interface to communicate with the hard disk controller which can also drive potentially up to 4 hard disk drives 5 ¼" Winchester, and parallel port for high-speed centronics-type printer interface.

2 channels of RS232C type interface (from 1 SIO), and 6 channels of RS422 interface (from the other 3 SIO's) to communicate with the user systems (TS-800, TS-801) are provided.

4K bytes of ROM and 1K bytes of SRAM are used for system initialization, diagnostics, boot, and floppy/hard disk control during the power-up or reset, and are not accessible by the users. After the initial program in the ROM is run upon power-up or reset, the dynamic memory area (address 00-16 K hex.) is switched on so that the whole 64K byte of main memory can be used by the user.

TS-806 uses a DMA controller to transfer data between memory and I/O devices, memory to memory, and also I/O device to I/O device.

A CTC chip (A40) is reserved to provide vectored interrupt capability. Another CTC chip (A39) generates (switch-selectable) baud rates for RS232C interface. 2 Channels of this CTC are reserved for time-of-day generation as a user implemented option.

4 LED indicators are used for diagnostic purposes to indicate faults in the hardware.

The system block diagram is shown in Figure 2, I/O port assignment is shown in Table 1 and the baud rate switch configuration is shown in Table 2.

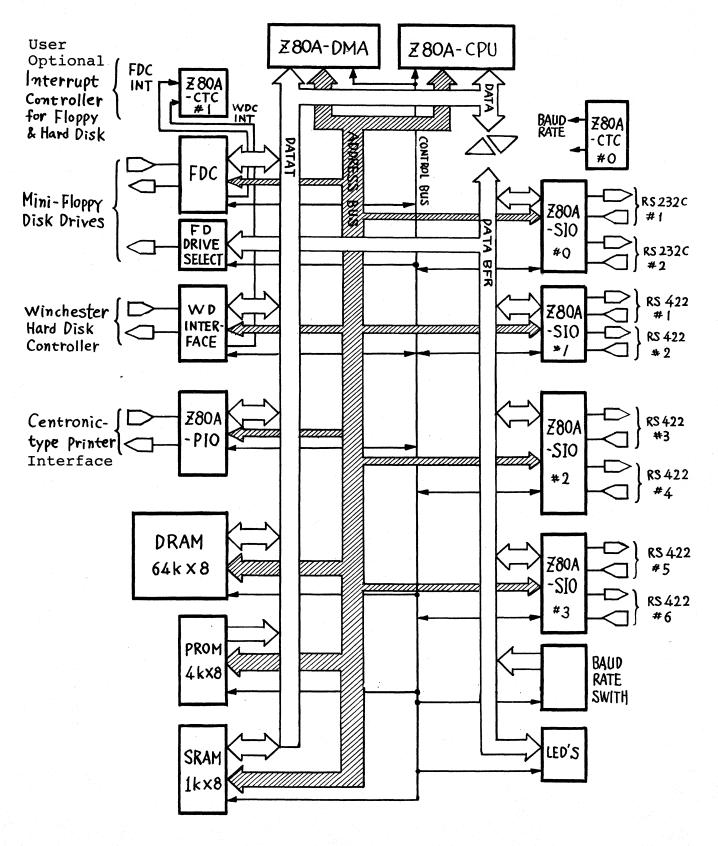


DIAGRAM 2 - BLOCK DESIGN OF TS-806

ABUS BIT #

I/O PORT

| - | | | | | | | | | |
|---|-----|---|---|----------|---|---|---|------|---|
| 7 | 6 | 5 | 4 | 3 | 2 | 1 | 0 | HEX. | |
| 0 | 0 | 0 | 0 | 0 | х | x | х | 00 | Indicator Load (LED Ref.) |
| 0 | 0 | 0 | 1 | 0 | х | 0 | 0 | 10 | SIO 1, Ch A Data Reg. |
| | | | | | | 1 | 0 | 12 | Com/Stat Reg |
| | | | | | | 0 | 1 | 11 | Ch B Data Reg. |
| | | | | | | 1 | 1 | 13 | Com/Stat Reg |
| 0 | 0 | 1 | 0 | 0 | х | 0 | 0 | 20 | SIO 2, Ch A Data Reg. |
| | | | | | | 1 | 0 | 22 | Com/Stat Reg |
| | | | | | | 0 | 1 | 21 | Ch B Data Reg. |
| | | | | | | 1 | 1 | 23 | Com/Stat Reg |
| 0 | 0 | 1 | 1 | 0 | х | 0 | 0 | 30 | SIO 3, Ch A Data Reg. |
| | | | | | | 1 | 0 | 32 | Com/Stat Reg |
| | | | | | | 0 | 1 | 31 | Ch B Data Reg. |
| | | | | | | 1 | 1 | 33 | Com/Stat Reg |
| 0 | 1 | 0 | 0 | 0 | х | х | х | 40 | WDC Reset (Winchester Disk Soft Reset) |
| 0 | 1 | 0 | 1 | 0 | Х | 0 | 0 | 50 | SIO 0, Ch A Data Reg. |
| | | | | | | 1 | 0 | 52 | Com/Stat Reg |
| | - , | | | | | 0 | 1 | 51 | Ch B Data Reg. |
| | | | | | | 1 | 1 | 53 | Com/Stat Reg |
| 0 | 1 | 1 | 0 | 0 | Х | Х | Х | 60 | Baud Rate Load |
| 0 | 1 | 1 | 1 | 0 | х | Х | Х | 70 | FDD Select |
| 1 | 0 | 0 | 0 | 0 | X | 0 | 0 | 80 | CTC Ch 0 |
| | | | | | | 0 | 1 | 81 | Ch 1 |
| | | | | | | 1 | Q | 82 | Ch 2 |
| | | | | | | 1 | 1 | 83 | Ch 3 |
| 1 | 0 | 0 | 1 | 0 | Х | X | х | 90 | DMA |
| 1 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | A0 | WDC Reg. 0 Data Reg. |
| | | | | | 0 | 0 | 1 | Al | Reg. 1 Error/Write Precompensation Reg. |
| | | | | | 0 | 1 | 0 | A2 | Reg. 2 Sector Count Reg. |
| | | | | | 0 | 1 | 1 | A3 | Reg. 3 Sector Number Reg. |
| | | | | | 1 | 0 | 0 | A4 | Reg. 4 Cylinder Low Reg. |
| | | | | | 1 | 0 | 1 | A5 | Reg. 5 Cylinder High Reg. |
| | | | | | 1 | 1 | 0 | A6 | Reg. 6 Size/Drive/Head Reg. |
| | | | | | 1 | 1 | 1 | A7 | Reg. 7 Status/Command Reg. |
| | | | | | | | | | |

ABUS BIT #

I/O PORT

| 7 | 6 | 5 | 4 | 3 | 2 | 1 | | HEX. | |
|----|---|---|---|---|---|---|---|------|----------------------------|
| 1 | 0 | 1 | 1 | 0 | Х | 0 | 0 | в0 | FDC Com/Stat Reg. |
| | | | | | | 0 | 1 | Bl | Track Reg. |
| | | | | | | 1 | 0 | В2 | Sector Reg. |
| | | | | | | 1 | 1 | В3 | Data Reg. |
| 1 | 1 | 0 | 0 | 0 | х | 0 | 0 | CO | PIC Ch 0 (CTC) User Option |
| | | | | | | 0 | 1 | Cl | Ch 1 |
| | | | | | | 1 | 0 | C2 | Ch 2 |
| | | | | | | 1 | 1 | C3 | Ch 3 |
| 1 | 1 | 0 | 1 | 0 | x | 0 | 0 | D0 | PIO Ch A Data Reg. |
| | | | | | • | 1 | 0 | D2 | Com/Stat Reg |
| | | | | | | 0 | 1 | Dl | Ch B Data Reg. |
| | | | | | | 1 | 1 | D3 | Com/Stat Reg |
| 1 | 1 | 1 | 0 | 0 | X | х | х | EO | Enable Dram |
| 1. | 1 | 1 | 1 | 0 | Х | Х | Х | FO | Disable Dram |
| | | | | | | | | | |

Table 1. I/O Port Assignment

Baud Rate Generation for SIO:

Switches for baud rates are upper 4 dip switches which can be set up as following table.

| | | Sw | itc | h P | os. |
|------|-------|----|-----|-----|-----|
| Baud | Rate | 4 | 3 | 2 | 1 |
| | | | | | |
| 19.2 | KBaud | 0 | 0 | 0 | 0 |
| 9.6 | KBaud | 0 | 0 | 0 | 1 |
| 4.8 | KBaud | 0 | 0 | 1 | 0 |
| 2.4 | KBaud | 0 | 0 | 1 | 1 |
| 1.2 | KBaud | 0 | 1 | 0 | 0 |
| 600 | Baud | 0 | 1 | 0 | 1 |
| 300 | Baud | 0 | 1 | 1 | 0 |
| 150 | Baud | 0 | 1 | 1 | 1 |
| 75 | Baud | 1 | 0 | 0 | 0 |

Table 2. Baud Rate Switch

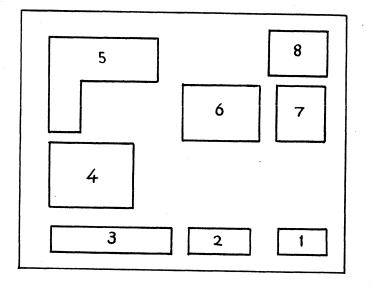
Dip Switch Description:

0 (Close) / 1 (Open)

| SW | 1: | Baud | Rate | 9 |
|----|----|----------|------|------------|
| | 2: | · · · · | | |
| | 3: | " | | |
| | 4: | " | | |
| | 5: | Unused | | |
| | 6: | Unused | | |
| | 7: | Unused | | |
| | 8: | Reserved | for | Diagnostic |

4. Circuit Description

Generally TS-806 circuitry can be divided into 8 sections as shown in Figure 3.



1. Reset

2. Clock

3. Memory

- 4. CPU & DMA
- 5. Serial I/O

6. Floppy Disk Control

7. Hard Disk Controller

8. Printer Interface

Figure 3 System Board Description

4.1 Reset

The chip associated with the reset circuit is a 74LS00 (A2), which debounces the reset signal when the reset switch is pressed.

A2-13 is normally low and Pin 1 is normally high. -Reset must be active for at least three clock cycles to properly reset the CPU. As long as -Reset remains active, the address and data bus float, and the control signal outputs are inactive.

The CPU returns to normal operation after two internal "T" cycles. -Reset clears the PC register so the register now contains a value of 0000.

4.2 Clock Generation

16 MHz signal from the oscillator (OSC) is divided down to a 4 MHz signal using a 93S16 (A7) and TR 2N2907 is used to drive the proper system clock level and again buffered for enough driving current.

The 1 MHz clock from the 93S16 is supplied to the floppy disk controller chip (FD 1793). To generate the baud rate, the 8 MHz signal from the 93S16 is fed into a 74LS163 (A8), to generate the frequency of 1/6.5 times the system clock for the CTC. Transmit clock for the SIO (RS422) is generated through another 74LS163 (A38) and 74S04 (A23) using system clock.

4.3 Memory

ROM and Static RAM:

A single 4K X 8 ROM and two 1K X 4 SRAM are used. ROM and SRAM are enabled upon power-up or reset. The enabling circuitry is composed of two 74LS138 (A21, A22) and other necessary chips of flip-flop and gates. By issuing an I/O instruction, we can either enable or disable the ROM and SRAM. When the ROM and SRAM are enabled, -CAS signal for dynamic RAM column address goes inactive disabling dynamic RAM (00H-16KH), so only ROM and SRAM are active. During read/write operations involving ROM and SRAM the 74LS74 (A36) generates one wait state to allow sufficient time for the read/write operation.

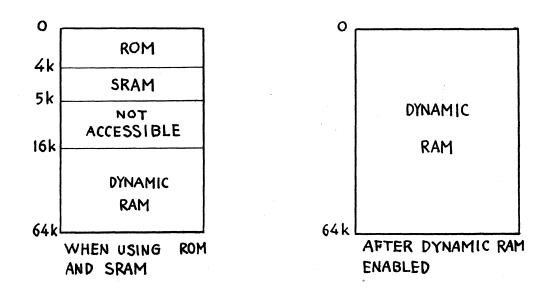


Figure 4 Memory Space Description

4.3 Eight DRAM chips (64Kb each) are used for main memory. Two 74S74 (A35,A37) and other related chips (A18, A19) are used to generate -CAS and -RAS signals using -M1 and -MREQ signals. Select column signal chooses either A0-A7 or A8-A15 as row and column address. As explained previously, when dynamic RAM is enabled it activiates -CAS signal so that all DRAM locations can be accessed by the user. During refresh, only -RAS signal is activated and one whole row is refreshed at one time.

The following diagrams show the timing of those signals.

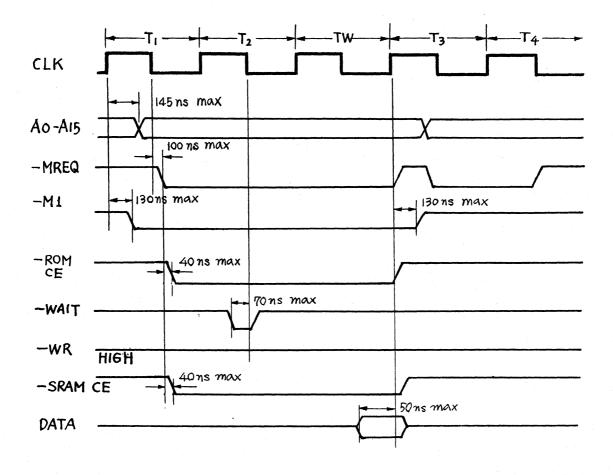
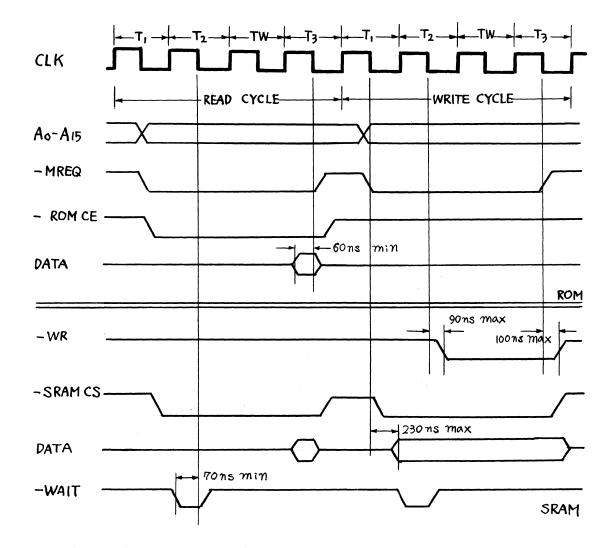
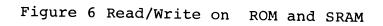
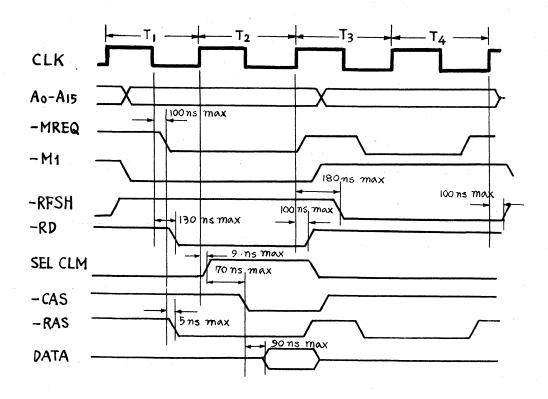
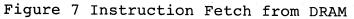


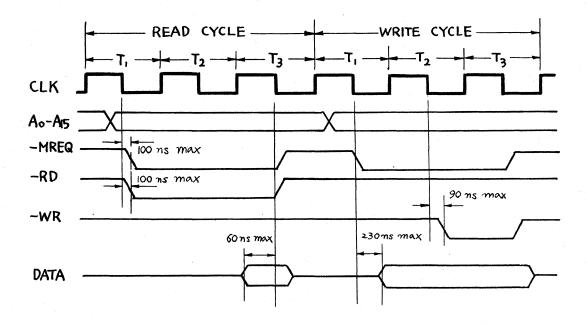
Figure 5 Instruction Fetch from ROM/SRAM

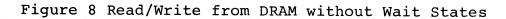










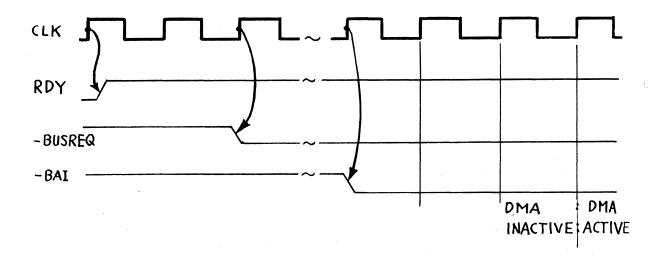


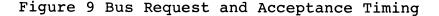
4.4 DMA Operation

The RDY line is monitored by the DMA to determine when a peripheral device is ready for read or write operation. When the DMA is enabled the RDY line indirectly controls DMA activity. -BUSREQ is used as output from the DMA chip to request for control of the address, data and control bus from the CPU. If CPU receives active -BUSREQ, the CPU will set these buses to high impedance state as soon as the current CPU machine cycle is terminated, and then sends out -BUSAK signal to indicate that DMA can control these signals.

Figure 9 illustrates the -BUSREQ and -BUSAK (BAI) timing. The RDY line is sampled on every rising edge of CLK as a level, not an edge. When the DMA detects a low on -BAI line for two consecutive rising edges of CLK, the DMA begins transferring data on the next rising edge of CLK.

Figure 10 and Figure 11 explain the timing diagram of RDY line in burst mode and in continuous mode of the DMA.





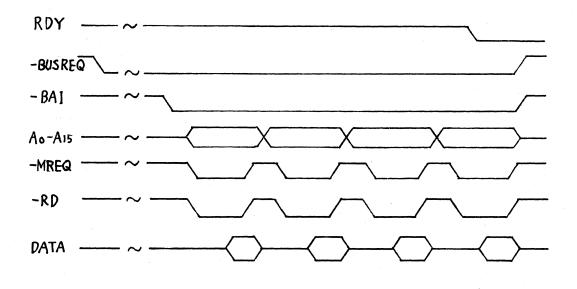


Figure 10 RDY Line in Burst Mode (Refer to I/O Operation for precise timing.)

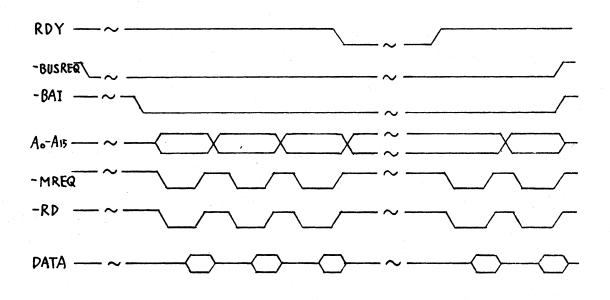
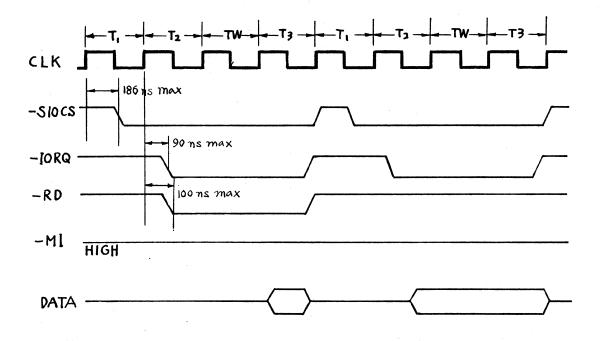


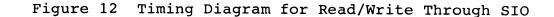
Figure 11 RDY Line in Continuous Mode (Refer to I/O Operation for precise timing)

4.5 I/O Operation

The Z80A CPU can address up to 256 I/O ports. The CPU uses only the lower 8 address lines for I/O addressing. The address lines A3 to A7 are decoded through two 74LS138 (A58,59). For I/O read/write, the -IORQ signal "NANDED" with RD/WR is used to provide -IORD/-IOWR signal. During the I/O operation, the CPU automatically inserts a single wait state (TW). This extra wait state allows sufficient time for an I/O port to complete a read/write operation.

One SIO (A82) is for RS232 interfaces and its related chips are 75188 (for output) and 75189 (for input). The other 3 SIO's (A61, A83, A84) are for RS422 interfaces with the user system; the related chips are 26LS31 (for output) and 26LS32 (for input). The timing of the signals generated by the CPU input instruction (to read data or the status byte) and output instruction (to write data or control bytes) are illustrated in Figure 12.





4.5 I/O Operation

FD 1793, WD1691, WD2143-01 and other related chips are used for floppy disk control. FD 1793 has two modes of operation according to the state of -DDEN (Pin 37). When -DDEN=1, single density is selected. The CLK input (Pin 24) is set at 1 MHz for mini floppy. DRQ line indicates that the data register contains assembled data in the read operation, or the data register is empty in the write operation. This signal is reset when serviced by the system through reading or loading the data register.

Figure 13 and 14 show timing diagram of those related signals.

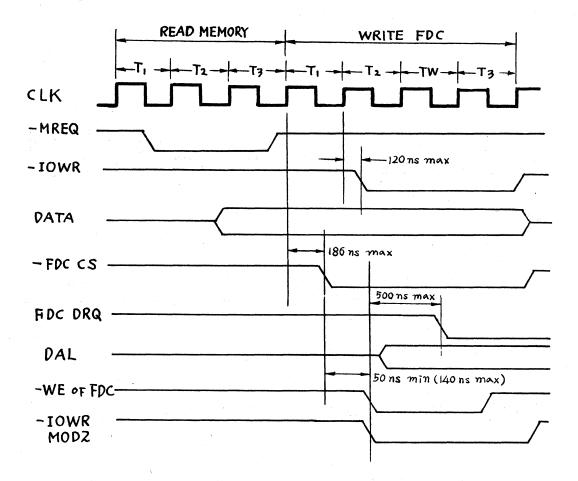


Figure 13 Timing Diagram of Memory to FDC Transfer (DMA Write Operation)

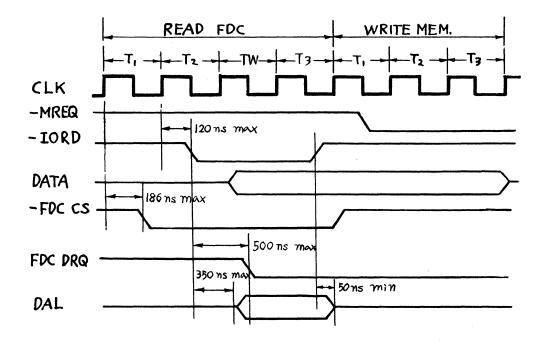
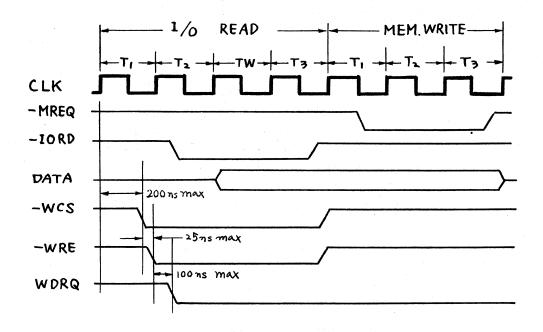
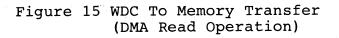


Figure 14 Timing Diagram of FDC to Memory Transfer (DMA read operation)

WD 1000 controller board is used for $5\frac{1}{4}$ " Winchester disk interface. As with the floppy disk controller, -WDC DRQ signal is used to control the RDY line of DMA by being "NANDED" with other I/O RDY signals through 74S133 (A71). -IORD signal decides the direction of data on 74LS245 (A91) between the system and controller. Figure 15 and 16 shows the timing diagram in WDC read/write operation using DMA.





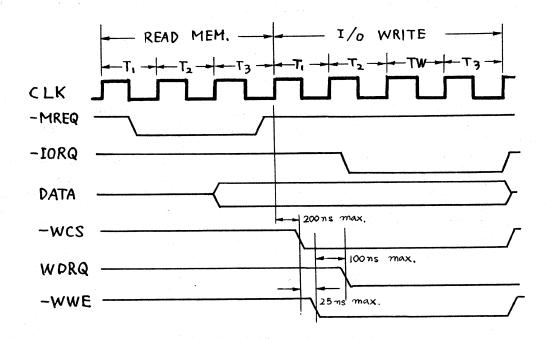


Figure 16 Memory to WDC Transfer (DMA Write Operation)

5. Connector Description

All the connector position is shown in Figure 17 and Table 3.

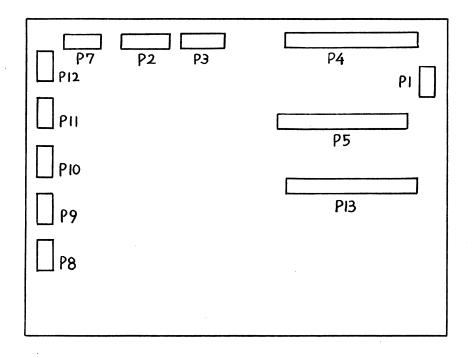


Figure 17 Connectors on the board

| Connector # | Description |
|-------------|--------------------------------------|
| P1 | Power |
| P2 | RS232C Service Terminal Interface |
| P3 | RS232C Serial Interface |
| P4 | Centronics-type Printer Interface |
| P5 | Winchester Disk Controller Interface |
| P6 | Reset, Power On Light |
| P7 | RS422 User Channel 1 |
| P8 | RS422 User Channel 2 |
| P9 | RS422 User Channel 3 |
| P10 | RS422 User Channel 4 |
| P11 | RS422 User Channel 5 |
| P12 | RS422 User Channel 6 |
| P13 | Floppy Disk Drive Interface |

Table 3. Connector Assignment

5.1 Power Connector (P1: 5 Pin)

| <u>Pin No.</u> | Description |
|----------------|-------------|
| 1 | -12 V |
| 2 | Unused |
| 3 | GND |
| 4 | +5 V |
| 5 | +12 V |

Table 4 Power Connector Configuration

5.2 RS232C Connector (P2,3: 25 Pin)

| <u>Pin No.</u> | Description (DCE P3) | |
|------------------------|---|--|
| 1 2 3 | Frame Ground Transmit Data (Receive Data P3) Receive Data (Transmit Data P3) | |
| 4 5 7 8 20 | Request to Send Clear to Send Signal Ground Data Carrier Detect Data Terminal Ready (Printer Busy | |

Polarity of Data Signals: - True Polarity of Control Signals: + True

Table 5 RS232C Connector Configuration

5.3 Floppy Disk Controller Connector

| <u>Pin No.</u> | Description |
|----------------|---|
| 6 | -Drive Select 3 |
| 8 | -Index/Sector |
| 10 | -Drive Select 0 |
| 12 | -Drive Select 1 |
| 14 | -Drive Select 2 |
| 16 | -Motor On |
| 18 | -Direction Select: - Toward Center + Off Center |
| 20 | -Step |
| 22 | -Composite Write Data |
| 24 | -Write Enable |
| 26 | -Track 00 |
| 28 | -Write Protected |
| 30 | -Composite Read Data |
| 32 | -Side 1 Select (Ignored by one side floppy disk drive) |

Pin 1, 2, 3, 4, and 34 are not used. All other odd number pins are grounded.

Table 6. FDC Connector Configuration

| <u>P</u> : | in No. | Description |
|------------|-------------------------------|-------------------------------|
| | 1 | WDAL 0 |
| | 3 | WDAL 1 |
| | 5 | WDAL 2 |
| | 7 | WDAL 3 |
| | 9 | WDAL 4 |
| | 11 | WDAL 5 |
| | 13 | WDAL 6 |
| | 15 | WDAL 7 |
| | 17 | WA O |
| | 19 | WA 1 |
| | 21 | WA 2 |
| | 23 | -WCS (Winchester Chip Select) |
| | 25 1 1 1 1 1 1 1 1 1 1 | -WWE (Winchester Read Enable) |
| | 27 | -WWAIT |
| | 29 | WINTRQ (Interrupt Request) |
| | 35 | WDRQ (Data Request) |
| | 37 | -WMR (Master Reset) |
| | 39 | |

Pin 31, 33 are not used. All even number pins are grounded.

Table 7. WDC Connector Configuration

5.5 Centronics-type Printer Interface Connector (P4: 40 Pin)

40 pin right angle header is used on the board, so the pin configuration on the board is different from the 36 pin connector on the rear panel but it will be matched when they are connected together.

Refer to Figure 18 for comparison of two connector pin assignment between 40 pin connector on the board and 36 pin connector on the rear panel which will be hooked up to the printer.

Table 8 describes the pin configuration of the connector on the board.

| CONNECTOR ON THI | BOARD | CONNEC | TOR ON | THE REAR | PANN | IEL |
|------------------|-------|--------|--------|----------|------|-----|
| 39 37 | 5 3 1 | 18 17 | | 3 | 2 | 1 |
| 40 38 | 6 4 2 | 36 35 | , | 21 | 20 | 19 |

Figure 18. Comparison of Two Connectors for Printer

| <u>Pin No</u> | n de la contra de l ■ ■ tra de la contra d | | | Description | |
|--|--|--|--------------------------------------|--|---|
| Rear Panel Connector | Board Po | ort/Data Bi | t | | |
| 1, 19* 2, 20 3, 21* 4, 22* 5, 23* 6, 24* 7, 25* 8, 26* 9, 27* 10, 28* 11, 29* 12 13 31, 30* 14 32 15 | 1, 2* 3, 4* 5, 6* 7, 8* 9, 10* 11, 12* 13, 14* 15, 16* 17, 18* 19, 20* 21, 22* 23 25 26, 24* (W6) 27 28 29 | DØ/Ø D1/Ø D1/1 D1/2 D1/3 D1/4 D1/5 D1/6 D1/7 BSTB/P10 DØ/5 DØ/4 DØ/6 DØ/1 | (W5) (W4) (W8) (W7) (W3) | -Data strobe Data 0 Data 1 Data 2 Data 3 Data 4 Data 5 Data 6 Data 7 -Ackn1g Busy ** PE (paper emp *** Select -Input Prime Gnd *** -Fault Unused | Output Output Output Output Output Output Output Output Input Input Input Input Input Input Input |
| 33 16 | 30 31 | DØ/2 | (W2) | **Light Detect GND | Input |
| 34 17 35 18 36 | 32 33 34 35 36 | DØ/7 | (W9) | LC Pulse Gnd Unused Unused Unused | Output |

* Second pin number indicates return ground.

Pin 37, 38, 39, and 40 are not connected.

**Ground if no connection
***Tie high if no connection

Table 8: Configuration of parallel Printer Interface

5.6 User system connector: RS422 (P7: 15 pin, P8-P12: 16 pin)

There are two types of connectors for RS422 on the system board. User 1 connector is D type 15 pin (P7) but from user 2 to user 6 (P8 to P12) connectors are 16 pin header type but these are connected to the D type on the rear panel eventually. The difference is only due to the different numbering order between those two types. Refer to Figure 19 for comparison.

| D Type | Header | Туре |
|------------------------------------|--------|--------------------|
| 8 7 6 5 4 3 2 15 14 13 12 11 10 | | 7 5 3 1 8 6 4 2 |

Figure 19 Two Different Type Connector Pin Order for RS422

a. P7 (D type: 15 Pin)

| Pin No. | Description | Pin No. | Description |
|---------|-------------|---------|--|
| | | | |
| | | | $(e_{k+1}, e_{k+1}) \in \mathbb{R}^{n} \to \mathbb{R}^{n}$ |
| 1 | Shield | 9 | -TXD |
| 2 | TXD | 10 | -RXD |
| 3 | RXD | 11 | -RTS |
| 4 | RTS | 12 | -CTS |
| 5 | CTS | 13 | TXC |
| 6 | -TXC | 14 | RXC |
| 7 | -RXC | 15 | Reserved for Manufacturer |
| 8 | Signal GND | | |

Table 9. D Type Configuration for RS422

b. P8 to P12 (header type: 16 pin)

| <u>Pin No</u> | · . | Description | | | Description |
|---------------|-------|-------------|--------------|-------|-------------|
| "D" Conn. | Board | | "D" Conn. | Board | |
| 1 | 1 | Shield | 9 | 2 | -TXD |
| 2 | 3 | TXD | 10 | 4 | -RXD |
| 3 | 5 | RXD | 11 | 6 | -RTS |
| 4 | 7 | RTS | 12 | 8 | -CTS |
| 5 | 9 | CTS | 13 | 10 | TXC |
| 6 | 11 | -TXC | 14 | 12 | RXC |
| 7 | 13 | -RXC | 15 | 14 | Unused |
| 8 | 15 | Signal GND | N/C | 16 | Unused |

Table 10. Header Type Configuration for RS422

5.7 Reset Power On Light Connector (P6: 3 Pin)

| <u>Pin No.</u> | Description |
|----------------|--------------------------|
| 1 | GND |
| 2 | Reset SW (normally high) |
| 3 | Reset SW (normally low) |

Table 11. Reset, Power On Light Connector Configuration

5.8 LED Description

4 LED's are used for diagnostic purposes. They are lit when the data line goes low. The corresponding port address is 00 the data byte assigned to each LED is as follows.

| LED No. | Data Byte |
|---------|-----------|
| 1 | FO |
| 2 | F1 |
| 3 | F2 |
| 4 | F3 |

The LED sequence for each test is as follows:

| Lit LED | Test In Process |
|---------------|---------------------------|
| 0 0 0 🕊 1 | MEM Test |
| o o 🕊 o 2 | DMA |
| o 🍟 o o 3 | WDC |
| Ψοοο4 | SIO |
| o o 📽 📽 1, 2 | PIO |
| ₩₩003,4 | FDC |
| ₩₩₩1, 2, 3, 4 | After finishing all tests |

6.0 Description for Jumpers

| W1: 1) Connect B and C when used with service terminal. 2) Connect A and B when used without service terminal. | | | | |
|---|--|---|--|--|
| W2-W9: | Use as follows depending on th | ne type of printer. | | |
| Jumper | Related Signal and Pin No. (Centronics-type) on rear panel | Description | | |
| W2 | Light detect (pin 33) | Connect when the printer interface does not need this signal. | | |
| W3 | -Fault (pin 32) | Connect when the printer interface does not need this signal. | | |
| W4 | Select (pin 13) | Connect when the printer interface does not need this signal. | | |
| W5 | Paper empty (pin 12) | Connect when the printer interface does not need this signal. | | |
| W6 | -Input prime return (pin 30) | Connect if GND is needed. | | |
| W7 | GND (pin 14) | Connect if GND is needed. | | |
| W8 | -Input prime (pin 31) | Connect when the printer interface needs this signal. | | |
| W9 | LC pulse (pin 34) | Connect when the printer interface needs this signal. | | |

7.0 POWER REQUIREMENTS

| | Sys. Bd | WDC | Hard Disk | Floppy Disk |
|------|------------|------------|------------|-------------|
| +5V | 5.5 A Max. | 2.5 A Max. | l.0 A Max | 0.5 A Max. |
| | 2.5 A Typ. | 2.1 A Typ. | .75 A Typ. | 0.45A Typ. |
| +12V | 80 mA Max. | | 5.0 A Max. | 2.2 A Max. |
| | 50 mA Typ. | | 1.5 A Typ. | 0.8 A Typ. |
| | | | | |
| -12V | 50 mA Max. | 50 mA Max. | | |
| | 30 mA Typ. | 25 mA Typ. | | |

8.0 WINCHESTER DISK CONTROLLER

8.1 Overview

The WD1000 hard disk controller is a discrete implementation of all the functions required to control SA1000/ST506-compatible Winchester (hard) disk drives via a standard data and control BUS. The controller is fabricated using a mix of high-speed bipolar and NMOS devices contained on a single, two-sided PC board.

The design of the circuitry makes use of a high-speed Microcontroller called the 8X300, newly developed NMOS support devices, Schottky devices, and low power Schottky devices. Together they work to achieve low component count and cost while maintaining high performance and reliability.

All I/O connections are made using standard ribbon cable connectors. Standard pin-out configurations for disk interface connectors are provided to permit direct pin-for-pin connection to ST506 compatible 5" drives.

All power for the board can be supplied from a single 5-volt power supply on a separate connector.

All host-to-disk data transfers are buffered by onboard RAM to achieve totally asynchronous transfers to and from the disk by the host.

The disk controller is built around five basic sections:

- 1. Processor functions
- 2. Serial data separation
- 3. Data conversion and checking
- 4. Serial data generation
- 5. Host interface functions

8.2 Processor Functions

All functions of the WD1000 controller are ultimately controlled by the onboard processor. Due to the high data rates associated with hard disk drives, processing of data and control of machine functions within the circuitry requires a processor capable of extremely fast execution speed. The processor used is the 8X300, a bipolar microcontroller particularily well suited for handling data efficiently at high rates.

The 8X300 operates at a basic clock rate of 8 MHz and performs all operations within 2 clock cycles giving it a speed of 4 MIPS (Million Instructions Per Second) or one instruction executed every 250 nS. The architecture of the processor is different from most other microprocessors in that no common data or address BUS is provided to be shared by RAM, ROM or peripheral devices.

Instructions are fetched from ROM via a dedicated instruction address and data BUS. The instruction address BUS (IAO-IA13) is capable of directly accessing 8K words of program storage. However, the WD1000 uses only the first 10 address lines, limiting onboard program storage to 1K words.

Program data is input to the 8X300 (U50) on the Instruction Data Operation. All BUS designations utilized by the 8X300 are reversed from the traditional LSB to MSB weighting. These bus lines have all been renamed on the schematic of the WD1000 to provide a more conventional designation system for the board.

8.2.1 Fast I/O Select

An extension byte has been added onto the instruction data memory to provide port access decoding on an instruction-by-instruction basis. This "Fast I/O Select" byte is not processed by the 8X300. It is decoded by auxiliary hardware to provide eight read strobes and eight write strobes which route data to the various devices distributed along the interface vector BUS.

The Fast I/O byte is latched into a 6-bit latch (U27) on the trailing edge of MCLK to ensure the data remains stable during the entire instruction. This data selects a read strobe and write strobe through two 1-of-8 decoders (U20 and U26) which are alternately enabled by the -WC control strobe produced by the 8X300. To provide edges on read strobes during sequential read operations from various ports, the red strobe decoder (U20) is always disqualified at the end of instruction by +MCLK' which is a delayed copy of +MCLK. This delay compensates for timing races through the Fast I/O latch (U27) and the control signals.

Because each decoder has a unique input, it is possible to select any read port with any write port during each instruction. Data is transferred between the processor and its ports on a separate 8-bit BUS called the "I/O" BUS. This BUS is active low. It must be noted that this BUS is in no way related to the instruction data BUS and can be thought of as simply an 8-bit bidirectional I/O BUS of the 8X300. In fact it has been renamed as I/O0-I/O7 to reflect this distinction.

8.2.2 Internal BUS Control

Several BUS control signals are produced by the 8X300 to identify and strobe the data in the I/O BUS. WC (write control) is a signal which determines the direction of the data to and from peripherals. When WC is false (during the first half cycle), data is being input to the 8X300 from the I/O BUS. When WC is true (during the second half cycle), data is being output from the 8X300 onto the I/O BUS. SC (select control) is a signal which becomes active during the second half cycle instead of WC if the I/O BUS contains an 8-bit I/O address. In the WD1000, both WC and SC are combined by a NOR gate (U24) to indicate all accesses to any port. This arrangement allows 8-bit immediate-data moves from the 8X300 to any output port within one instruction, instead of the normal 5-bit immediate moves provided for by the instructions set.

All instruction fetches occur late in the second cycle of the preceding instruction. This time is marked by the generation of a 65 nS (nominal) active high pulse called MCLK which occurs every instruction. MCLK is used to latch data prior to being input on the I/O BUS to insure stability during reads. It is also used to disqualify any read strobes that would otherwise remain true during the second clock cycle of any instruction which does not write to a port.

Two additional BUS control signals are produced by the 8X300. They are RIGHT BANK SELECT (-RB) and LEFT BANK SELECT (-LB). These are not used, however, in the WD1000 due to the implementation of the Fast I/O Select logic.

8.2.3 Reset Circuit

The 8X300 is held in reset for approximately 40 mS after initial power on. This is accomplished by an RC network (R42, C68 and CR3) which drives a Schmitt trigger (U31) to provide a proper rise/fall time on the -RESET line of the 8X300. Alternate reset of the processor can be accomplished by dropping -MR (J5 pin 39) whenever the host wishes to reset the controller. A Schmitt trigger (U31) is provided with a 4.7K pull-up (R43) to buffer the -MR input from the host. -RESET also propagates to the drive control latches (U52 and U48) and their associated line drivers (U54-56) and host interface WAIT (U320, DRQ and INTRQ latches (U30), ensuring proper initialization of these functions during power-up and subsequent resets from the host.

8.2.4 Processor Power Supply

Power is supplied to the 8X300 from the +5 Volt (Vcc) power BUS. Due to the internal operation of the 8X300, an on-chip voltage reference is provided to produce bias to an external pass transistor (Q4). This drops Vcc to the 8X300 to approximately +3.0 Volts. All signals into and out of the 8X300 are internally level shifted to be TTL compatible.

8.2.5 Read and Write Ports

Throughout the circuit, output ports are formed by D type latches using write strobes (WRO-7) to latch data into the ports. Reading of ports is universally accomplished by using read strobes (RDO, RD2, RD4-6) that enable selected tri-state output devices on the I/O BUS. Additionally, two read strobes are used to clock the host DRQ and INTRQ latches (U30) and one read strobe is left unused as a "dummy" port for instructions not requiring data from a port. This insures troublefree operation of the Fast I/O port decoders.

8.2.6 Read/Write Memory

Since the 8X300 does not permit data to be saved or retrieved from dedicated program storage, RAM must be installed on the I/O BUS. RAM must, therefore, be accessed via the I/O BUS by I/O instructions like all other port accesses. To provide for addressing the RAM, three latch counters (U40, 45 and 46) are connected to the I/O BUS to receive and store addresses required to access the RAM (U33 and U39).

8.2.6.1 RAM Addressing

The RAM address BUS (RAO-RA9) uniquely addresses 1 of 1024 memory locations. As each counter chip reaches a count of 0, it sets a borrow condition to the next higher counter which is decremented at the end of the next access to RAM. When all bits of the address have been reset, the -ROVF bit on the last counter (U40) is reset providing an overflow status which can be read by the processor on U43. By setting various beginning address values, -ROVF can be used to mark the end of any RAM access loop from 1 to 1024 bytes in length. In the WD1000, this function is used for setting sector buffer lengths of 128, 256, or 512 bytes.

8.2.6.2 Sector Buffering

All data read from the disk or written to the disk is passed through the RAM to provide buffering required for asynchronous data transfer between the host and disk. The counters are postdecremented which means that effective addresses are stable to the RAM by at least the instruction prior to the actual access. This preselection feature effectively reduces RAM access time to the output enable and propagation time of the RAM for read operations and the width of the minimum -WR strobe pulse for write operations.

8.2.6.3 RAM Accessing

RAM access is initiated by -RCS which is the logical OR (by U25) of -RDO and -WRO generated by the Fast I/O decoders (U20 and U26). Data to be read from RAM will be placed on the I/O BUS whenever -RCS is low and -WC is high. Data is written into a selected RAM cell on the trailing edge of -WC if -RCS is low. During writes, both -WC and -RCS will be low for at least 120 nS so that data setup time requirements are met.

8.2.6.4 Scratchpad Operations

Because the RAM address counters can be preset, direct reads and writes to a specific address are possible. This function is used for scratchpad storage during program execution. This mode of RAM access requires 2 or 3 instruction cycles for each random access to the RAM as opposed to 1 for sequential access using the

post-decrement feature.

8.2.7 MAC Control Port

Basic control of the various functional sections of the WD1000 is accomplished by a dedicated 6-bit control port called MAC CNTRL (U34). MAC CNTRL enables the functions of the WAIT control circuitry (-WAEN), CRC generation (-CRCIZ), gating of read data into data separation circuitry (RGATE), selection of read or write functions (-WRITE), control of CRC check word output (-IBLA) and AM detection (SRCH). MAC CNTRL output states are latched into the port by a write strobe (WR7). Additionally, any time MAC CNTRL is loaded with a new byte, the lower two data bits (I/O0-1) are strobed into the upper two address counter/latch bits (RA8,9).

All remaining ports are distributed among the basic functional sections of the WD1000 and are described in detail within the discussion of those functions.

8.3 Serial Data Separation

The WD1000 controller contains on-board circuitry to process incoming MFM data from the drive by a process called data separation. Here, some background information may be helpful:

In order to provide maximum data-recording density for maximum storage efficiency, data is recorded on the disk using a Modified Frequency Modulation (MFM) technique. This technique requires clock bits to be recorded only when two successive data bits are missing in the serial data stream. This reduces the total number of bits required to record a given amount of information on the disk. This results in an effective doubling of the amount of data capacity, hence the term "double density".

The fact that clock bits are not recorded with every data bit cell requires circuitry that can remain in sync with data during the absence of clock bits. Synchronous decoding of MFM data streams requires the decoder circuitry to synthesize clock bit timing when clocks are missing and synchronize to clock bits when they are present. This is accomplished by using a phase-locked oscillator employing an error amplifier/filter to sync onto and hold a specific phase relationship to the data and clock bits in the data stream. The synthesized clock called RCLK can then be used to separate data bits from clock bits and to shift the resultant serial data into registers for parallelization into bytes.

8.3.1 Incoming Data Selection

In the WD1000, serial data is input from up to four radiallyconnected drives via a quad RS 422 differential receiver (U15). The receiver converts differential input data to TTL levels for use by the controller. The data from the selected drive is then routed to the data separation circuitry by a four-section AND/OR/INVERT gate (U23). At this point data and clocks are still combined and appear as 50 nS (nominal) active high pulses spaced at intervals of 1, 1.5, or 2 times the RCLK period. This data is presented to the input of another AND/OR/INVERT gate (U18) which will gate either MFM data or a reference clock into the first stage of the VCO error amplifier circuitry.

8.3.2 Reference Clock

The reference clock is derived from the write clock crystal oscillator (Q3, U17 and associated circuitry). This oscillator uses a fundamental crystal cut to oscillate at 4 times the RCLK frequency. The 4X output is then divided by U17 to produce both a 2X clock (2XDR) which is used as a reference and a 1X clock (WCLK) which is used to produce MFM write data for the disk. The crystal (Y1) frequency is 20.000 MHz for ST506 compatible drives or 17.360 MHz for SA1000 compatible drives.

8.3.3 Clock Gating

The gating of the reference and MFM data into the data separator is dependent on the condition of the read gate (RGATE) signal and the spacing of the data on the serial stream after RGATE is brought true. The techniques used to separate data from clocks make it necessary to run the VCO at a rate twice the data clock (RCLK) rate. The VCO is therefore set to an open loop frequency of 2X RCLK. Any variations in this rate due to variations in disk rotational speed must be compensated for by the VCO. Instantaneous shifts in data due to the effects of adjacent bit cells on the disk and minor noise must be ignored. Also, the response of the VCO must be adjusted to effectively override missing clock bits which occur as a result of the MFM recording technique.

The resultant compromise between response and reject requirements of the VCO causes the VCO to have a tendency to become locked onto harmonics of the data rate rather easily. This is likely to occur if the VCO is connected to a data stream over a field of data which has data bits spaced at 1.5 or 2 times the actual RCLK time intervals.

To prevent this from occurring, the VCO is always held locked onto a stable clock running at 2 X RCLK frequency whenever the controller is not actually reading data. Read data is switched to the VCO error detector only when it is known that the data stream frequency is equal to the RCLK frequency. This can occur only when the data is a solid stream of all ones or all zeros.

8.3.4 High Frequency Detector

The switching function is initiated immediately after the RGATE goes true. It will only switch read data into the VCO after 16 consecutive ones or zeros (high frequency) are detected by a one-shot (U10) and counter (U12) connected directly to the raw MFM data.

The one-shot is adjusted for a pulse width of 1.25 times the RCLK period. This is a 250 +/-10 nS for ST506 compatible drives and 287 +/-10 nS for SA1000 compatible drives. These adjustments of the DRUN one-shot (U10) provide tolerance of up to 1/4 RCLK period in jitter on the MFM data bits while still being able to distinguish MFM zeros or ones from other data patterns.

Each clock or data bit on the serial stream triggers the oneshot. If the time between successive triggers is less than the one-shot time constant, the one-shot remains retriggered.

As the one-shot is triggered by data stream bits, so is the up/down counter (U12), whose count mode is controlled by the state of the one-shot outputs. While the one-shot is being retriggered, the counter counts up.

When any data bit fails to reach the one-shot before its time constant is over, the one-shot resets and in turn clears the counter. Only when 16 successive retriggers occur can the counter reach its terminal count. At this time, the counter overflow goes true and sets up the -DRUN latch output (U13 pin 8) low which switches read data in and reference clock out. An AND/OR/INVERT gate (U18) performs the switching. -DRUN is read through U42 by the 8X300 to determine the condition of the MFM data stream.

The data and clocks are now connected to the first stage of the data separator. The heart of the data separator is the VCO (U2 and associated circuitry) and the error amplifier and filter (U1). As previously stated, the VCO runs at a frequency twice that of the RCLK rate.

8.3.5 VCO

RCLK is produced by the VCO through a divide-by-two counter (U8). The VCO is a discrete LC oscillator with a shunt capacitor formed by a hyper-abrupt tuning diode (CR1). The diode varies its capacitance in accordance with the amount of reverse DC bias applied to its PN junction. As the bias decreases, the capacitance increases, pulling the oscillator down in frequency. Conversely, as the bias increases, the oscillator frequency rises.

The VCO performs conventionally with one exception. To help the VCO lock onto the incoming signals more quickly, an external timing signal freezes the output of the VCO in the high state. This is done by disqualifying U2 in the VCO feedback circuit and by removing bias from the transistor (Q2) which provides loop gain in the oscillator. -PHASEUP performs this function and is present just after the switch over from reference clock to MFM data is made.

The width of -PHASEUP is directly related to the difference in timing between the positive going transition of the VCO output and the positive transition of the second data/clock bit of the MFM data stream. -PHASEUP causes the output of the VCO to rise in phase with the MFM data from the drive. This allows the VCO to adjust its phase slightly to center data/clocks within the RCLK rather than performing a frequency acquisition to lock onto the data stream. The phase acquisition is much faster and easier to achieve and results in vastly improved performance.

8.3.6 Error Amplifier

The error amplifier and filter control the VCO. The error amplifier is a balanced differential amplifier whose output either sources or sinks current to the filter stage. The error amplifier output is pulse-width modulated by the phase detector (U6, 7).

Whenever the VCO is running too slow, the error amplifier receives pulses from data bits before pulses from the VCO clock. This causes the error amplifier to produce pump-up pulses to the filter. The filter integrates these pulses producing an average increase in the voltage to the cathode of the hyper-abrupt tuning diode (CR1). This effectively increases the reverse bias on the diode which reduces its capacitance and therefore increases the VCO frequency slightly to match the phase of the incoming data.

Whenever the VCO is running too fast, the error amplifier produces pump-down pulses to the filter. The diode then receives decreased reverse bias, more capacitance, and lower VCO frequency.

The operating point of the tuning diode (CR1) is initially set for an open loop VCO frequency of two times RCLK by setting -OSC ADJUST and monitoring the VCO output. This adjustment places the initial bias through the bias divider (R18-20, R22 and C8-9) at approximately -2.8 V to -3.2 V. At this setting the most responsive region of the diode is being used giving higher gain in the VCO. To keep the initial bias voltage close to three volts with varying disk data rates, the VCO inductor (L1) is 3.3 UH for 5 MHz drives and 3.9 UH for 4.34 MHz drives.

The VCO is forced to match the phase of the incoming data. Once the VCO is close to the phase of the incoming data, the pump pulses become very small or are missing completely. It must be noted, however, that a slight error is always be present because, without pumps, the filter floats and the VCO drifts. The overall gain of the error amplifier and the VCO maintains this very small error, resulting in very close tracking between the VCO ouput phase and the incoming data phase.

As previously stated, great care is taken to insure that the VCO starts on the same phase as the incoming data. If this were not the case, the error amplifier would produce very large pumps in an attempt to pull the VCO onto frequency and phase. Due to the gain of the error amplifier and the required characteristics of the filter, the integrated pump pulses would over-compensate, causing the VCO to overshoot in its attempt to lock-on. This

action would continue in a diminishing fashion until lock-on occurred. Unfortunately, the data sync fields it was trying to acquire would be over by the time the VCO finally acquired lockon. -PHASEUP is, therefore, extremely important to the overall ability of the data separator to function reliably.

8.3.7 Sample on Phase Detector

The circuitry which feeds the error amplifier is called the Sample on Phase Detector. This circuit consists of several D latches (U7) and a delay line (DL1). The function of the circuit is to provide time windows during which the leading edges of the incoming MFM data can be compared to the leading edges of the VCO clock. These windows are approximately 50 nS in length. The windows are initiated by a leading edge of any data bit as it enters the detector. They are terminated by that same data bit, edge-delayed by a net 50 nS (60 nS in the delay line minus approximately 10 nS in propagation delays).

When both the delayed data bit and the nearest VCO edge arrive at the detector, the detector is reset until the next data bit arrives on the MFM data stream. The delayed data bit sets its half of the detector latches to produce a pump-up condition at the error amplifier. The VCO clock edge sets its half of the detector to produce a pump-down condition. When the circuit is balanced, either both pumps are on or both pumps are off, producing no net pump-up or down.

8.3.8 Window Extension

Once the VCO has been locked onto the phase of the incoming data, the actual separation of data and clocks can occur. This is done by using a technique called window extention. This technique causes data bits to first have their leading edges shifted into the center of the RCLK half cycles and then to be latched or extended until the next rising edge of the RCLK. The shift is accomplished by tapping the data off the sample on-phase detector delay line at the 60 nS tap and inverting the VCO clock to the RCLK divider (U8). The delayed data clocks a pair of latches (U9). The 'data' latch has its D input and CLEAR connected to +RCLK and the 'clock' latch has its D input and CLEAR connected to -RCLK.

If an MFM data bit enters the latches while -RCLK is high, it will be extended as a data bit. If -RCLK is high, it will be extended as a clock bit. Due to this extension technique, bits can jitter approximately 1/4 the RCLK period without being lost. The output of each latch is then further extended by feeding directly into another stage of latches (U3) and clocked on alternate edges of RCLK. The final outputs of the data extension/separation stage are two separate signals, one consisting solely of NRZ data, and the other NRZ clocks. NRZ data and clocks are finally in a form suitable for processing by subsequent circuitry within the WD1000.

8.3.9 Clock Detection

Due to the nature of MFM data encoding, it is impossible to know exactly if MFM bits are data or clocks. This ambiguity results in having to create circuitry to assume that bits on -RCLK are actually data bits until the VCO is locked on and a unique data/clock pattern is detected. This is accomplished by holding the VCO to RCLK divider (U8) reset until it is fairly certain that bits on the data stream are actually clocks belonging to a field of zero data.

Once this assessment has been made, the processor releases the AM detector (U14) by raising the SEARCH signal. This signal releases a latch (U13) which removes -DHOLD from the RCLK divider (U8) on the next rising edge of a MFM data bit so that CLOCKS will be on the -RCLK phase and DATA will be on the +RCLK phase. The processor makes its assessment of the state of the data stream solely on the occurrence of a significant run of zeros which are detected by the one-shot (U10) in the DRUN circuit. Once released, the phase of RCLK vs data and clocks remains stable throughout the read of an ID field or data field. Whenever SEARCH is dropped, the VCO to RCLK divider is once again reset and no RCLKS are produced.

8.4 Data Conversion and Checking

MFM data which has been separated to form NRZ data and clocks is processed through specialized circuitry to prepare it for parallel processing by the 8X300. This processing consists of 3 functional circuits.

- 1) AM detection (U14)
- 2) Serial to parallel conversion (U29)
- 3) CRC checking circuit (U19)

Each function is discussed separately but many interdependencies exist.

8.4.1 AM Detection

As previously stated, it is impossible to know whether serial data bits are actually data or clock bits by just looking at the data stream. It is also impossible to determine byte boundaries. This problem is solved by a uniquely recorded data/clock pattern called an Address Mark (AM). The AM consists of a data pattern of HEX 'Al' with a missing clock pattern of HEX 'OA'. Normally, a data byte of HEX 'Al' requires a clocking pattern of HEX 'OE'. In fact, due to the rules of MFM data encoding, an alternating clock pattern such as HEX 'A' or HEX '5' cannot exist legally.

The AM is used to uniquely identify the start of a field of information (Data or ID field) within each sector. Preceding each AM on the disk there is always a long run of "zero" data. Zeros have a clock bit for every RCLK. When attempting to read information from the disk, the WD1000 first acquires phase lock over a field of zeros. After this acquisition is achieved, the processor releases the AM detector (U14) by raising the SEARCH control line (SRCH) on the MAC CNTRL port (U34). Because of the circuitry associated with the VCO to RCLK divider, the -RDAT output of the data separator (U3 pin 8) will be high and the -CLKS output (U3 pin 6) will be low. -RCLK will be the shifting clock for -RDAT and +RCLK will be the shifting clock for -CLKS. These 4 signals are routed into the AM detector.

Inside the AM detector, the -RDAT is shifted into an 8-bit synchronous serial shift register and clocked on the falling edge of -RCLK. -CLKS are shifted into a similar shift register on the falling edge of +RCLK. The output stage of the -RDAT register is dumped into an 'Al' comparator and the output stage of the -CLKS register is dumped into an 'OA' comparator. AM detection occurs when both detectors are true, thereby making the relationship between data and clocks known. It is also known that data is being clocked by -RCLK so -CLKS can actually be discarded; their only purpose was in detecting AM. The -AMDET signal is used as a synchronization signal to start subsequent conversion circuitry. The -AMDET signal remains true until the processor again deasserts the SEARCH control line.

8.4.2 Serial to Parallel Conversion

After an AM has been detected, the Serial to Parallel converter (U29) takes over. NRZ data and -RCLK are used to shift data bits into an 8-bit serial to parallel shift register. As each bit is shifted, a divide-by-8 counter circuit is incremented. After every eight bit of data is shifted, the counter produces an overflow pulse, marking byte boundaries in the serial data stream. The overflow bit from the counter resets the counter, clocks the data from the shift register into an 8-bit parallel latch, and sets a tri-state flag register called BDONE. The flac can be read by the processor to see if any converted data is ready to be read from the latches.

When the processor sees BDONE in the true state, it services the device by gating data onto the I/O BUS using read strobe 4 (RD4) in conjunction with a tri-state buffer (U36). The act of reading the latches also clears off the pending BDONE flag. As successive bytes are processed, the BDONE is serviced by the processor as data becomes available.

Outputs from the serial to parallel device also include -SHFTCLK and -DOUT. -SHFTCLK is actually -RCLK propagated through the device. -DOUT and -SHFTCLK are tri-stated along with BDONE, and are active only when -WRITE is high, indicating a read mode of operation. -DOUT and -SHFTCLK are routed to the CRC generator checker device.

8.4.3 CRC Checking Circuit

Data recorded on magnetic media is prone to several types of errors which could render data unusable if some form of error detection were not employed. On the WD1000, a Cyclic Redundancy Check (CRC) is performed on all data transfers from the disk. The CRC is an error detection code consisting of 16 additional bits which are appended to every ID field and data field on the disk. These bits are produced by dividing the data stream serially with a large polynomial. This division produces a unique 16-bit value for any information passed through the CRC generator.

As data is being read from the disk, the CRC generator recomputes the original CRC bits. After the last two bytes (containing the original recorded CRC) are read, the value in the CRC generator must always be zero. When this happens, the data was correctly read and the controller will not flag an error. If, however, the CRC generator is not zeroed after it has checked all bytes of the recorded data, then the controller will flag the data as erroneous and enter into a retry condition. If after attempting to correctly read the data 16 times, the controller still cannot get correct data, the read is aborted and the host is informed that the data in the buffer is questionable.

The WS1000 uses the same device to generate and check CRCs for data being written and read on the disk. The polynomial used is:

X16+X12+X5+1 (commonly called the CRC-CCITT polynomial)

During read operations, the processor polls the condition of the DRUN circuitry. When DRUN is true, it begins to search for an address mark. Once the AM is located, the processor starts to read parallel data which has been converted from NRZ data by the serial-to-parallel device. The processor terminates this activity when it has received the information it is looking for or if an error is detected.

While the processor is reading the parallel data, the CRC generator is reconstructing the CRC check value. The CRC generator is initialized by the processor setting -CRCIZ low for at least 250 nS during the search for the AM. -CRCIZ is originated on the MAC CNTRL port (U34). Upon receiving the -CRCIZ signal, the CRC generator/checker presets all 16 of its internal polynomial division shift registers to logic ones. It also arms an internal latch which enables the checking function on the leading edge of the first non-zero data to enter the device.

Prior to an AM there is a field of zeros (all data bits low), so the first non-zero data bit into the device will always be the most significant bit of the AM (HEX Al).

Once enabled by the first non-zero data bit, the CRC device shifts succeeding data bits into a feedback shift register string with exclusive OR gates tied to the feedback nodes on the first, fifth, twelvth and sixteenth registers. As each RCLK occurs, the registers divide the incoming data and a unique pattern of ones and zeros appear across the registers.

When the last bit of an ID or DATA field is processed, the pattern in the registers should be equivalent to the 16 bits appended to the fields during original recording. The appended bits are also entered into the CRC device. If all of the bits in the appended field are identical to the bits in the registers, then the exclusive OR gates in the register string will have flipped all of the ones to zeros and the CRC will have been satisfied.

The output of each register stage is tied to a 16-bit wide comparator which goes true when all of its inputs are zeros. The output of the comparator is retimed to remove any decoding slivers and is output as CRCOK. The processor can read CRCOK through U43 to see if a CRC error occurred.

After the CRC bits are processed, the data stream contains at least one more byte of zeros. It is the nature of the CRC polynomial that if no bits are set to ones in the registers, then none are flipped if a constant input of zeros is shifted into the registers. This provides a convenient latching function for the CRCOK flag which remains true for at least 1 byte after the last CRC check byte, giving the processor time to read the flag.

The data, clock and BDONE are supplied to the CRC device on a 3bit minibus. During read operations, the Serial-to-Parallel device (U29) sources these lines because the WRITE control line from MAC CNTRL (U34) is low. This enables the tri-state drivers on these lines. The Parallel-to-Serial device (U37) has its tristate drivers disabled.

8.5 Serial Data Generation

The WD1000 records data on the disk in MFM format. To produce the proper data format, the WD1000 uses several specialized devices. These devices process the parallel data supplied by the host into a serial MFM data stream. The data supplied by the host is temporarily stored in the buffer RAM until the correct sector is located for the data to be written.

The process of writing is the opposite of reading except that the data separator circuitry is not required and the generation of the MFM data stream is produced by synchronous clocking techniques.

The functional sections of the serial data generation section are listed below:

- 1) Parallel to Serial conversion (U37)
- 2) CRC generation (U19)
- 3) MFM and precompensation (U30)

8.5.1 Parallel to Serial Conversion

Parallel data is converted into a serial NRZ data stream by the Parallel-to-Serial device (U37). The processor enables this conversion by lowering the -WRITE signal on MAC CNTRL (U34). -WRITE causes the tri-state buffers present on the parallel-toserial device to become active, supplying the CRC device with data, clocks, and BDONE strobes.

The processor presents parallel data on the I/O bus along with the -WR4 write strobe which latches the data into the parallel port on the trailing edge of the strobe. The write strobe also resets any pending BDONE. Inside the parallel-to-serial device, the parallel latches are loaded into a serial shift register on every eighth WCLK transistion. As the data is transferred to the shift registers, the BDONE status flag is set. The processor reads this flag through U43 to determine when to write the next parallel byte to the device. The timing of the parallel accesses is at a rate 1/8 that of the bit rate of the NRZ data stream. For ST506 compatible drives the byte timing is 1.6 us and for SA1000 drives it is 1.84 us.

The output of the last register in the shift string is brought out of the device as an NRZ serial data stream. The shifting clock is also brought out as SHFCLK to be used as the clock for the CRC device.

Whenever it is desired to write a repetitive string of identical data bytes, the processor can simply ignore the BDONE flag and permit the device to reload the data from its latches over and over again for as long as required to generate the field. This feature of the device is used in writing certain fields used in formatting.

8.5.2 CRC Generation

The CRC generator/checker (U19) is used to generate the CRC bits and to append them to the end of the data being written to the disk. This is the complementary function to that performed during reads. The operation of the polynomial generator is identical to read operations except that at the end of the data field, the processor sets a signal which causes the device to output the computed CRC after the data instead of reading the CRC and checking it.

The initial state of the shift registers within the device is forced to all ones by the processor pulsing -CRCIZ for approximately 250 nS while the parallel-to-serial device is outputting all zeros on the NRZ data line. At that time, a latch is set which holds the registers at ones until the first non-zero data bit enters the device. The first non-zero bit is the MSB of the AM (HEX Al) of the data field to be written. When the processor decides that enough zeros have been written to satisfy the sync field requirements, it stores a HEX Al in the parallelto-serial device. At the proper time (in sync with BDONE) the parallel-to-serial device begins to send the MSB of the AM to the CRC device. This starts the CRC polynomial generator and the CRC will be computed.

As the processor writes the last byte of data to the parallel-toserial device, it drops the -lBLA (1 Byte Look-ahead) signal on MAC CNTRL port (U34). This signal causes the CRC generator to begin dumping the computed CRC onto the NRZ data stream at the conclusion of the last date (synchronized with the BDONE signal). In this fashion, the device is able to append the proper CRC information to the end of a field of data. -lBLA is maintained at a low state for the duration of the unloading process which lasts for 16 bit times.

During the unloading process, the CRC registers back-fill with zeros. This feature is handy because by leaving -lBLA low for additional time, zeros will always be written after the CRC which is a requirement for the proper operation of the CRC device during read operations. The NRZ data with CRC appended is then sent to the MFM generator device.

8.5.3 MFM Generation

The conversion from NRZ write data to MFM write data takes place in the MFM/Precomp device. This device accepts NRZ data and a complimentary WCLK. It also produces MFM data and clocks by sending the data through circuitry which decides when and where to write clocks on the data stream under the MFM encoding rules. The proper encoding of the data into MFM requires the device to apply three rules to the data.

- If the current data cell contains a data bit, then no clock bit is generated.
- If the previous data cell contained a data bit, then no clock bit is generated.
- 3) If the previous data cell and the present data cell are vacant, then a clock bit in the current clock cell is produced.

The terms "data cell" and "clock cell" are defined by the state of the WCLK. While WCLK is low, it is a data cell. While WCLK is high, it is a clock cell. Therefore, both clock and the data cells are 1/2 the period of WCLK or 100 nS for ST506 compatible drives and 115 nS for SA1000 drives. Also, a clock and data bit can never occur within the same WCLK period and legal spacings for bits can be 1, 1.5, or 2 times the WCLK period only.

These rules are implemented within the device by shift registers that hold the next two last and present data bits and combinatorial logic. The state of WCLK is considered and the appropriate bit cells are filled and combined on the MFMW output line of the device. This line is subject to decoding slivers, so it is run through a re-timing latch (U21) to clean it up.

8.5.3.1 Write Precompensation

The MFM data stream is now totally compatible with the recording rules and may be sent to suitable line drivers for transmission to the drive except for one modification. Due to the decreasing radius on the physical surface of the disk, the inside tracks have less circumference and therefore exhibit an increase in recording flux density over the outside tracks. This increase in flux density aggravates a problem in magnetic recording known as 'dynamic bit shift'.

Dynamic bit shift comes about as the result of one bit on the disk (a flux reversal) influencing an adjacent bit. The effect is to shift the leading edge of both bits closer together or further apart than recorded. The net result is that enough jitter is added to the data recorded on the inside tracks to make them harder to recover without error. "Write precompensation" is a method which can be applied to reduce the effect of this shift on the data.

Precompensation is a way of predicting which direction a particular bit will be shifted and then intentionally writing that bit out of position in the opposite direction to the expected shift. This is done by examining the next two data bits, the last and the present bits to be written, and producing three signals depending on what these bits are. The three signals are EARLY, LATE and NOMINAL. They are used in conjunction with a delay line to cause the leading edge of a data/clock bit to be written early, late, or on time. As with MFMW, these signals are subject to decoding slivers and must be retimed by U21.

The processor can enable or disable the generation of these signals by controlling the RWC (Reduce Write Current) line from the MAC CNTRL port (U34). When RWC is high, precomp is in effect. When RWC is low, no precomp is generated and the NOM output of the device is held true.

The delay line actually performs the precomp with the help of an AND/OR/INVERT gate (U22). MFMW pulses are applied to the input of the delay line and, depending on which of the three precomp signals is present, the AND/OR/INVERT gate selects a different tap on the delay line. Nominal data is tapped from the second tap, early data from the first tap, and late data from the third tap. From the AND/OR/INVERT gate, the MFMW data is sent to the input of an RS 422 driver (U16) where it is converted to a differential form and then is sent to the drive.

The AND/OR/INVERT gate has one other function. If the controller is not writing, the WGI (Write Gate Internal) signal is low. This is inverted by U38 and applied to the AND/OR/INVERT gate's fourth section. This resulting high input effectively inhibits the gate from accepting MFMW data.

8.6 Host Interface

All data transfers between the host and the WD1000 take place over an eight-bit bi-directional bus (J5) consisting of eight Data Access Lines (DALO-7). The source or destination register inside the WD1000 is selected by the three address lines (A0-2). All accesses to the WD1000 are controlled by Card Select (CS-), Read enable (RE-), or Write Enable (WE-). Since the access time for any particular read or write operation varies, the WD1000 provides a not-ready signal (WAIT-). For systems using interrupts and/or DMA, the WD1000 provides INTerrupt reQuest (INTRQ) and Data ReQuest (DRQ).

Accessing the WD1000 is like accessing variable speed RAM. The host must provide a valid address in A0-2 along with a CS-. Immediately or after a short set-up time, the host may assert REor WE-. If access time on the WD1000 is over 100 nS, then WAITis asserted. The host <u>must</u> keep all address lines and strobes stable while WAIT- is true. When the WD1000 de-asserts WAIT-, the data has been accepted on a write or the data is on the DAL bus on a read.

8.6.1 Wait Enable

Since most of the registers in the WD1000 are not implemented in hardware, it takes the 8X300 a finite amount of time to fetch the requested data on a read or to store data on a write. This time varies depending on the amount of processing the 8X300 must do to access the desired register. After the data has been written or read, the WD1000 de-asserts the WAIT- line, allowing the host to terminate the current bus cycle.

The generation of the WAIT- signal is controlled by a bit in the MAC latch (U34) called WAit ENable (WAEN- will be asserted). On each bus access, the host <u>must</u> drop the Card Select (CS-) line on J5. The leading edge of CS- clocks the wait control latch (U32), transferring the WAEN- state through the latch, qualifying the wait drivers (U44, 54). This clocking action is required to ensure that WAIT- will not be asserted in the middle of any bus access already in progress. After the wait latch has been clocked, CS- (BIC or BOC in some installations) causes WAIT- to be asserted to the host.

The WAIT- line is released on the trailing edge of any read or write strobe to the communications latch (U49). This release is caused by the logical OR of RD6- and WR6- on U25 which presets the wait latch (U32) to a non-wait request condition. The WAITsignal is stretched to the trailing edge of the RD6- or WR6- by U2.

If WAEN- is de-asserted, the WD1000 generates no waits at all. In this case, the host reads the dummy status written to the communications latch by the 8X300. This feature is used by the microcode to simulate a busy condition when the host reads the status register in non-interrupt driven systems. When the WD1000 becomes un-busy, the WAEN- line is asserted and operations on the host interface bus are monitored once again.

8.6.2 Bus Gating

During all accesses by the host, one of two signals are produced to gate the bus. During read operations, CS- and RE- are ANDed, producing Bus Output Control (BOC-). This signal gates the contents of the communication latch (U49) onto the DAL bus. During write operations, CS- and WE- produce Bus Input Control (BIC). This signal latches the state of the DAL lines into an internal R/S latch.

8.6.3 Register Selection

The combination of a host read or write operation along with the WREQ- signal being asserted, generates a signal called the Card Select ACcess (CSAC). The 8X300 samples this signal at U43 every 250 nS, and, if asserted, reads the status of A0-2 and WE- (U43). The state of A0-2 and WE- determines which register is to be accessed (A0-2) and in what direction that access will take place.

8.6.4 Interrupts and DRQs

The WD1000 produces INTerrupt ReQuests (INTRQ) to signal the end of all disk operations and Data ReQuests (DRQ) to signal data ready to DMA controllers. INTRQ and DRQ originate on the MFM generator (U30) as an auxillary function of the chip. The WD1000 sets INTRQ using INTCLK- and sets DRQ using DRQCLK-, both of which are produced by U20.

Interrupts are cleared by HSAC- (Host Select ACcess) and AO, when the host reads the Status register, issues a command, or accesses the Sector Number register. DRQs are cleared when the host accesses the Data or cylinder Low registers. DRQs are re-issued for each byte to be transferred. HSAC- is a 200 nS version of the CSAC- signal. During Power-On Reset of Master Reset (MR-), INTRQ is set and DRQ is reset.

8.7 Maintenance and Adjustments

The WD1000 requires no scheduled preventive maintenance. There are a few adjustments associated with the data separation circuitry that may need to be adjusted if a drive with a different data rate is installed. Remember, the inductor Ll must be the proper value for the data rate being used and Yl must be selected for a fundamental frequency of four times the data rate.

8.7.1 VCO Adjustments

Data separation circuitry on the WD1000 uses a voltage-controlled oscillator (VCO) which phaselocks onto incoming data and provides a clock suitable for separating data and clock bits on an MFM encoded data stream. The VCO must be adjusted using the following procedures.

- 1. Ground the cathode of the tuning diode (CR1) to the closest accessible ground using a low inductance-shorting cable. This cable should consist of the shortest piece of wire possible to make the connection.
- Connect a frequency counter to the VCO buffered output on TP9.
- 3. Apply power to the board. Ensure a logic "1" exists at TP17. A logic "0" on TP17 inhibits the VCO and makes adjustments impossible. If a logic "1" is not present, verify that the DRUN circuitry (U10-U13) is functioning and adjusted properly. Refer to Section 8.7.4 for adjustments to DRUN.
- 4. Vary the OSC ADJ pot (R22) to verify that the range of adjustment values listed in the following table are obtainable. After the range has been verified, adjust R22 to the final setting listed in the table.

| Yl Frequency | Identity | Range | Final Setting |
|--------------|-------------|--------------|-------------------|
| 20.000 MHz | Ll = 3.3 uh | 9.0-11.0 MHz | 10.0 MHz +/-1 KHz |
| 17.360 MHz | Ll = 3.9 uh | 7.5-9.5 MHz | 8.68 MHz +/-1 KHz |

5. Turn off power to the board. Disconnect all test jumpers and test equipment.

8.7.2 Error Amplifier Adjustments (Static)

The phase detection technique used to correct the frequency and phase of the VCO employs a balanced sample and hold error amplifier. To ensure reliable operation of the data separator, the error amplifier must be properly balanced. The balance adjustment must be made using the following procedures.

- 1. Apply power to the board. Ground TP20 or U7 pin 4 to turn off the right half of the error amplifier (U1). In this state current only flows in the left half of the amplifier.
- 2. Connect a 100-ohm resistor between Ul pin 8 and ground.
- Adjust the BAL pot (R1) until a reading of 0 V +/- 20 mV is observed on TP5.

- 4. Remove the ground from TP20. Install a ground to the PUP1 signal line accessible on U7 pin 1. Connect Vcc to U7, pin 4 to turn off the left half of the error amplifier (the right half of the error amplifier will be on).
- 5. Verify a reading of 0 V +/- 50 mV on TP6. This indicates that the side-to-side balance of the error amplifier is within tolerance.
- 6. Re-adjust Rl until the reading at TP5 is 0 V +/- 20 mV.
- 7. Turn off the power to the board. Disconnect all jumpers and test equipment.

8.7.3 Error Amplifier Adjustments (Dynamic)

After the static balance adjustments are performed, the error amplifier should be adjusted for balanced dynamic operation.

This adjustment requires that the controller be constantly reading data on the innermost cylinder of a formatted drive.

- 1. While reading data, monitor TP9 (-OSC) with a frequency counter.
- 2. Adjust the error amplifier balance pot (Rl) until the most stable display reading is obtained. This indicates that the VCO is being locked on every attempt to read data.
- 3. Turn off power to the board.
- 4. Disconnect all test equipment.

8.7.4 DRUN Adjustments

To facilitate the process of acquiring phase-lock on data being read from a disk, a hardware detector is utilized to indicate when the read/write head of the drive is over a recorded field of all ones or all zeros. The detector depends on the timing of a one-shot (U10) which is adjustable by the DRUN pot (R26). R26 must be adjusted according to the following procedures:

The DRUN adjustment is made with the WD1000 in an operating test configuration with a host, drive and power source.

- 1. Once the proper setup is made, apply power to the WD1000 and all test equipment.
- Monitor TP14 (-DRUN) with a 10X oscilloscope probe while attempting to read a sector of data from the drive. The scope should be set to trigger on a high to low transition.
- 3. While observing TP14, adjust R26. The period of the DRUN single shot should be adjusted to 1.25 times the period of RCLK.

- 4. Turn off power to the WD1000.
- 5. Disconnect all test equipment.





Features 31 32 33 SYSTEN 34 35 36 37 ADDRESS BUS 38 39 40 CONTROL ਜ BUS USACK DATA BUS 29

Pin Descriptions **A₀-A₁₅.** Address Bus (output, active High, 3-state). A₀-A₁₅ form a 16-bit address bus. The Address Bus provides the address for memory data bus exchanges (up to 64K bytes) and for I/O device exchanges.

BUSACE. Bus Acknowledge (output, active Low). Bus Acknowledge indicates to the requesting device that the CPU address bus, data bus, and control signals MREQ, IORQ, RD, and WR have entered their highimpedance states. The external circuitry can now control these lines.

BUSREQ. Bus Request (input, active Low). Bus Request has a higher priority than NMI and is always recognized at the end of the current machine cycle. BUSREQ forces the CPU address bus, data bus, and control signals MREQ, IORQ, RD, and WR to go to a highimpedance state so that other devices can control these lines. BUSREQ is normally wire-ORed and requires an external pullup for these applications. Extended BUSREQ periods due to extensive DMA operations can prevent the CPU from properly refreshing dynamic RAMs.

 $D_0-D_7.$ Data Bus (input/output, active High, 3-state). D_0-D_7 constitute an 8-bit bidirectional data bus, used for data exchanges with memory and I/O.

HALT. Halt State (output, active Low). HALT indicates that the CPU has executed a Halt instruction and is awaiting either a nonmaskable or a maskable interrupt (with the

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mask enabled) before operation can resume. While halted, the CPU executes NOPs to maintain memory refresh.

INT. Interrupt Request (input, active Low). Interrupt Request is generated by I/O devices. The CPU honors a request at the end of the current instruction if the internal softwarecontrolled interrupt enable flip-flop (IFF) is enabled. INT is normally wire-ORed and requires an external pullup for these applications.

IORQ. Input/Output Request (output, active Low, 3-state). IORQ indicates that the lower half of the address bus holds a valid I/O address for an I/O read or write operation. IORQ is also generated concurrently with MI during an interrupt acknowledge cycle to indicate that an interrupt response vector can be placed on the data bus.

 $\label{eq:main_state} \begin{array}{|c|c|c|c|} \hline \textbf{M}\textbf{I}. & Machine Cycle One (output, active Low). \\ \hline \textbf{M}\textbf{I}, together with $\overline{\texttt{MREQ}}$, indicates that the current machine cycle is the opcode fetch cycle of an instruction execution. $\overline{\texttt{M1}}$, together with $\overline{\texttt{IORQ}}$, indicates an interrupt acknowledge cycle.} \end{array}$

MREQ. Memory Request (output, active Low, 3-state). MREQ indicates that the address bus holds a valid address for a memory read or memory write operation.

NMI. Non-Maskable Interrupt (input, active Low). NMI has a higher priority than INT. NMI is always recognized at the end of the current instruction, independent of the status of the interrupt enable flip-flop, and automatically forces the CPU to restart at location 0066H.

RD. Memory Read (output, active Low, 3-state). RD indicates that the CPU wants to read data from memory or an I/O device. The addressed I/O device or memory should use . this signal to gate data onto the CPU data bus. **RESET.** Reset (input, active Low). **RESET** initializes the CPU as follows: it resets the interrupt enable flip-flop. clears the PC and

Registers I and R, and sets the interrupt status to Mode 0. During reset time, the address and data bus go to a high-impedance state, and all control output signals go to the inactive state. Note that RESET must be active for a minimum of three full clock cycles before the reset operation is complete.

RFSH. Refresh (output, active Low). RFSH, together with MREQ, indicates that the lower seven bits of the system's address bus can be

used as a refresh address to the system's dynamic memories.

WAIT. Wait (input, active Low). WAIT indicates to the CPU that the addressed memory or I/O devices are not ready for a data transfer. The CPU continues to enter a Wait state as long as this signal is active. Extended

The 280 microprocessor has one of the most powerful and versatile instruction sets available in any 8-bit microprocessor. It includes such unique operations as a block move for fast, efficient data transfers within memory or between memory and I/O. It also allows operations on any bit in any location in memory.

Instruction

Set

8-Bit

Load

Group

The following is a summary of the Z80 instruction set and shows the assembly language mnemonic, the operation, the flag status, and gives comments on each instruction. The Z80 CPU Technical Manual (03-0029-01) and Assembly Language Programming Manual (03-0002-01) contain significantly more details for programming use. The instructions are divided into the

following categories:

□ 8-bit loads

 Rotates and shifts
 Bit set, reset, and test operations
 Jumps
 Calls, returns, and restarts
 Input and output operations
 A variety of addressing modes are implemented to permit efficient and fast data transfer between various registers, memory locations, and input/output devices. These addressing modes include:
 Immediate
 Indexed
 Register
 Modified page zero
 Register indirect

WAIT periods can prevent the CPU from refreshing dynamic memory properly.

WR. Memory Write (output, active Low,

□ 8-bit arithmetic and logic operations

General-purpose arithmetic and CPU

memory or I/O location.

16-bit arithmetic operations

control

3-state). WR indicates that the CPU data bus

holds valid data to be stored at the addressed

| D 16-bit | | | | | | | | b .c | _ | | Relat | | | | □ Implied |
|-----------------------|--|-------|-----|------|----------|-----|-----|-------------|-----|-----------------------------|-------|------|-------------------|---------|---------------------|
| L Excha | nges, block tr | ansie | ers | , aı | na | 300 | arc | ne | 8 | | sxte | nded | | | D Bit |
| Macanonic | Symbolic Operation | | z | | Fie H | gt | P/1 | . 14 | с | Opcode 78 543 21 |) Hea | | No.of M Cycles | | Commonte |
| LD r, r' LD r, n | r - r' r - n | : | : | X | : | X | | : | : | 01 r r 00 r 11 | | 1 2 | 1 2 | 1 | r. r' Reg. 000 B |
| | | | | | | | | | | - n - | | | - | | 001 C |
| LD r, (HL) | r - (HL) | : | : | X | : | X | : | : | : | 01 7 11 | | 1 | 25 | 7 19 | 010 D 011 E |
| LD r, (IX + d) | r (IX + d) | • | • | ^ | • | ^ | • | • | • | 01 r 10 | | 3 | 5 | 19 | 100 H |
| | | | | | | | | | | - d | | | | | 101 L |
| LD r, (I¥ + d) | r (IY + d) | • | • | X | • | x | • | • | • | 11 111 10 01 r 11 | | з | 5 | 19 | 111 A |
| | | | | | | | | | | - d - | · . | | | | |
| LD (HL), # | (HL) - r | • | ٠ | X | ٠ | x | ٠ | • | . • | 01 110 r | | 1 | 2 | 7 | |
| LD (IX + d), r | (IX+d) — r | • | • | X | • | x | • | • | • | 11 011 10 01 110 r | | 3 | 5 | 19 | |
| | | | | | | | | | | - d - | | | | | |
| LD (I¥ + d), r | (IY + d) ← r | • | ٠ | x | • | x | ٠ | ٠ | ٠ | 11 111 10 | I FD | 3 | 5 | 19 | |
| | | | | | | | | | | 01 110 r d | | | | | |
| LD (HL), n | (HL) n | • | ٠ | x | • | х | ٠ | • | • | 00 110 11 | 36 | 2 | 3 | 10 | |
| | . n | | | | | | | | | - n - | | | | | |
| LD (1 X +d), n | $(\mathbf{IX} + \mathbf{d}) \leftarrow \mathbf{n}$ | • | • | x | • | * | • | • | • | 11 011 10 | | 4 | 5 | 19 | |
| | | | | | | | | | | - d - | | | | | |
| LD (I¥ + d), n | 11 1 | | | x | | x | | | | 11 111 10 | - | 4 | 5 | 19 | |
| LD (11 + 0), h | (11 + a) ← n | • | • | ^ | • | ^ | | • | • | 00 110 110 | | • | 5 | 19 | |
| | | | | | | | | | | - d - | | | | | |
| LD A, (BC) | A - (BC) | | | x | | x | | | | 00 001 010 | 0.8 | 1 | 2 | 7 | |
| LD A. (DE) | A - (DE) | • | • | x | | X | • | • | | 00 011 010 | | 1 | 2 | 7 | |
| LD A, (nn) | A (nn) | • | ٠ | x | • | x | ٠ | ٠ | • | 00 111 010 | 38 | 3 | 4 | 13 | |
| | | | | | | | | | | - n - - n - | | | | | |
| LD (BC), A | (BC) - A | • | • | X | • | x | • | ٠ | • | 00 000 010 | | 1 | 2 | 7 | |
| LD (DE), A | (DE) - A | • | ٠ | X | • | X. | • | ٠ | | 00 010 010 | | 1 | 2 | 7 13 | |
| LD (nn), A | (nn) — Å | • | • | x | • | X | •. | • | • | 00 110 010 | 34 | 3 | 4 | 13 | |
| | | | | | | | | | | - n - | | | | | |
| LD A, I | A - I | 1 | 1 | x | 0 | X | IFF | 0 | • | 11 101 101 | | 2 | 2 | 9 | |
| LD A. R | A - B | | 1 | x | 0 | x | IFF | 0 | | 01 010 111 | | 2 | 2 | 9 | |
| | | • | | | | | | | | 01 011 111 | 5F | | - | - | |
| LD I, A | 1 - A | • | • | X | • | x | • | ٠ | • | 11 101 101 01 01 01 000 111 | | 2 | 2 | 9 | |
| LD R, A | R - A | • | • | x | • | x | • | • | • | 11 101 101 | | 2 | 2 | 9 | |
| | | | | | | | | | | 01 001 111 | | | | | |

IFF the content of the interrupt enable flip-flop, (IFF) is memonic tables, see Symbolic Notation and symi copied into the P/V flag.

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| 16-Bit Load Group | Mnemonic | Symbolic Operation | 8 | z | | Fle H | | P/V | | | Opcode 76 543 210 Hex | Bytes | No.of M Cycles | States | Comments |
|----------------------|---|---|----------------------------------|-------------------------|-------------|----------|-------------|--------|-------|----------|--|-------|---|----------|---|
| | LD dd, nn | dd - nn | • | • | X | • | X | • | • | • | 00 dd0 001 - n - - n - | 3 | 3 | 10 | dd Pair 00 BC 01 DE |
| | LD IX, nn | IX — nn | • | | X | • | x | • | • | • | 11 011 101 DD 00 100 001 21 | 4 | 4 | 14 | 10 HL 11 SP |
| | LD IY, nn | IY - nn | • | • | x | • | X | • | • | • | 11 111 101 FD 00 100 001 21 | 4 | 4 | 14 | |
| | LD HL, (nn) | H - (nn+1) L - (nn) | • | • | X | • | x | • | • | • | 00 101 010 2A | 3 | 5 | 16 | |
| | LD dd, (nn) | | • | • | x | • | X | • | • | • | - n - 11 101 101 ED 01 dd1 011 - n - | 4 | 6 | 20 | |
| | LD IX, (nn) | $IX_{H} - (nn+1)$ $IX_{L} - (nn)$ | • • • | • | x | • | x | • | • | • | - n - 11 011 101 DD 00 101 010 2A - n - | 4 | 6 | 20, | |
| | LD IY, (nn) | $IY_H - (nn+1)$ $IY_L - (nn)$ | • | •, | x | • | X | • | • | • | 11 111 101 FD 00 101 010 2A | 4 | 6 | 20 | |
| | LD (nn), HL | (nn+1) - H (nn) - L | | • | x | • | x | • | • | • :' | 00 100 010 22 - n - | 3 | 5 | 16 | |
| | LD (nn), dd | (nn+1) - ddH (nn) - ddL | • | • | x | •. | x | • | • | • | 11 101 101 ED 01 dd0 011 - n - - n - | 4 | 6 | 20 | |
| | LD (nn), 1X | $(nn+1) - IX_H$ $(nn) - IX_L$ | • | • | x | • | X | • | • | • | 11 011 101 DD 00 100 010 22 | 4 | 6 | 20 | |
| | LD (nn), IY | $(nn+1) - IY_H$ $(nn) - IY_L$ | • | • | x | • | x | • | • | • | - n - 11 111 101 FD 00 100 010 22 - n - | 4 | 6 | 20 | |
| | LD SP, HL LD SP, IX | SP - HL SP - IX | : | : | X X | : | X X | : | : | : | ← n → 11 111 001 F9 11 011 101 DD 11 111 001 F9 | 1 | 1 2 | 6 10 | |
| | LD SP, IY | SP - 1Y | • | • | X | • , | x | • | • | • | 11 111 101 FD | 2 | 2 | 10 | |
| | PUSH qq | (SP - 2) - qqL (SP - 1) - qqH SP - SP - 2 | • | ÷, | x | • | x | • | • * | • | 11 111 001 F9 11 gg0 101 | 1 | 3 | 11 | 99 Pair 00 BC 01 DE 10 HL |
| | PUSH IX | $(SP-2) - IX_L$ $(SP-1) - IX_H$ | • | • | X | • | x | • | • | • | 11 011 101 DD 11 100 101 E5 | 2 | 4 | 15 | 11 AF |
| | PUSH IY | $(SP-2) = IY_L$ $(SP-1) = IY_H$ SP = SP = 2 | • | • | X | | | | • | • | 11 111 101 FD 11 100 101 E5 | 2 | 4 | 15 | |
| | POP qq | $qq_H - (SP+1)$ $qq_L - (SP)$ SP - SP + 2 | • | • | x | | | • | • | • | 11 qq0 001 | 1 | 3 | 10 | |
| | POP IX | $IX_H - (SP + 1)$ $IX_L - (SP)$ SP - SP + 2 | • | • | X | • | X | • | • | • | 11 011 101 DD 11 100 001 E1 | 2 | 4 | - 14 | |
| | POP IY | $IY_{H} - (SP+1)$ $IY_{L} - (SP)$ SP - SP + 2 | , **** , **** | •., | X. | • | X | • | • | • | 11 111 101 FD 11 100 001 E1 | 2 | 4 | 14 | |
| | NOTES: dd is a qq is a (PAIR) e.g. | ny of the register pairs BC ny of the register pairs AF H, (PAIR)L refer to high o BCL = C, AFH = A. | , DE, HI , BC, Di rder and | L, SP E, HI I low | order | eigh | t bite | of the | e reg | jister j | air respectively, | | | | |
| Exchange. | EX DE, HL | DE - HL AF - AF | : | • | X X X | | | • | : | : | 11 101 011 EB 00 001 000 08 | 1 | 1 | 4 | · · · · · · · · · · · · · · · · · · · |
| Block Fransfer, | EX AF, AF' EXX | AF = AF BC = BC' DE = DE' HL = HL' | • | ÷ | Ŷ | • | X X X | : | • | • | 11 011 001 D9 | i | i | 4 | Register bank and auxiliary register bank exchange |
| Block Search | EX (SP), HL | $H \rightarrow (SP + 1)$ | • | ٠ | X | ٠ | x | ٠ | • | • | 11 100 011 E3 | 1 | 5 | 19 | |
| roups | EX (SP), IX | $L = (SP)$ $IX_{H} = (SP+1)$ $IX_{L} = (SP)$ $IY_{H} = (SP+1)$ $IY_{H} = (SP+1)$ | • | • | x | • • | x | • | • | • | 11 011 101 DD 11 100 011 E3 | 2 | 6 | 23 | |
| | EX (SP), IY | 11L = (SP) | • | • | X | • | X | • | • | • | 11 111 101 FD 11 100 011 E3 | 2 | 6 | 23 | |
| | LDI | (DE) - (HL) DE - DE + 1 -HL - HL + 1 BC - BC - 1 | • | | X | 0 | X | 1 | 0 | • | 11 101 101 ED 10 100 000 A0 | 2 | 4 ₁₀ -1 1 - 1 - 1 1 - 1 | 16 | Load (HL) into (DE), increment the pointers and decrement the byte counter (BC) |
| | LDIR | (DE) - (HL) DE - DE + 1 HL - HL + 1 BC - BC - 1 | • | • | x | 0 | x | 0 | 0 | • | 11 101 101 ED 10 110 000 B0 | 22 | 5 4 | 21 16 | If BC $\neq 0$ If BC =0 |
| | | Repeat until BC = 0 | | | | | | | | | | | | | |

| xchange. lock | Mnemonic | Symbolic Operation | s | z | | Flo H | ige j | P/V | N | с | 0 78 | pcoc 543 | ie 210 | Hex | No.of Bytes | No.of M Cycles | No.of T States | Comments |
|---|--|--|--|---|---|--|---|---|---|---|---|------------------------------------|--|------------|---------------------------------|----------------------------|---------------------------------|---|
| ransfer, | | | | | | | | 0 | | | | | | | | | | |
| lock Search | LDD | (DE) - (HL) | • | • | x | 0 | х | 1 | 0 | • | | | 101 | | 2 | 4 | 16 | |
| roups | | DE - DE - 1 HL - HL - 1 | | | | | | | | | 10 | 101 | 000 | A8 . | | | | |
| Continued) | | BC - BC - 1 | | | | | | | | | | | | | | | | |
| | LDDR | (DE) - (HL) | | | X | 0 | х | 0 | 0 | | п | 101 | 101 | ED | 2 | 5 | 21 | If BC ≠ 0 |
| | | DE - DE - 1 | | | | | | | | | | | 000 | | 2 | 4 | 16 | If BC = 0 |
| | | HL - HL - 1 BC - BC - 1 | | | | | | | | | | | | | | | | |
| | | Repeat until | | | | | | | | | | | | | | | | |
| | | BC = 0 | | Q | | | | 0 | | | | | | | | | | |
| | CPI | A - (HL) | 1 | 1 | X | 1 | X | 1 | 1 | • | | | 101 | | 2 | 4 | 16 | |
| | | HL HL + I BC BC - I | | | | | | | | | 10 | 100 | 001 | AI | | | | |
| | | | | 0 | | | | 0 | | | | · | | | | | | |
| | CPIR | A - (HL) | ' | 1 | х | 1 | х | 1 | 1 | • | 11 | 101 | 101 | ED | 2 | 5 | 21 | If BC \neq 0 and A \neq (HL) |
| | | HL - HL + 1 | | | | | | | | | 10 | 110 | 001 | Bl | 2 | 4 | 16 | If BC = 0 or |
| | | BC ← BC - 1 Repeat until | | | | | | | | | | | | | | | | $\mathbf{A} = (\mathbf{HL})$ |
| | | A = (HL) or | | | | | | | | | | | | | | | | |
| | | BC = 0 | | C | | | | 0 | | | | | | | | | | |
| | CPD | A - (HL) | 1 | ĩ | | 1 | x | | 1 | • | п | 101 | 101 | ED | 2 | .4 | 16 | |
| | | HL - HL - 1 BC - BC - 1 | | | | | | | | | 10 | 101 | 001 | A 9 | | | | |
| | | | | 0 | | | | 0 | | | | | | | | | | |
| | CPDR | A – (HL) | 1 | 1 | x | 1 | х | t | 1 | • | 11 | 101 | 101 | ED | 2 | 5 | 21 | If BC ≠ 0 and A ≠ (HL) |
| | | HL - HL - I | | | | | | | | | 10 | ш | 001 | B9 | 2 | 4 | 16 | If BC = 0 or |
| | | BC - BC - 1 | | | | | | | | | | | | | | | | $\mathbf{A} = (\mathbf{HL})$ |
| | | Repeat until A = (HL) or | | | | | | | | | | | | | | | | |
| | | | | | | | | | | | | | | | | | | |
| | | BC = 0 leg is 0 if the result of B is 1 if A = (HL), other | | | herwi | e P/ | / = 1 | • . | | | | | | | | | | |
| | (2) Z filog ADD A, r | leg is 0 if the result of B is 1 if A = (HL), other A ← A + r | | | x | | x | v | 0 | | | 000 | | | 1 | 1 | 4 | rReg. |
| Arithmetic | 2 fleg | leg is 0 if the result of B is 1 if A = (HL), other | wise Z = | 0. | | | x | v | 0 0 | | | 000 | 110 | | 12 | 1 2 | 4 7 | 000 B |
| Arithmetic and Logical | (2) 2 tieg ADD A, r ADD A, n | leg is 0 if the result of B is 1 if $A = (HL)$, other $A \leftarrow A + r$ $A \leftarrow A + n$ | wise Z = 1 1 | : 0. 1 1 | x x | 1 | x x | v v | 0 | 1 | 11 | 000 n | 110 | | 2 | 2 | 7 | 000 B 001 C 010 D |
| Arithmetic and Logical | (2) Z tieg ADD A, r ADD A, n ADD A, (HL) | leg is 0 if the result of B is 1 if A = (HL), other A - A + r A - A + n A - A + n | wise Z = 1 1 | : 0. 1 1 | x x x | 1 1 1 | x x x | v v v | 0 | 1 1 | 11 | | 110 - 110 | | 2 | 2 | 7 | 000 B 001 C 010 D 071 E |
| Arithmetic and Logical | (2) Z tieg ADD A, r ADD A, n ADD A, (HL) | leg is 0 if the result of B is 1 if $A = (HL)$, other $A \leftarrow A + r$ $A \leftarrow A + n$ | wise Z = 1 1 | : 0. 1 1 | x x | 1 | x x | v v v | 0 | 1 1 | 11 | n 0000 011 | 110 - 110 101 | DD | 2 | 2 | 7 | 000 B 001 C 010 D 071 E 100 H |
| Arithmetic and Logical | 2 flag ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) | leg us 0 if the result of B is 1 if $A = (HL)$, other A - A + r A - A + n A - A + (HL) A - A + (IX+d) | vise Z = I I I I | 1 1 1 1 | x x x x | 1 1 1 1 | x x x x | v v v v | 0 | 1 | 10 10 11 10 | n 000 011 000 d | 110 110 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 071 E |
| rithmetic Ind Logical | 2 flag ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) | leg is 0 if the result of B is 1 if A = (HL), other A - A + r A - A + n A - A + n | wise Z = 1 1 | : 0. 1 1 | x x x | 1 1 1 1 | x x x | v v v v | 0 | 1 | 11 10 11 10 | n 000 011 000 d 111 | 110 110 101 110 | DD | 2 1 3 | 2 | 7 | 000 B 001 C 010 D 071 E 100 H 101 L |
| Arithmetic and Logical | 2 flag ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) | leg us 0 if the result of B is 1 if $A = (HL)$, other A - A + r A - A + n A - A + (HL) A - A + (IX+d) | vise Z = I I I I | 1 1 1 1 | x x x x | 1 1 1 1 | x x x x | v v v v | 0 | 1 | 10 10 11 10 11 | | 110 110 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 071 E 100 H 101 L |
| Arithmetic and Logical | 2 flag ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) | leg us 0 if the result of B is 1 if $A = (HL)$, other A - A + r A - A + n A - A + (HL) A - A + (IX+d) | vise Z = I I I I | 1 1 1 1 | x x x x | 1 1 1 1 | x x x x | v v v v | 0 | 1 | 10 10 11 10 11 | n 000 011 000 d 111 | 110 110 101 110 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 C 011 E 100 H 101 L 111 A s is any of r. n. |
| Arithmetic and Logical | (1) 2 (ling ADD A, r ADD A, (HL) ADD A, (IX + d) ADD A, (IY + d) | lag us 0 it the result of B u = it A = (HL), other A = A + r A = A + r A = A + n A = A + (HL) A = A + (IX + d) A = A + (IY + d) | wise Z = | 1 1 1 1 | x x x x x | 1 | x x x x x | v v v v | 0 0 0 | 1 | 10 11 10 11 10 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A 111 A 111 A |
| irithmetic and Logical | () Z Ileg ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) ADD A, (IY + d) ADC A, s | lag is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + r A = A + r A = A + (HL) A = A + (HL) A = A + (IY + d) A = A + s + CY | wise Z = | 1 1 1 1 | x x x x x x | 1 1 1 1 1 | x x x x x x | V V V V V | 0 0 0 | 1 | 10 11 10 11 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A 114 A 115 A 116 A 117 A 118 A 119 B 111 A |
| Arithmetic and Logical | ADD A, r ADD A, n ADD A, (HL) ADD A, (HL) ADD A, (IX + d) ADD A, (IY + d) ADC A, s SUB s | leg is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + n A = A + n A = A + (HL) A = A + (IX + d) A = A + (IY + d) A = A + s + CY A = A - s | wise Z = | 1 1 1 1 | x x x x x x x x | | x x x x x x x x | V V V V V V | 0 0 0 1 | 1 1 1 1 | 10 11 10 11 10 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| Arithmetic and Logical | ADD A, r ADD A, n ADD A, n ADD A, (IL) ADD A, (IX + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s | leq is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + r A = A + n A = A + (HL) A = A + (HL) A = A + (IX + d) A = A + (IY + d) A = A + cY A = A - c | wise Z = | 1 1 1 1 | x x x x x x x x x | | x x x x x x x x x x | v v v v v v v v v | 0 0 0 1 1 | 1 1 1 1 1 | 10 11 10 11 10 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| rithmetic Ind Logical | () 2 (lisq) ADD A, r ADD A, n ADD A, (HL) ADD A, (IX + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s AND s | leg is 0 if the result of B is i if $A = (HL)$, other A = A + r A = A + r A = A + (HL) A = A + (HL + d) A = A + (IX + d) A = A + (IY + d) A = A + s + CY A = A - s - CY A = A - s | wise Z = | 1 1 1 1 | x x x x x x x x x x x x | 1 1 1 1 1 1 | x x x x x x x x x x x x | V V V V V V V V P | 0 0 0 1 1 0 | 1 1 1 1 1 1 1 0 | 10 10 11 10 11 10 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| rithmetic Ind Logical | (d) 2 flieg ADD A. r ADD A. dD A. (HL) ADD A. (IX + d) ADC A. s SUB s SBC A. s AND s OR s | leg is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + n A = A + n A = A + (HL) A = A + (HL + d) A = A + (IY + d) A = A + s + cY A = A - s = cY A = A - s A = A - s | wise Z = | 1 1 1 1 | x x x x x x x x x x x x x x x | 1 1 1 1 1 1 1 0 | x x x x x x x x x x x x x x x | V V V V V V P P P | 0 0 0 1 1 0 0 | 1 1 1 1 1 1 1 1 | 10 10 11 10 11 10 10 10 | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| Arithmetic and Logical | ADD A, r ADD A, n ADD A, (IL) ADD A, (IX+d) ADD A, (IY+d) ADD A, (IY+d) ADC A, s SUB | | wise Z = | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x | 1 1 1 1 1 1 0 0 | x x x x x x x x x x x x x x x | V V V V V V V P P P V | 0 0 0 1 1 0 0 | 1 1 1 1 0 0 0 | | | 110 110 101 101 101 110 | | 2 1 3 | 2 2 5 | 7 7 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| rithmetic Ind Logical | (b) 2 lisq ADD A, r ADD A, n ADD A, (HL) ADD A, (IX+d) ADD A, (IY+d) ADC A, s SUB s SBC A, s AND s OR s COP s INC (FL) | $\begin{array}{l} \text{Aeg a 0 (if the result of B}\\ \text{as } i \text{ of } A = (HL), \text{ other}\\ \text{A} = A + r\\ \text{A} = A + r\\ \text{A} = A + r\\ \text{A} = A + (HL)\\ \text{A} =$ | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x x x x | 1 1 1 1 1 1 0 0 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 1 0 0 | 1 1 1 1 1 0 0 0 0 1 1 1 | 10 10 11 10 11 10 10 11 10 10 10 10 10 1 | | 110 | FD | 2 1 3 3 | 2 5 5 1 3 | 7 19 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| Arithmetic and Logical | ADD A, r ADD A, n ADD A, (IL) ADD A, (IL+d) ADD A, (IX+d) ADD A, (IX+d) ADD A, (IX+d) ADC A, s SUB s SBC A, s SNB s SBC A, s AND s OR s CP s IINC r | leg is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + r A = A + r A = A + r A = A + (HL) A = A + (HL) A = A + (HL + d) A = A + (IX + d) A = A + s + CY A = A - s = A = A - s = A = A - s = A = A + s = A | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x x x x | 1 1 1 1 1 0 0 1 1 | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 1 0 0 1 0 | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | 10 10 11 10 10 11 10 10 10 10 10 10 10 1 | | 110 | | 2 1 3 3 | 2 2 5 5 | 7 19 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| Arithmetic and Logical | A DD A, r A DD A, n A DD A, (HL) ADD A , (IX + d) ADD A , (IX + d) ADD A , (IX + d) ADC A , s SUB s SBC A , s SUB s SBC A , s SDC A , s SDC A , s SDC S SDC A , s SDC S SDC C , s SDC S SDC C , s SDC S SDC S S SDC S SDC S SDC S SDC S SDC S SDC S | $ \begin{array}{l} \mbox{leq} \ a \ 0 \ (i \ ha \ result \ o \ B \) \\ \mbox{leq} \ a \ 0 \ (i \ ha \ result \ o \ B \) \\ \mbox{leq} \ A = A + (RL), \ other \ \\ \mbox{des} \ A = A + (RL) \\ \mbox{des} \ A = A + (RL) \\ \mbox{des} \ A = A + (RL) \\ \mbox{des} \ A = A + (IX + d) \\ \mbox{des} \ A = A $ | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x x x x | 1 1 1 1 1 1 0 0 1 1 1 1 1 1 1 1 1 1 1 1 | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 1 0 0 | 1 1 1 1 0 0 0 1 1 | 11 10 11 10 11 10 10 10 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| Arithmetic and Logical | (b) 2 lisq ADD A, r ADD A, n ADD A, (HL) ADD A, (IX+d) ADD A, (IY+d) ADC A, s SUB s SBC A, s AND s OR s COP s INC (FL) | leg is 0 if the result of B is 1 if $A = (HL)$, other A = A + r A = A + (HL) A = A + (HY + d) A = A + (Y + d) A = A + s + CY A = A - s A = A - s A = A + s A = | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 1 0 0 1 0 0 | 1 1 1 1 0 0 0 1 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 | 7 19 19 19 19 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A |
| 3-Bit Arithmetic Ind Logical Group | Theorem ADD A, r ADD A, n ADD A, (HL) ADD A, (IY + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s AND s OR s XOR s CP s INC (IX + d) INC (IY + d) | $ \begin{array}{l} \text{ ag a 0 (I the result of B \\ \text{ is } i \text{ d } A = (HL), \text{ other} \\ \hline A = A + r \\ A = A + r \\ A = A + (HL) \\ A = A + (HL) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + S + CY \\ A = A - S \\ A = A + S \\ A = A - S \\ A = A + S \\ A =$ | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 0 1 0 0 0 0 | 1 1 1 1 1 0 0 0 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 19 19 19 | 000 B 001 C 010 D 011 E 100 H 101 L 111 A * Is any of r. n, (HL) (IX - d), (IY - d) as aboven for ADD matrixuton. The indicated bits replace the EOM in the ADD set above. |
| Arithmetic and Logical | A DD A, r A DD A, n A DD A, (HL) ADD A , (IX + d) ADD A , (IX + d) ADD A , (IX + d) ADC A , s SUB s SBC A , s SUB s SBC A , s SDC A , s SDC A , s SDC S SDC A , s SDC S SDC C , s SDC S SDC C , s SDC S SDC S S SDC S SDC S SDC S SDC S SDC S SDC S | $ \begin{array}{l} \text{ ag a 0 (I the result of B \\ \text{ is } i \text{ d } A = (HL), \text{ other} \\ \hline A = A + r \\ A = A + r \\ A = A + (HL) \\ A = A + (HL) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + $ | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 0 1 0 0 0 0 | 1 1 1 1 1 0 0 0 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 19 19 19 | 000 B 001 C 001 C 001 E 100 H 101 L 111 A * is any of r, n, (HL), (X + d), (HT + d) as aboven for ADD matruction. The indicated bias replace the [503] in the ADD matruction. The indicate bias replace the [503] in the ADD ast above. m is any of r, (HL), |
| Arithmetic and Logical | Theorem ADD A, r ADD A, n ADD A, (HL) ADD A, (IY + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s AND s OR s XOR s CP s INC (IX + d) INC (IY + d) | $ \begin{array}{l} \text{ ag a 0 (I the result of B \\ \text{ is } i \text{ d } A = (HL), \text{ other} \\ \hline A = A + r \\ A = A + r \\ A = A + (HL) \\ A = A + (HL) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + S + CY \\ A = A - S \\ A = A + S \\ A = A - S \\ A = A + S \\ A =$ | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 0 1 0 0 0 0 | 1 1 1 1 1 0 0 0 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 19 19 19 | 000 B 001 C 011 E 100 H 101 L 111 A * is any of r, n, (HL), (X*-d), (Y+d) as above. The indicated bits replace the EOD and above. m is any of r, (HL), (X*-d), (Y*-d), (Y*-d) |
| rithmetic Ind Logical | Theorem ADD A, r ADD A, n ADD A, (HL) ADD A, (IY + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s AND s OR s XOR s CP s INC (IX + d) INC (IY + d) | $ \begin{array}{l} \text{ ag a 0 (I the result of B \\ \text{ is } i \text{ d } A = (HL), \text{ other} \\ \hline A = A + r \\ A = A + r \\ A = A + (HL) \\ A = A + (HL) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + S + CY \\ A = A - S \\ A = A + S \\ A = A - S \\ A = A + S \\ A =$ | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 0 1 0 0 0 0 | 1 1 1 1 1 0 0 0 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 19 19 19 | 000 B 001 C 001 D 001 E 100 H 101 E 100 H 101 L 111 A * is any of r, 0; (H), 0; X + d), (Y + d) as aboven for ADD matrixetons. The indicated bits replace the EOD at above. m is any of r, 0; (H), 0; (X + d), 0; Y + d), as boven for INC. DEC same format |
| irithmetic and Logical | O Z fleq ADD A, r ADD A, n ADD A, (HL) ADD A, (IY + d) ADD A, (IY + d) ADC A, s SUB s SBC A, s AND s OR s XOR s CP s INC (IX + d) INC (IY + d) | $ \begin{array}{l} \text{ ag a 0 (I the result of B \\ \text{ is } i \text{ d } A = (HL), \text{ other} \\ \hline A = A + r \\ A = A + r \\ A = A + (HL) \\ A = A + (HL) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + (IY + d) \\ \hline A = A + (IY + d) \\ A = A + S + CY \\ A = A - S \\ A = A + S \\ A = A - S \\ A = A + S \\ A =$ | wise Z = 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 | | x x x x x x x x x x x x x x x x x x x | | x x x x x x x x x x x x x x x x x x x | V V V V V V V V V V V V V V V V V V V | 0 0 0 1 1 0 0 0 1 0 0 0 0 | 1 1 1 1 1 0 0 0 1 | 11 10 11 10 11 10 11 10 11 10 10 | | | FD | 2 1 3 3 1 1 3 | 2 5 5 1 3 6 | 7 19 19 19 19 19 | 000 B 001 C 001 C 001 D 01 D 11 B 101 L 111 A 111 <t< td=""></t<> |

| Seneral- Purpose | Mnomonic | Symbolic Operation | 8 | z | | Fla H | iĝa | P/V | N | с | Opcode 76 543 210 H | lex 1 | No.of Bytes | No.of N Cycles | No.of T States | Comments |
|----------------------------------|------------------|---|--------|------|-----------------------|------------|--------|-----|---|---|---|----------|----------------|-------------------|-------------------|--|
| Arithmetic and CPU Control | DAA | Converts acc. content into packed BCD following add or subtract with packed | 1 | 1 | x | 1 | X | P | • | 1 | 00 100 111 2 | 27 | 1 | 1 | 4 | Decimal adjust accumulator. |
| FO Control Froups | CPL | BCD operands. A - Ā | • | • | x | 1 | x | • , | 1 | • | 00 101 111 2 | 2F | 1 | ı | 4 | Complement accumulator (one's complement). |
| | NEG | A - 0 - A | 1 | r | x | ŧ, | X | v | 1 | 1 | 11 101 101 H 01 000 100 4 | ED | 2 | 2 | 8 | complement). Negate acc. (two's complement). |
| | CCF | CY - CY | • | • | X | x | X | • | 0 | 1 | 00 111 111 3 | BF | 1 | 1 | 4 | Complement carry flag. |
| | SCF | CY - 1 No operation | : | : | X X X X X | • | X X | : | • | 1 | 00 110 111 3 | 00 | 1 | 1 | 4 | Set carry flag. |
| | HALT DI + | CPU halted IFF - 0 | : | : | X | : | X X | : | : | : | 01 110 110 1 | 73 | 1 | 1 | 4 | |
| | EI + IM 0 | IFF - 1 Set interrupt | : | : | X X | : | X | : | : | : | 11 111 011 1 11 101 101 1 | Ð | 1 2 | 1 2 | 4 8 | |
| | IM I | mode 0 Set interrupt mode 1 | • | • | X | • | X | • | • | • | 01 000 110 4 11 101 101 1 01 010 110 1 | ED | 2 | 2 | 8 | |
| | IM 2 | Set interrupt mode 2 | • | • | x | • | X | • | • | • | 11 101 101 I 01 011 110 1 | ED | 2 | 2 | 8 | |
| | CY in | dicates the interrupt enable fli dicates the carry flip-flop. licates interrupts are not samp | | | end o | of El c | or Di | | | | | | | | | |
| 6-Bit | ADD HL, 58 | HL - HL+ss | • | • | х | x | x | • | 0 | 1 | 00 ssl 001 | | 1 | 3 | 11 | ss Reg. 00 BC |
| Arithmetic Group | ADC HL, ss | HL - HL+ss+CY | ı | 1 | x | x | x | v | 0 | • | 11 101 101 1 01 ssl 010 | ED | 2 | 4 | 15 | 01 DE 10 HL 11 SP |
| • | SBC HL, se | HL - HL-ss-CY | ı | 1 | x | x | x | v | 1 | ı | 11 101 101 1 01 ss0 010 | ED | 2 | 4 | 15 | 11 SP |
| | ADD IX. pp | IX — IX + pp | • | • | x | x | x | • | 0 | | 11 011 101 1 01 pp1 001 | DD | 2 | 4 | 15 | pp Reg. 00 BC 01 DE 10 IX 11 SP |
| | ADD IY, rr | IY – IY + rr | • | • | x | . x | x | • | 0 | ı | 11 111 101 1 00 rr1 001 | D | 2 | 4 | 15 | rr Reg. 00 BC 01 DE 10 IY |
| | INC ss INC IX | ss - ss + 1 IX - IX + 1 | : | • | X X | • | X X | : | : | : | 00 ss0 011 11 011 101 I | ac | 12 | 1 2 | 6 10 | 11 SP |
| | INC IN | IY - IY + 1 | | | x | • | x | • | • | • | 00 100 011 1 | 23 7D | 2 | 2 | 10 | |
| | DEC se | ss - ss - 1 IX - IX - 1 | • | • | X X | • | X | • | • | • | 00 100 011 1 00 ssl 011 11 011 101 I | | 1 | 1 | 6 | |
| | DEC IX DEC IY | IX - IX - I $IY - IY - I$ | : | : | x x | : | X X | | : | : | 00 101 011 101 1 | 2B | 2 2 | 2 | 10 10 | |
| | | ny of the register pairs BC, D any of the register pairs BC, D ny of the register pairs BC. D | E, HL | . SP | | | | | | | 00 101 011 2 | 2B | | | | |
| | rr is al | ny of the register pairs BC. D | E. IY. | SP. | | | | | | | | | | | | |
| Rotate and Shift Group | RLCA | CY- 7-0+ | • | • | X | 0 | x | • | 0 | ; | 00 000 111 | 07 | 1 | : 1 | 4 | Rotate left circular accumulator. |
| | RLA | -CY70 | • | • | X | 0 | X | • | 0 | 1 | 00 010 111 | 17 | 1 | 1 | 4 | Rotate left accumulator. |
| | RRCA | | • | • | x | 0 | х | • | 0 | t | 00 001 111 | 0F | 1 | 1 | 4 | Rotate right circular accumulator. |
| | RRA | | • | • | x | 0 | Х | • | 0 | 1 | 00 011 111 | 1F | 1 | 1 | 4 | Rotate right accumulator. |
| | RLC r | | 1 | I | x | 0 | X | Ρ | 0 | 1 | 11 001 011 00 000 r | СВ | 2 | 2 | 8 | Rotate left circular register r. |
| | RLC (HL) | | ١, | 1 | x | 0 | х | Ρ | 0 | ı | 11 001 011 00 000 110 | СВ | 2 | 4 | 15 | r Reg. 000 B 001 C |
| | RLC (IX + d) | CY + 7 - 0 + r.(HL),(IX + d),(IY + d) | 1 | 1 | x | 0 | x | Ρ | 0 | 1 | 11 011 101 11 001 011 - d - 00 000 110 | DD CB | 4 | 6 | 23 | 001 C 010 D 011 E 100 H 101 L 111 A |
| | RLC (IY+d) | | ł | 1 | x | 0 | х | P | 0 | , | 11 111 101 11 001 011 ← d → | FD CB | 4 | 6 | 23 | |
| | | | | | | | | | | | 00 000 110 | | | | | Instruction format and states are as |
| | | CY - 7 - 0 - n = r.(HL).(IX + d).(IY + d) | ı. | 1 | Х | 0 | Х | Ρ | 0 | ł | 010 | | | | | shown for RLC's. To form new |

| lotate and shift Group | Mnemonic | Symbolic Operation | 8 | z | | Flo | ige | P/V | N | с | 76 | 543 2 | 10 | Hex | No.of Bytee | No.of M Cycles | No.of T States | Comments | |
|---------------------------|-----------------|--|-------|-------|------|------|-----|-----|---|---|-----------|-----------------------|-----------|-----------|----------------|-------------------|-------------------|---|--|
| Continued) | | m=r.(HL).(IX + d).(IY + d) | ı | : | x | 0 | x | ₽ | 0 | 1 | I | 011 | | | | | | | |
| | SLA m | | t | ł | x | 0 | x | P | 0 | ı | | 100 | | | | | | | |
| | SRA m | 7 | 1 | ı | x | 0 | x | P | 0 | 1 | | 101 | | | | | | | |
| | SRL m | m = r,(HL),(IX + d),(IY + d) | 1 | ı | x | 0 | x | P | 0 | 1 | | 111 | | | | | | | |
| | RLD [| 7-43-0 7-43-0 A (HL) | 1 | ı | x | 0 | x | P | 0 | • | | 101 1 101 1 | | ED 6F | 2 | 5 | 18 | Rotate digit left and right between the accumulator | |
| | RRD (| 7-4]3-0][7-4]3-0 A (HL) | ľ | 1 | x | 0 | x | P | 0 | • | 11 01 | 101 1 100 1 | 01 11 | ED 67 | 2 | 5 | 18 | and location (HL). The content of the upper half of the accumulator is unaffected. | |
| lit Set, Reset | BIT b, r | 2 - ī _b | x | , | x | 1 | x | x | 0 | • | 11 | 001 0 | 11 | СВ | 2 | 2 | 8 | r <u>Beg.</u> 000 B | |
| ind Test | BIT b, (HL) | $Z = (\overline{HL})_{\mathbf{b}}$ | x | ı | x | 1 | x | x | 0 | • | 01 | 001 0 | ii - | СВ | 2 | 3 | 12 | 000 B 001 C 010 D 011 E | |
| Group | ВІТ Ь, (IX + d) | $\mathbf{b} \ \mathbf{Z} \leftarrow (\mathbf{\overline{IX} + d})_{\mathbf{b}}$ | X | ı | x | 1 | x | x | 0 | ٠ | 11 | 011 1 001 0 d | - 11 - | DD. CB | 4 | 5 | 20 | 100 H 101 L | |
| | | | | | | | | | | | | Б | | | | | | 111 A b Bit Tested | |
| | BIT b. (IY + d) | $I_{\rm b}$ Z = $(\overline{\rm IY} + {\rm d})_{\rm b}$ | x | 1. | x | 1 | x | x | 0 | • | 11 | 111 1 001 0 | n - | FD CB | 4 | 5 | 20 | 000 0 | |
| | | | | | | | | | | | 01 | d b l | 10 | | | | | 010 2 011 3 100 4 101 5 110 6 | |
| | SET b, r | г ь +- 1 | | | x | | x | | • | • | | 001 0 | | СВ | 2 | 2 | 8 | 111 7 | |
| | SET b, (HL) | (HL) _b - 1 | | | x | | x | • | | | [] | ь 001 с | r 011 | СВ | 2 | 4 | 15 | | |
| | |) (IX+d) _b - 1 | | | x | | x | | | | | ь I 011 I | 10 | | 4 | 6 | 23 | | |
| | 5E1 D, (IA+0 |) (IX+U/b = 1 | • | • | î | • | î | - | | | 11 | 001 с d | - | CB | • | Ū | | | |
| | SET b, (IY + d |) $(IY + d)_b = 1$ | • | • | x | • | x | • | • | • | 11 | лі 001 с а ь | 01 011 | FD CB | 4 | 6 | 23 | | |
| | RES b, m | $m_b \leftarrow 0$ m = r. (HL), (IX + d), (IY + d) | • | • | x | • | x | • | • | • | 0 | В | | | | | | To form new opcode replace [1] of SET b, s with [0]. Flags and time states for SET instruction. | |
| | NOTES: The | notation m _b indicates bit b (0 | to 7) | or lo | atio | n m. | | | | | | | | | | | | | |
| Jump | JP nn | PC - nn | • | • | x | • | x | • | • | • | 11 | 000 - n - n | 011 | СЗ | 3 | 3 | 10 | | |
| Group | JP cc, nn | If condition cc is true PC — nn, otherwise continue | • | • | x | • | x | • | • | • | | cc | 010 | | 3 | 3 | 10 | cc Condition 000 NZ non-zero 001 Z zero 010 NC non-carry 011 C carry 100 PO parity odd 101 101 PE parity even | |
| | JR e | PC - PC+e | | | x | | x | | | | 0 | 011 | 000 | 18 | 2 | 3 | 12 | 110 P sign positive 111 M sign negative | |
| | JR C, e | If C = 0, | | | x | | x | | | | ō | • • - 2 0 111 | 000 | | 2 | 2 | 7 | If condition not met. | |
| | | continue If C = 1. | | | | | | | | | - | •-2 | - | | 2 | 3 | 12 | If condition is met. | |
| | IR NC, • | $PC \leftarrow PC + e$ If $C = 1$, continue | • | • | x | • | x | • | • | • | | 0 110 - e-2 | | 30 | 2 2 | 2 3 | 7 12 | If condition not met. | |
| | JP Z, e | If C = 0, PC PC + e If Z ≠ 0 | | | x | | x | | | | ~ | 0 101 | | 28 | 2 | 2 | 7 | If condition not met. | |
| | 31° Z, e | continue | • | | * | • | ^ | | • | • | | e - 2 | ! - | 20 | 2 | 2 | 12 | If condition is met. | |
| | 10 117 | If $Z = 1$, PC - PC + e | | | x | | x | | | | ~ | 001 00 | m | 20 | 2 | 2 | 7 | If condition is met. | |
| | JR NZ, e | If Z = 1, continue If Z = 0 | • | • | ^ | | ^ | | | • | | - e-: | | 10 | 2 | 2 | 12 | If condition is met. | |
| | IP (HL) | II Z = 0, PC - PC + • PC - HL | | | y | | x | | | | 1 | 1 101 | 001 | E9 | 1 | 3 1 | 4 | a condition is inst. | |
| | JP (IX) | PC - IX | | | x | | x | | | | | 011 | | | 2 | 2 | 8 | | |
| | 1P (1X) | PC - 1X | • | • | X | • | x | • | • | • | 1 | 1 101 | 001 | E9 | 2 | 2 | 6 | | |

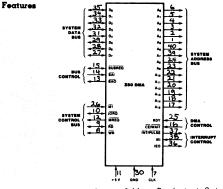
| (Continued) | Masmonie | Symbolic Operation | 8 | z | | Play H | | P/¥ | N | c | Opcode 78 543 210 | Nex | No.of Bytes | Ho.of M Cycles | No.of T States | Comments | |
|--------------------------|--|---|-----------------------------|---------------------------|---------------|-----------------------|--------------|-------------|-----|--------|--|----------|----------------|--------------------------------|-------------------|---|--|
| | JP (IY) | PC - IY | • | • | x | • | x | • | • | • | 11 111 101 11 101 001 | FD | 2 | 2 | 8 | | |
| | DINZ, • | B - B - 1 If B = 0, continue | • | • | x | • | x | • | • | • | 00 010 000 | 10 | 2 | 2 | 8 | If $\mathbf{B} = 0$. | |
| | | If B ≠ 0, PC - PC++ | | | | | | | | | | | 2 | 3 | 13 | lf B ≠ 0. | |
| | NOTES: e repre e le a s e-2 in by 2 | cents the extension in the re- ligned two's complement nur the opcode provides an effect prior to the addition of e. | lative nber in ective | eddre n the i addre | esing stor | mode < - pc + e | 126, 44 F | 129 C 10 | >. | emente | d | | | | | | |
| Call and leturn Group | CALL nn | $(SP-1) = PC_H$ $(SP-2) = PC_L$ PC = nn | • | • | x | • | x | • | • | • | 11 001 101 | CD | 3 | 5 | 17 | | |
| | CALL cc, nn | If condition cc is false continue, | • | • | x | • | x | • | • | • | 11 cc 100 - n - - n - | | 3 3 | 3 5 | 10 17 | If cc is false. | |
| | | otherwise same as CALL nn | | | | | | | | | | | | | | | |
| | RET | $PC_L - (SP)$ $PC_H - (SP + 1)$ | • | • | x | • | x | • | • | • | 11 001 001 | C9 | 1 | 3 | 10 | | |
| | RET oc | If condition cc is false continue. | • | • | x | • | x | • | • | •. | 11 cc 000 | | 1 1 | 1 3 | 5 11 | If cc is false. If cc is true. | |
| | | otherwise same as RET | | | | | | | | | | | • | Ū | | cc Condition | |
| | RETI | Return from | • | • | x | • | x | • | • | • | 11 101 101 01 001 101 | ED | 2 | 4 | 14 | 001 Z zero 010 NC non-carry 011 C carry 100 PO parity odd | |
| | RETN ¹ | interrupt Return from non-maskable interrupt | • | • | X | • | x | • | • | • | 01 001 101 11 101 101 01 000 101 | ED | 2 | 4 | 14 | 100 PO parity odd 101 PE parity even 110 P sign positive 111 M sign negative | |
| | RST p | $\begin{array}{l} (SP-1) = PC_H\\ (SP-2) = PC_L\\ PC_H = 0\\ PC_L = p \end{array}$ | • | • | x | • | x | • | • " | • | n i ni | | 1 | 3 | -11 | t p 000 00H 001 08H 010 10H | |
| | | | | | | | | | | | | | | | | 011 18H 100 20H 101 28H 110 30H 111 38H | |
| | NOTE: 'RETN I | oeds IFF ₂ - IFF ₁ | | | | | | | | | | | | | | | |
| input and | IN A. (n) | A - (n) | • | • | x | • | x | • | • | • | 11 011 011 | DB | 2 | 3 | п | n to Ag ~ Ag | |
| Output Group | IN r, (C) | r ← (C) if r = 110 only the flags will be affected | 1 | י ה | x | T | x | P | 0 | • | 11 101 101 01 r 000 | ED | 2. | 3 | 12 | Acc. to $A_8 \sim A_{15}$ C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$ | |
| | INI | (HL) - (C) B - B - 1 | X | | x | x | x | x | 1 | • | 11 101 101 10 100 010 | ED A2 | 2 | 4 | 16 | C to A ₀ ~ A ₇ B to A ₈ ~ A ₁₅ | |
| | INIR | HL = HL + 1 (HL) = (C) B = B - 1 HL = HL + 1 | x | 1 | x | x | x | X | 1 | • | 11 101 101 10 110 010 | ED B2 | 2 2 | 5 (∐fB≠0) 4 | 21 16 | C to A ₀ ~ A ₇ B to A ₈ ~ A ₁₅ | |
| | | Repeat until B = 0 | | _ | | | | | | | | | 2 | (lf B = 0) | 10 | | |
| | IND | (HL) - (C) B - B - 1 | x | Ť, | x | x | x | x | 1 | • | 11 101 101 10 101 010 | ED AA | 2 | 4 | 16 | C to A ₀ ~ A ₇ B to A ₈ ~ A ₁₅ | |
| | INDR | HL - HL - 1 $(HL) - (C)$ $B - B - 1$ $HL - HL - 1$ Repeat until | x | 1 | x | x | X | x | ļ | • | 11 101 101 10 111 010 | ED BA | 2 | 5 (lf B≠0) 4 (lf B=0) | 21 16 | C to A ₀ - A ₇ B to A ₈ - A ₁₅ | |
| | OUT (n), Å | B = 0 (n) - A | | • | x | • | x | • | | • | 11 010 011 | D3 | 2 | 3 | ú | n to Ag ~ A7 | |
| | OUT (C), r | (C) - r | • | • | x | • | x | • | • | • | 11 101 101 01 r 001 | ED | 2 | 3 | 12 | n to $A_0 \sim A_7$ Acc. to $A_8 \sim A_{15}$ C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$ | |
| | OUTI | | X | Ŷ | x | x | x | x | 1 | • | 11 101 101 10 100 011 | ED A3 | 2 | 4 | 16 | C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$ | |
| | OTIR | HL - HL + 1 $(C) - (HL)$ $B - B - 1$ $HL - HL + 1$ Repeat until | X | 1 | x | X | x | x | 1 | • | 11 101 101 10 110 011 | ED B3 | 2 2 | 5 (If B≠0) 4 (If B=0) | 21 16 | C to $A_0 - A_7$ B to $A_8 - A_{15}$ | |
| | OUTD | B = 0 (C) - (HL) B - B - 1 | x | ۵ ۱ | x | x | X | x | 1 | • | 11 101 101 10 101 011 | | 2 | 4 | 16 | C to $A_0 \sim A_7$ B to $A_8 \sim A_{15}$ | |

NOTE: If the result of B - 1 is zero the Z flag is set, otherwise it is reset.

| Input and Output Group | Mnemonic | Symbolic Operation | | 8 | z | | Fla H | ga | P/V | M | | Opcode 8 543 210 Hex | No.ef Bytes | No.of M Cycles | | Comments | 2 |
|---------------------------|--|--|--|--|---|--------------------------------------|--|--------------------------|---|--|-----------------|--|--|--|--|---|----------------------|
| (Continued) | OTDR | $\begin{array}{l} (C) \leftarrow (HL) \\ B - B - 1 \\ HL \leftarrow HL - 1 \\ Repeat until \\ B = 0 \end{array}$ | | x | 1 | X | X | x | x | 1 | | 11 101 101 ED 10 111 011 | 2 2 | 5 (liB≠0) 4 (liB=0) | 16 | C to A ₀ ~ A ₇ B to A ₈ ~ A ₁₅ | |
| Summary of Flag | Instruction | | Dy S | z | | н | | P/V | N | Ър С | Com | aments | | | | | |
| Operation | ADD A. s; AD | | 1 | 1 | x | ŧ. | X | ۷ | 0 | 1 | | add or add wit | | | | | |
| • | SUB s; SBC A AND s | , s; CP s; NEG | 1 | ; | X X | 1 | X X | V P | 1 | 01 | | | sct with | carry, col | mpare and | d negate accumulator. | |
| | OR s, XOR | | ÷ | i i | x | ò | x | Ρ | ŏ | ŏ | | cal operations. | | | | | |
| | INC a | | 1 | 1 | X X | 1 | X | v | 0 | : | | t increment. decrement. | | | | | |
| | DEC s ADD DD, ss | | : | : | x | 1 X | X X | ÷ | 0 | ÷ | | add. | | | | | |
| | ADC HL, ss | | 1 | 1 | х | х | х | ۷ | Ó | i | | it add with carr | | | | | |
| | SBC HL, 56 RLA, RLCA, R | DA. DOCA | 1 | 1 | X X | X | X | ۷. | 1 | 1 | | at subtract with te accumulator. | carry. | | | | |
| | RL m; RLC m; | | 1 | ī | î. | ŏ | î. | P | ő | ÷ | | te and shift loca | tions. | | | | |
| | RRC m; SLA | im; | | | | | | | | | | | | | | | |
| | SRA m; SRL RLD; RRD | m | 1 | | x | 0 | х | Р | 0 | | Rota | te digit left and | right. | | | | |
| | DAA | | i. | i. | X X | 1 | х | P | • | 1 | Deci | imal adjust accu | mulator | | | | |
| | CPL | | : | : | x | 1 | X X | : | 1 | 1 | | plement accum carry. | ulator. | | | | |
| | SCF | | : | : | x | x | x | • | ŏ | ÷ | | oplement carry. | | | | | |
| | IN r (C) | | 1 | 1 | X | 0 | х | Ρ | 0 | •. | | it register indire | | | | | |
| | INI, IND, OUT INIR: INDR; O | | X | 1 | X X | X | X X | X | 1 | ::} | Bloc | k input and out | put.Z = | 0 ıf B ≠ | 0 otherw | ise Z = 0. | |
| | LDI; LDD | TIN, OTON | | χ. | х | ô | х | î | 0 | •) | Bloc | k transfer instru | ctions 1 | | BC ≠ 0 | otherwise P/V = 0. | |
| | LDIR; LDDR | | X X | X | x | 0 X | X | 0 | 0 | :' | | | | | | wherwise Z = 0. P/V = | 1 |
| | CPI; CPIR; CP | D; CPDR | x | 1 | | x | | | 4 | • | if I | BC # 0, otherw | ise P/V | = 0. | | | |
| | LDA, I, LDA, BIT 5, s | , R | ¦ X | 1 1 | X X | 0 1 | X X | IFF X | 0 0 | : | The The | content of the ir state of bit b of | terrupt location | enable flij s is copie | p-flop (IF ed into the | F) is copied into the P/V 2 flag. | fleg. |
| Symbolic | Symbol | | | Öpe | | ion | | | | | | Symbol | | | | Operation | |
| Notation | | ign flag. S = | | | | | the | res | ult | is 1. | | 1 | The | lac is a | | according to the r | esult of th |
| Notation | | ero flag. Z = | | | | | | | | | is 0. | | | ation. | | | |
| | | arity or overflo | | | | | | | | | | • | The | lag is u | nchang | ed by the operatio | n. |
| | C C | V) share the sa | me | flag: | Lo | gica | al or | era | atio | ns af | lect | 0 | | | | the operation. | |
| | | his flag with th | | | | | | | | | | 1 | The i | lag is s | et by th | e operation. | |
| | | rithmetic oper | | | | | | | | | | х | | lag is a | | | |
| | | verflow of the | resu | lt. I | f P/ | V ho | olds | pa | rity, | P/V | = | v | P/V 9 | lug affe | cted ac | cording to the ove | rflow res |
| | | | | one | erati | on i | s ev | en, | , P/' | V = | 0 if | | | a operat | | | |
| | 1 | if the result of | t the | | | wor | flow | | | | | Р | | | | cording to the par | ity result |
| | 1 r | esult is odd. If | P/V | hol | | | | | | uonfl | ow. | | | peration | | | |
| | l r ti | esult is odd. If he result of the | P/V ope | hol rati | on p | огос | luce | | | | | | | | | | D D 17 |
| | l r ti H H | esult is odd. If he result of the falf-carry flag. | P/V ope H = | hol rati = 1 | on p if th | oroc e ac | luce dd o | r si | ubtr | act | | r | Any | one of t | he CPU | registers A, B, C, | D, E, H, |
| | l T H H | esult is odd. If he result of the falf-carry flag. operation produ | P/V e ope H = uced | hol erati = 1 a c | on p if th arry | oroc e ac | luce dd o | r si | ubtr | act | | r s | Any Any | one of ti 8-bit lo | he CPU cation f | or all the addressir | D, E, H, ng modes |
| | l T H H E | esult is odd. If he result of the falf-carry flag. operation produ- bit 4 of the account | P/V e ope H = uced umu | hol erati = 1 a c ator | on p if th arry r. | e ac int | luce dd o o or | r si bo | ubtr orro | act w fro | m | s | Any Any allow | one of t 8-bit lo red for 1 | he CPU cation f the part | or all the addressir icular instruction. | ng modes |
| | l H H N A | esult is odd. If he result of the Half-carry flag. operation produ- bit 4 of the acco Add/Subtract fl | P/V ope H = uced umul lag. | hol erati = 1 a c ator | on p if th arry r. | e ac int | luce dd o o or | r si bo | ubtr orro | act w fro | m | | Any Any allow Any | one of t 8-bit lo ed for 1 16-bit lo | he CPU cation f the part ocation | or all the addressir icular instruction. for all the address | ng modes |
| | l H H N K | esult is odd. If he result of the Half-carry flag. operation produ- bit 4 of the accu- Add/Subtract fl ion was a subtr | P/V e ope H = uced umul lag. ract. | hol erati a c lator N = | on p if th arry r. 1 i | e ac int | luce id o o or e pr | r si bo evi | ubtr orro ous | act w fro ope: | m a- | S 55 | Any Any allow Any allow | one of t 8-bit loo ed for t 16-bit lo ed for t | he CPU cation f the part ocation that inst | or all the addressir icular instruction. for all the address ruction. | ng modes ing mode |
| | H H H N H K K K K K K K K K K K K K K K | esult is odd. If he result of the Half-carry flag. operation produ- bit 4 of the accu- Add/Subtract fl ion was a subtr I and N flags a | P/V e ope H = uced umul lag. ract. are u | hol rati a c lator N = sed | on p if th arry r. : 1 i in c | oroc e ac int f the conj | luce dd o o or e pr | r si bo evi | ubtro orro ous a wit | act wfro ope: hth | om 'a- e | s ss ii _ | Any Any allow Any allow Any | one of t 8-bit loved for 1 16-bit loved for 1 one of t | he CPU cation f the part ocation that inst that inst | or all the addressir icular instruction. for all the address | ng modes ing mode |
| | I H H N H & N H & N H & N H & N H & M H d | esult is odd. If he result of the Half-carry flag, operation produ- bit 4 of the accu- Add/Subtract fl ion was a subtr I and N flags a lecimal adjust | P/V e ope H = uced umul lag. ract. are u instr | hol rati a c lator N = sed uction | on p if th arry : 1 i in c on (| int f the DA | luce dd o o or e pr unct A) to | evi ion | ubtr orro ous wii rope | act wfro ope: th th erly | a- e cor- | s ss ii ~ R | Any Any allow Any allow Any Refre | one of the sed for the l6-bit low red for the red for the one of the resh court | he CPU cation f the part ocation that inst the two inter. | or all the addressir icular instruction. for all the address ruction. index registers IX | ng modes ing mode |
| | l H H N H & N H & N H & N H & N H & N H & N H H H H | esult is odd. If he result of the falf-carry flag. operation produ- bit 4 of the accr Add/Subtract fl ion was a subtr f and N flags a lecimal adjust ect the result ii | P/V ope H = uced umul lag. ract. are u instr nto p | hol rati a c lator N = sed uction pack | on p if th arry in c on (ced | int f the conj DA. BCI | luce dd o o or e pr unct A) to D for | r si bo evi ior | ubtr orro ous n wit rope at fo | act w fro ope: th th erly llow | a- e cor- | s ss ii R n | Any Any allow Any allow Any Refre 8-bit | one of the ed for the l6-bit low red for the red for the one of the resh court value in | he CPU cation f the part ocation that inst that inst the two nter. n range | or all the addressir icular instruction. for all the address iruction. index registers IX < < 0, 255 > . | ng modes ing mode |
| | I H H H H & N H & N H & N H & N H & N H M H K H H H H H H H H H H H H H H H H | esult is odd. If he result of the falf-carry flag. opperation produ- oit 4 of the accr Add/Subtract fl ion was a subtr I and N flags a lecimal adjust ect the result is ddition or sub | P/V e ope H = uced umul lag. ract. are u instr nto p tract | hol rati a c lator N = sed uction | on p if th arry in c on (ced | int f the conj DA. BCI | luce dd o o or e pr unct A) to D for | r si bo evi ior | ubtr orro ous n wit rope at fo | act w fro ope: th th erly llow | a- e cor- | s ss ii ~ R | Any Any allow Any allow Any Refre 8-bit | one of the ed for the l6-bit low red for the red for the one of the resh court value in | he CPU cation f the part ocation that inst that inst the two nter. n range | or all the addressir icular instruction. for all the address ruction. index registers IX | ng modes ing mode |
| | I H H H H & N H & N H & N H & N H & F G G F F | esult is odd. If he result of the falf-carry flag. operation produ- bit 4 of the accr Add/Subtract fl ion was a subtr f and N flags a lecimal adjust ect the result ii | P/V e ope H = uced umul lag. ract. are u instr nto p tract | hol ration a c lator N = sed uction ion | on p if th arry : 1 i in c on (ced usir | f the conj DA. BCI | luce dd o o or e pr unct A) to pera | evi ior ma | ous ous rope at fo s wi | act w fro ope: th th erly llow ith | e cor- ng | s ss ii R n | Any Any allow Any allow Any Refre 8-bit | one of the ed for the l6-bit low red for the red for the one of the resh court value in | he CPU cation f the part ocation that inst that inst the two nter. n range | or all the addressir icular instruction. for all the address iruction. index registers IX < < 0, 255 > . | ng modes ing mode |

4





Pin Description **A₀-A₁₅.** System Address Bus (output, 3-state). Addresses generated by the DMA are sent to both source and destination ports (main memory or I/O peripherals) on these lines.

BAI. Bus Acknowledge In (input, active Low). Signals that the system buses have been released for DMA control. In multiple-DMA configurations, the BAI pin of the highest priority DMA is normally connected to the Bus Acknowledge pin of the CPU. Lower-priority DMAs have their BAI connected to the BAO of a higher-priority DMA.

BAO. Bus Acknowledge Out (output, active Low). In a multiple-DMA configuration, this pin signals that no other higher-priority DMA has requested t¹ $_{2}$ system buses. BAI and BAO form a daisy chain for multiple-DMA priority resolution over bus control.

BUSREQ. Bus Request (bidirectional, active Low, open drain). As an output, it sends requests for control of the system address bus, data bus and control bus to the CPU. As an input, when multiple DMAs are strung together in a priority daisy chain via BAI and BAO, it senses when another DMA has requested the buses and causes this DMA to refrain from bus requesting until the other DMA is finished. Because it is a bidirectional pin, there cannot be any buffers between this DMA and any other DMA. It can, however, have a buffer between it and the CPU because it is unidirectional into the CPU. A pull-up resistor is connected to this pin.

CE/WAIT. Chip Enable and Wait (input, active Low). Normally this functions only as a CE line, but it can also be programmed to serve a WAIT function. As a CE line from the CPU, it becomes active when WR and IORQ

Z8410 Z80° DMA Direct Memory Access Controller

are active and the I/O port address on the system address bus is the DMA's address, thereby allowing a transfer of control or command bytes from the CPU to the DMA. As a WAIT line from memory or I/O devices, after the DMA has received a bus-request acknowledge from the CPU, it causes wait states to be inserted in the DMA's operation cycles thereby slowing the DMA to a speed that matches the memory or I/O device.

CLK. System Clock (input). Standard Z-80 single-phase clock at 2.5 MHz (Z-80 DMA) or 4.0 MHz (Z-80A DMA). For slower system clocks, a TTL gate with a pullup resistor may be adequate to meet the timing and voltage level specification. For higher-speed systems, use a clock driver with an active pullup to meet the V_{IH} specification and risetime requirements. In all cases there should be a resistive pullup to the power supply of 10K ohms (max) to ensure proper power when the DMA is reset.

D₀-D₇. System Data Bus (bidirectional, 3-state). Commands from the CPU, DMA status, and data from memory or I/O peripherals are transferred on these lines.

IEI. Interrupt Enable In (input, active High). This is used with IEO to form a priority daisy chain when there is more than one interruptdriven device. A High on this line indicates that no other device of higher priority is being serviced by a CPU interrupt service routine.

IEO. Interrupt Enable Out (output, active High). IEO is High only if IEI is High and the CPU is not servicing an interrupt from this DMA. Thus, this signal blocks lower-priority devices from interrupting while a higherpriority device is being serviced by its CPU interrupt service routine.

INT/PULSE. Interrupt Request (output, active Low, open drain). This requests a CPU interrupt. The CPU acknowledges the interrupt by pulling its IORQ output Low during an MI cycle. It is typically connected to the INT pin of the CPU with a pullup resistor and tied to all other INT pins in the system. This pin can also be used to generate periodic pulses to an external device. It can be used this way only when the DMA is bus master (i.e., the CPU's BUSREQ and BUSACK lines are both Low and the CPU cannot see interrupts).

IORQ. Input/Output Request (bidirectional, active Low, 3-state). As an input, this indicates that the lower half of the address bus holds a valid I/O port address for transfer of control or status bytes from or to the CPU, respectively;

this DMA is the addressed port if its \overline{CE} pin and its \overline{WR} or \overline{RD} pins are simultaneously active. As an output, after the DMA has taken control of the system buses, it indicates that the 8-bit or 16-bit address bus holds a valid port address for another I/O device involved in a DMA transfer of data. When \overline{IORQ} and $\overline{M1}$ are both active simultaneously, an interrupt acknowledge is indicated.

Mi. Machine Cycle One (input, active Low). Indicates that the current CPU machine cycle is an instruction fetch. It is used by the DMA to decode the return-from-interrupt instruction (RETI) (ED-4D) sent by the CPU. During twobyte instruction fetches, MI is active as each opcode byte is fetched. An interrupt acknowledge is indicated when both MI and IORQ are active.

MREQ. Memory Request (output, active Low, 3-state). This indicates that the address bus holds a valid address for a mémory read or write operation. After the DMA has taken control of the system buses, it indicates a DMA

Programming The Z-80 DMA has two programmable fun-

damental states: (1) an enabled state, in which it can gain control of the system buses and direct the transfer of data between ports, and (2) a disabled state, in which it can initiate neither bus requests nor data transfers. When the DMA is powered up or reset by any means. it is automatically placed into the disabled state. Program commands can be written to it by the CPU in either state, but this automatically puts the DMA in the disabled state. which is maintained until an enable command is issued by the CPU. The CPU must program the DMA in advance of any data search or transfer by addressing it as an I/O port and sending a sequence of control bytes using an Output instruction (such as OTIR for the Z-80 CPU)

Writing. Control or command bytes are written into one or more of the Write Register groups (WR0-WR6) by first writing to the base register byte in that group. All groups have base registers and most groups have additional associated registers. The associated registers in a group are sequentially accessed by first writing a byte to the base register containing register-group identification and pointer bits (1's) to one or more of that base register's associated registers.

This is illustrated in Figure 8b. In this figure, the sequence in which associated registers within a group can be written to is shown by the vertical position of the associated registers. For example, if a byte written to the DMA contains the bits that identify WR0 (bits D0, D1 and D7), and also contains 1's in the bit positions that point to the associated "Port A Starting Address (low byte)" and "Port A Starting Address (low byte)," then the next two bytes written to the DMA will be stored in these two registers, in that order.

transfer request from or to memory.

RD. Read (bidirectional, active Low, 3-state). As an input, this indicates that the CPU wants to read status bytes from the DMA's read registers. As an output, after the DMA has taken control of the system buses, it indicates a DMA-controlled read from a memory or I/O port address.

RDY. Ready (input, programmable active Low or High). This is monitored by the DMA to determine when a peripheral device associated with a DMA port is ready for a read or write operation. Depending on the mode of DMA operation (Byte, Burst or Continuous), the RDY line indirectly controls DMA activity by causing the BUSREQ line to go Low or High.

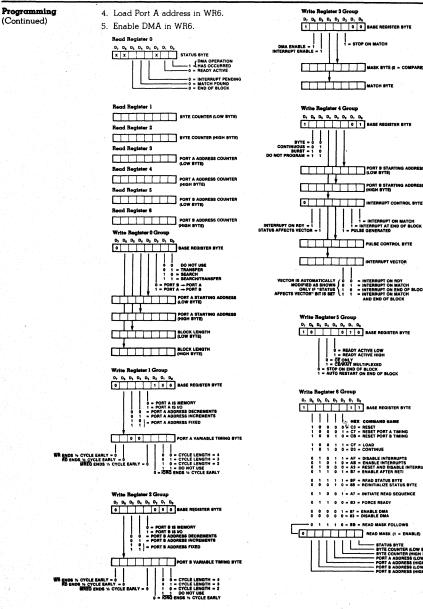
WR. Write (bidirectional, active Low, 3-state). As an input, this indicates that the CPU wants to write control or command bytes to the DMA write registers. As an output, after the DMA has taken control of the system buses, it indicates a DMA-controlled write to a memory or I/O port address.

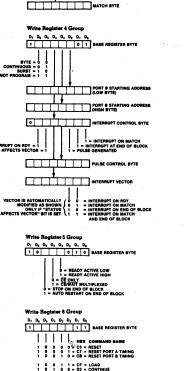
Reading. The Read Registers (RR0-RR6) are read by the CPU by addressing the DMA as an I/O port using an Input instruction (such as INIR for the Z-80 CPU). The readable bytes contain DMA status, byte-counter values, and port addresses since the last DMA reset. The registers are always read in a fixed sequence beginning with RRO and ending with RR6. However, the register read in this sequence is determined by programming the Read Mask in WR6. The sequence of reading is initialized by writing an Initiate Read Sequence or Set Read Status command to WR6. After a Reset DMA. the sequence must be initialized with the Initiate Read Sequence command or a Read Status command. The sequence of reading all registers that are not excluded by the Read Mask register must be completed before a new Initiate Read Sequence or Read Status command.

Fixed-Address Programming. A special circumstance arises when programming a destination port to have a fixed address. The load command in WR6 only loads a fixed address to a port selected as the source, not to a port selected as the destination. Therefore, a fixed destination address must be loaded by temporarily declaring it a fixed-source address and subsequently declaring the true source as such, thereby implicitly making the other a destination.

The following example illustrates the steps in this procedure, assuming that transfers are to occur from a variable address source (Port A) to a fixed address destination (Port B):

- Temporarily declare Port B as source in WR0.
- 2. Load Port B address in WR6.
- 3. Declare Port A as source in WR0.





- STOP ON MATCH

1 = AF = DISABLE INTERRUPTS 0 = AB = ENABLE INTERRUPTS 0 = A3 = RESET AND DISABLE 1 = B7 = ENABLE AFTER RETI

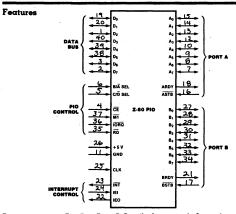
0 1 1 1 1 = BF = READ STATUS BYTE 0 0 0 1 0 = 88 = REINITIALIZE STATUS BYTE 0 1 0 0 1 = A7 = INITIATE READ SEQUENCE 0 1 1 0 0 = 83 = FORCE READY

0 0 0 0 1 = 87 = ENABLE DWA 0 0 0 0 = 83 = DISABLE DWA

- 0 1 1 1 0 - BB - READ MASK FOLLOWS READ MASK (1 = ENABLE)

| STATUS BYTE BYTE COUNTER (LOW BYTE) BYTE COUNTER (HIGH BYTE) PORT A ADDRESS (LOW BYTE) |
|--|
| PORT & ADDRESS (LOW BYTE PORT & ADDRESS (HIGH BYTE PORT & ADDRESS (LOW BYTE PORT & ADDRESS (HIGH BYTE |

Zilog



Pin Description

Ag-A7. Port A Bus (bidirectional, 3-state). This 8-bit bus transfers data, status, or control information between Port A of the PIO and a peripheral device. Ao is the least significant bit of the Port A data bus.

ARDY. Register A Ready (output, active High). The meaning of this signal depends on the mode of operation selected for Port A as follows:

Output Mode. This signal goes active to indicate that the Port A output register has been loaded and the peripheral data bus is stable and ready for transfer to the peripheral device.

Input Mode. This signal is active when the Port A input register is empty and ready to accept data from the peripheral device.

Bidirectional Mode. This signal is active when data is available in the Port A output register for transfer to the peripheral device. In this mode, data is not placed on the Port A data bus, unless ASTB is active.

Control Mode. This signal is disabled and forced to a Low state.

ASTB. Port A Strobe Pulse From Peripheral Device (input, active Low). The meaning of this signal depends on the mode of operation selected for Port A as follows:

Output Mode. The positive edge of this strobe is issued by the peripheral to acknowledge the receipt of data made available by the PIO.

Input Mode. The strobe is issued by the peripheral to load data from the peripheral into the Port A input register. Data is loaded into the PIO when this signal is active.

Bidirectional Mode. When this signal is active, data from the Port A output register is gated onto the Port A bidirectional data bus. The positive edge of the strobe acknowledges the receipt of the data

Control Mode. The strobe is inhibited internally.

Z8420 Z80° PIO Parallel Input/Output Controller

Bn-B7. Port B Bus (bidirectional, 3-state). This 8-bit bus transfers data, status, or control information between Port B and a peripheral device. The Port B data bus can supply 1.5 mA at 1.5 V to drive Darlington transistors. Bo is the least significant bit of the bus.

B/A. Port B Or A Select (input, High = B). This pin defines which port is accessed during a data transfer between the CPU and the PIO. A Low on this pin selects Port A; a High selects Port B. Often address bit An from the CPU is used for this selection function.

BRDY. Register B Ready (output, active High), This signal is similar to ARDY, except that in the Port A bidirectional mode this signal is High when the Port A input register is empty and ready to accept data from the peripheral device.

BSTB. Port B Strobe Pulse From Peripheral Device (input, active Low). This signal is similar to ASTB, except that in the Port A bidirectional mode this signal strobes data from the peripheral device into the Port A input register.

C/D. Control Or Data Select (input, High = C). This pin defines the type of data transfer to be performed between the CPU and the PIO. A High on this pin during a CPU write to the PIO causes the Z-80 data bus to be interpreted as a command for the port selected by the B/\overline{A} Select line. A Low on this pin means that the Z-80 data bus is being used to transfer data between the CPU and the PIO. Often address bit A1 from the CPU is used for this function.

CE. Chip Enable (input, active Low). A Low on this pin enables the PIO to accept command or data inputs from the CPU during a write cycle or to transmit data to the CPU during a read cycle. This signal is generally decoded from four I/O port numbers for Ports A and B, data, and control. CLK. System Clock (input). The Z-80 PIO uses

the standard single-phase Z-80 system clock.

Do-D7. Z-80 CPU Data Bus (bidirectional, 3-state). This bus is used to transfer all data and commands between the Z-80 CPU and the Z-80 PIO. Do is the least significant bit.

IEI. Interrupt Enable In (input, active High). This signal is used to form a priority-interrupt daisy chain when more than one interruptdriven device is being used. A High level on this pin indicates that no other devices of higher priority are being serviced by a CPU interrupt service routine.

Pin Description (Continued)

High). The IEO signal is the other signal required to form a daisy chain priority scheme. It is High only if IEI is High and the CPU is not servicing an interrupt from this PIO. Thus this signal blocks lower priority devices from interrupting while a higher priority device is being serviced by its CPU interrupt service routine.

IEO. Interrupt Enable Out (output, active

INT. Interrupt Request (output, open drain, active Low). When INT is active the Z-80 PIO is requesting an interrupt from the Z-80 CPU. **IORQ.** Input/Output Request (input from Z-80 CPU, active Low). IORQ is used in conjunction with B/A, C/D, CE, and RD to transfer commands and data between the Z-80 CPU and the Z-80 PIO. When CE, RD, and IORQ are active, the part addressed hw RA transfer

active, the port addressed by B/\overline{A} transfers data to the CPU (a read operation). Conversely, when CE and IORQ are active but \overline{RD} is not, the port addressed by B/\overline{A} is written into from the CPU with either data or control

Programming Mode 0, 1, or 2. (Byte Input, Output, or Bidirectional). Programming a port for Mode

0, 1, or 2 requires two words per port. These words are: **A Mode Control Word**. Selects the port operating mode

(Figure 6). This word may be written any time. An Interrupt Vector. The Z-80 PIO is designed for use with

the Z-80 CPU in interrupt Mode 2 (Figure 7). When interrupts are enabled, the PIO must provide an interrupt vector.

Mode 3. (*Bit Input/Output*). Programming a port for Mode 3 operation requires a control word, a vector (if interrupts are enabled), and three additional words, described as follows:

I/O Register Control. When Mode 3 is selected, the mode control word must be followed by another control word that sets the I/O control register, which in turn defines which port lines are inputs and which are outputs (Figure 8).

Interrupt Control Word. In Mode 3, handshake is not used. Interrupts are generated as a logic function of the input signal levels. The interrupt control word sets the logic conditions and the logic levels required for generating an interrupt. Two logic conditions or functions are available: AND (if all input bits change to the active level, an interrupt is triggered), and OR (if any one of the input bits changes to the active level, an interrupt is triggered). Bit D₆ sets the logic function, as shown in Figure 9. The active level is controlled by Bit D₅.

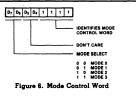
Mask Control Word. This word sets the mask control register, allowing any unused bits to be masked off. If any bits are to be masked, then D_4 must be set. When D_4 is set, the next word written to the port must be a mask control word (Figure 10).

Interrupt Disable. There is one other control word which can be used to enable or disable a port interrupt. It can be used without changing the rest of the interrupt control word (Figure 11).

information, as specified by C/\overline{D} . Also, if IORQ and \overline{M} are active simultaneously, the CPU is acknowledging an interrupt; the interrupting port automatically places its interrupt vector on the CPU data bus if it is the highest priority device requesting an interrupt.

MI. Machine Cycle (input from CPU, active Low). This signal is used as a sync pulse to control several internal PIO operations. When both the MI and RD signals are active, the Z-80 CPU is fetching an instruction from memory. Conversely, when both MI and IORQ are active, the CPU is acknowledging an interrupt. In addition, MI has two other functions within the Z-80 PIO: it synchronizes the PIO interrupt logic; when MI occurs without an active RD or IORQ signal, the PIO is reset.

RD. Read Cycle Status (input from Z-80 CPU, active Low). If RD is active, or an L/O operation is in progress, RD is used with B/A, C/D, CE, and IORQ to transfer data from the Z-80 PIO to the Z-80 CPU.



D7 D6 D5 D4 D3 D2 D1 0



Figure 7. Interrupt Vector Word

D7 D8 D5 D4 D3 D2 D1 D0

6 SETS BIT TO OUTPUT 1 SETS BIT TO INPUT Figure 8. I/O Register Control Word

D7 D8 D5 D4 0 1 1 1

| IDENTIFIES INTERRUPT CONTROL WORD |
|---|
| $D_4 = 0$ NO MASK WORK FOLLOWS $D_4 = 1$ MASK WORD FOLLOWS |
| D5 = 0 ACTIVE LEVEL IS LOW D5 = 1 ACTIVE LEVEL IS HIGH |
| $D_6 = 0$ INTERRUPT ON OR FUNCTION $D_6 = 1$ INTERRUPT ON AND FUNCTION |
| D7 = 0 INTERRUPT DISABLED D7 = 1 INTERRUPT ENABLED* |

*NOTE: THE PORT IS NOT ENABLED UNTIL THE INTERRUPT ENABLE IS FOLLOWED BY AN ACTIVE M1.

Figure 9. Interrupt Control Word



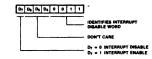


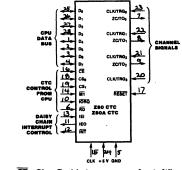
Figure 10. Mask Control Word

Zilog

Features

Pin

Description



CE. Chip Enable (input, active Low). When enabled the CTC accepts control words, interrupt vectors, or time constant data words from the data bus during an I/O write cycle; or transmits the contents of the down-counter to the CPU during an I/O read cycle. In most applications this signal is decoded from the eight least significant bits of the address bus for any of the four I/O port addresses that are mapped to the four counter-timer channels.

CLK. System Clock (input). Standard singlephase Z-80 system clock.

CLE/TRG₀-CLE/TRG₃. External Clock/Timer Trigger (input, user-selectable active High or Low). Four pins corresponding to the four Z-80 CTC channels. In counter mode, every active edge on this pin decrements the down-counter. In timer mode, an active edge starts the timer.

 CS_0-CS_1 . Channel Select (inputs active High). Two-bit binary address code selects one of the four CTC channels for an I/O write or read (usually connected to A_0 and A_1).

D₀-D₇. System Data Bus (bidirectional, 3-state). Transfers all data and commands between the Z-80 CPU and the Z-80 CTC.

IEI. Interrupt Enable In (input, active High). A High indicates that no other interrupting devices of higher priority in the daisy chain are being serviced by the Z-80 CPU.

Figure 11. Interrupt Disable Word **Z8430 Z80° CTC Counter/ Timer Circuit**

IEO. Interrupt Enable Out (output, active High). High only if IEI is High and the Z-80 CPU is not servicing an interrupt from any Z-80 CTC channel. IEO blocks lower priority devices from interrupting while a higher priority interrupting device is being serviced.

INT. Interrupt Request (output, open drain, active Low). Low when any Z-80 CTC channel that has been programmed to enable interrupts has a zero-count condition in its down-counter.

IORQ. Input/Output Request (input from CPU, active Low). Used with CE and RD to transfer data and channel control words between the Z-80 CPU and the Z-80 CTC. During a write cycle, IORQ and CE are active and RD inactive. The Z-80 CTC does not receive a specific write signal; rather, it internally

generates its own from the inverse of an active RD signal. In a read cycle, IORO, CE and RD are active; the contents of the down-counter are read by the Z-80 CPU. If IORO and MI are both true, the CPU is acknowledging an interrupt request, and the highest priority interrupting channel places its interrupt vector on the Z-80 data bus.

Mi. Machine Cycle One (input from CPU, active Low). When **Mi** and **IORQ** are active, the Z-80 CPU is acknowledging an interrupt. The Z-80 CPC then places an interrupt vector on the data bus if it has highest priority, and if a channel has requested an interrupt (INT).

RD. Read Cycle Status (input, active Low). Used in conjunction with IORQ and CE to transfer data and channel control words between the Z-80 CPU and the Z-80 CTC.

RESET. Reset (input active Low). Terminates all down-counts and disables all interrupts by resetting the interrupt bits in all control registers; the ZC/TO and the Interrupt outputs go inactive; IEO reflects IEI; D_0 - D_7 go to the high-impedance state.

ZC/TO₀-ZC/TO₂. Zero Count/Timeout (output, active High). Three ZC/TO pins corresponding to Z-80 CTC channels 2 through 0 (Channel 3 has no ZC/TO pin). In both counter and timer modes the output is an active High pulse when the down-counter decrements to zero.

Programming

Each Z-80 CTC channel must be pro-

grammed prior to operation. Programming consists of writing two words to the I/O port that corresponds to the desired channel. The first word is a control word that selects the operating mode and other parameters; the second word is a time constant, which is a binary data word with a value from 1 to 256. Å time constant word must be preceded by a channel control word.

After initialization, channels may be reprogrammed at any time. If updated control and time constant words are written to a channel during the count operation, the count continues to zero before the new time constant is loaded into the counter.

If the interrupt on any Z-80 CTC channel is enabled, the programming procedure should also include an interrupt vector. Only one vector is required for all four channels, because the interrupt logic automatically modifies the vector for the channel requesting service.

A control word is identified by a 1 in bit 0. A 0 in bit 2 indicates a time constant word is to follow. Interrupt vectors are always addressed to Channel 0, and identified by a 0 in bit 0.

Addressing. During programming, channels are addressed with the channel select pins CS_1 and CS_2 . A 2-bit binary code selects the appropriate channel as shown in the following table.

| | Channel | CS_1 | CS ₀ | |
|------|---------|--------|-----------------|--|
| | 0 | 0 | 0 | |
| | 1 | 0 | 1 | |
| | 2 | 1 | 0 | |
| | 3 | 1 | 1 | |

Reset. The CTC has both hardware and software resets. The hardware reset terminates all down-counts and disables all CTC interrupts by resetting the interrupt bits in the control registers. In addition, the ZC/TO and Interrupt outputs go inactive, IEO reflects IEI, and

D₀-D₇ go to the high-impedance state. All channels must be completely reprogrammed after a hardware reset.

The software reset is controlled by bit 1 in the channel control word. When a channel receives a software reset, it stops counting. When a software reset is used, the other bits in the control word also change the contents of the channel control register. After a software reset a new time constant word must be written to the same channel.

If the channel control word has both bits D_1 and D_2 set to 1, the addressed channel stops operating, pending a new time constant word. The channel is ready to resume after the new constant is programmed. In timer mode, if $D_3 = 0$, operation is triggered automatically when the time constant word is loaded.

Channel Control Word Programming. The channel control word is shown in Figure 5. It sets the modes and parameters described below.

Interrupt Enable. D₇ enables the interrupt, so that an interrupt output (\overline{INT}) is generated at zero count. Interrupts may be programmed in either mode and may be enabled or disabled at any time.

Operating Mode. D₆ selects either timer or counter mode. Prescaler Factor. (Timer Mode Only). D₅

selects factor—either 16 or 256.

Trigger Slope. D₄ selects the active edge or slope of the CLK/TRG input pulses. Note that reprogramming the CLK/TRG slope during operation is equivalent to issuing an active edge. If the trigger slope is changed by a control word update while a channel is pending operation in timer mode, the result is the same as a CLK/TRG pulse and the timer starts. Similarly, if the channel is in counter mode, the counter decrements.

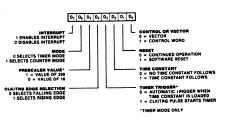


Figure 5. Channel Control Word

Programming Trigger Mode (Timer Mode Only). D_3 selects (Continued) the trigger mode for timer operation. When D

the trigger mode for timer operation. When D_3 is reset to 0, the timer is triggered automatically. The time constant word is programmed during an I/O write operation, which takes one machine cycle. At the end of the write operation there is a setup delay of one clock period. The timer starts automatically (decrements) on the rising edge of the second clock pulse (T₂) of the machine cycle following the write operation. Once started, the timer runs continuously. At zero count the timer reloads automatically and continues counting without interruption or delay, until stopped by a reset.

When D_3 is set to 1, the timer is triggered externally through the CLK/TRG input. The time constant word is programmed during an I/O write operation, which takes one machine cycle. The timer is ready for operation on the rising edge of the second clock pulse (T2) of the following machine cycle. Note that the first timer decrement follows the active edge of the CLK/TRG pulse by a delay time of one clock cycle if a minimum setup time to the rising edge of clock is met. If this minimum is not met, the delay is extended by another clock period. Consequently, for immediate triggering, the CLK/TRG input must precede T₂ by one clock cycle plus its minimum setup time. If the minimum time is not met, the timer will start on the third clock cycle (T_3) .

Once started the timer operates continuously, without interruption or delay, until stopped by a reset.

Time Constant to Follow. A 1 in D_2 indicates that the next word addressed to the selected channel is a time constant data word for the time constant register. The time constant word may be written at any time.

A 0 in D₂ indicates no time constant word is to follow. This is ordinarily used when the channel is already in operation and the new channel control word is an update. A channel will not operate without a time constant value. The only way to write a time constant value is to write a control word with D₂ set.

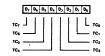


Figure 6. Time Constant Word

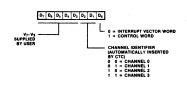


Figure 7. Interrupt Vector Word

Software Reset. Setting D₁ to 1 consets a software reset, which is described in the Reset section. Control Word. Setting Dn to 1 identifies the

word as a control word.

Time Constant Programming. Before a channel can start counting it must receive a time constant word from the CPU. During programming or reprogramming, a channel control word in which bit 2 is set must precede the time constant word to indicate that the next word is a time constant. The time constant word can be any value from 1 to 256 (Figure 6). Note that 00_{16} is interpreted as 256. In time mode, the time interval is controlled

by three factors:

- The system clock period (ϕ)
- The prescaler factor (P), which multiplies the interval by either 16 or 256
- The time constant (T), which is programmed into the time constant register

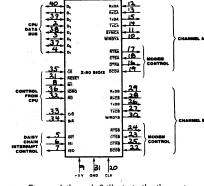
Consequently, the time interval is the product of $\phi \times P \times T$. The minimum timer resolution is $16 \times \phi$ (4 μ s with a 4 MHz clock). The maximum timer interval is $256 \times \phi \times 256$ (16.4 ms with a 4 MHz clock). For longer intervals timers may be cascaded.

Interrupt Vector Programming. If the Z-80 CTC has one or more interrupts enabled, it can supply interrupt vectors to the Z-80 CPU. To do so, the Z-80 CTC must be pre-programmed with the most-significant five bits of the interrupt vector. Programming consists of writing a vector word to the I/O port corresponding to the Z-80 CTC Channel 0. Note that Do of the vector word is always zero, to distinguish the vector from a channel control word. D_1 and D_2 are not used in programming the vector word. These bits are supplied by the interrupt logic to identify the channel requesting interrupt service with a unique interrupt vector (Figure 7). Channel 0 has the highest priority.





Features



Pin Description Figures 1 through 6 illustrate the three pin configurations (bonding options) available in the SIO. The constraints of a 40-pin package make it impossible to bring out the Receive Clock (RxC), Transmit Clock (TxC), Data Terminal Ready (DTR) and Sync (SYNC) signals for both channels. Therefore, either Channel B lacks a signal or two signals are bonded together in the three bonding options offered:

Z-80 SIO/2 lacks SYNCB

Z-80 SIO/1 lacks DTRB

 Z-90 SIO/0 has all four signals, but TxCB and RxCB are bonded together

The first bonding option above (SIO/2) is the preferred version for most applications. The pin descriptions are as follows:

B/ $\overline{\mathbf{A}}$. Channel A Or B Select (input, High selects Channel B). This input defines which channel is accessed during a data transfer between the CPU and the SIO. Address bit A_0 from the CPU is often used for the selection function.

C/D. Control Or Data Select (input, High selects Control). This input defines the type of information transfer performed between the CPU and the SIO. A High at this input during a CPU write to the SIO causes the information on the data bus to be interpreted as a command for the channel selected by B/\overline{A} . A Low at C/\overline{D} means that the information on the data bus is data. Address bit A_1 is often used for this function.

CE. Chip Enable (input, active Low). A Low level at this input enables the SIO to accept command or data input from the CPU during a write cycle or to transmit data to the CPU during a read cycle.

Z8440 Z80° SIO Serial Input/Output Controller

CLK. System Clock (input). The SIO uses the standard Z-80 System Clock to synchronize internal signals. This is a single-phase clock.

CTSA. CTSB. Clear To Send (inputs, active Low). When programmed as Auto Enables, a Low on these inputs enables the respective transmitter. If not programmed as Auto Enables, these inputs may be programmed as general-purpose inputs. Both inputs are Schmitt-trigger buffered to accommodate slowrisetime signals. The SIO detects pulses on these inputs and interrupts the CPU on both logic level transitions. The Schmitt-trigger buffering does not guarantee a specified noiselevel margin.

D₀-D₇. System Data Bus (bidirectional, 3-state). The system data bus transfers data and commands between the CPU and the Z-80 SIO. D₀ is the least significant bit.

DCDA, **DCDB**. Data Carrier Detect (inputs, active Low). These pins function as receiver enables if the SIO is programmed for Auto Enables; otherwise they may be used as general-purpose input pins. Both pins are Schmitt-trigger buffered to accommodate slow; risetime signals. The SIO detects pulses on these pins and interrupts the CPU on both logic level transitions. Schmitt-trigger buffering does not guarantee a specific noise-level margin.

DTRA. DTRB. Data Terminal Ready (outputs, active Low). These outputs follow the state programmed into Z-80 SIO. They can also be programmed as general-purpose outputs. In the Z-80 SIO/1 bonding option, DTRB is

omitted. **IEI.** Interrupt Enable In (input, active High).

This signal is used with IEO to form a priority daisy chain when there is more than one interrupt-driven device. A High on this line indicates that no other device of higher priority is being serviced by a CPU interrupt service routine.

IEO. Interrupt Enable Out (output, active High). IEO is High only if IEI is High and the CPU is not servicing an interrupt from this SIO. Thus, this signal blocks lower priority devices from interrupting while a higher priority device is being serviced by its CPU interrupt service routine.

INT. Interrupt Request (output, open drain, active Low). When the SIO is requesting an interrupt, it pulls INT Low.

IORQ. Input/Output Request (input from CPU, active Low). IORQ is used in conjunction with

Pin Description (Continued) B/\overline{A} , C/\overline{D} , \overline{CE} and \overline{RD} to transfer commands and data between the CPU and the SIO. When \overline{CE} , \overline{RD} and \overline{IORQ} are all active, the channel selected by B/\overline{A} transfers data to the CPU (a read operation). When \overline{CE} and \overline{IORQ} are active but \overline{RD} is inactive, the channel selected by B/\overline{A} is written to by the CPU with either data or control information as specified by C/\overline{D} . If \overline{IORQ} and \overline{MI} are active simultaneously, the CPU is acknowledging an interrupt and the SIO automatically places its interrupt vector on the CPU data bus if it is the highest priority device requesting an interrupt.

MI. Machine Cycle (input from Z-80 CPU, active Low). When MI is active and RD is also active, the Z-80 CPU is fetching an instruction from memory; when MI is active while IORQ is active, the SIO accepts MI and IORQ as an interrupt acknowledge if the SIO is the highest priority device that has interrupted the Z-80 CPU.

RxCA, **RxCB**. Receiver Clocks (inputs). Receive data is sampled on the rising edge of RxC. The Receive Clocks may be 1, 16, 32 or 64 times the data rate in asynchronous modes. These clocks may be driven by the Z-80 CTC Counter Timer Circuit for programmable baud rate generation. Both inputs are Schmitttrigger buffered (no noise level margin is specified).

In the Z-80 SIO/0 bonding option, \overline{RxCB} is bonded together with \overline{TxCB} .

RD. Read Cycle Status (input from CPU, active Low). If \overline{RD} is active, a memory or I/O read operation is in progress. \overline{RD} is used with B/\overline{A} , \overline{CE} and \overline{IORQ} to transfer data from the SIO to the CPU.

RxDA. RxDB. *Receive Data* (inputs, active High). Serial data at TTL levels.

RESET. Reset (input, active Low). A Low RESET disables both receivers and transmitters, forces TxDA and TxDB marking, forces the modem controls High and disables all interrupts. The control registers must be rewritten after the SIO is reset and before data is transmitted or received.

RTSA, **RTSB**. Request To Send (outputs, active Low). When the RTS bit in Write Register 5 (Figure 14) is set, the \overline{RTS} output goes Low. When the RTS bit is reset in the Asynchronous mode, the output goes High after the transmitter is empty. In Synchronous modes, the \overline{RTS} pin strictly follows the state of the RTS bit. Both pins can be used as general-purpose outputs.

SYNCA. SYNCB. Synchronization (inputs/outputs, active Low). These pins can act either as inputs or outputs. In the asynchronous receive mode, they are inputs similar to \overline{CTS} and \overline{DCD} . In this mode, the transitions on these lines affect the state of the Sync/Hunt status

bits in Read Register 0 (Figure 13), but have no other function. In the External Sync mode, these lines also act as inputs. When external synchronization is achieved, SYNC must be driven Low on the second-rising edge of RxC after that rising edge of RxC on which the last bit of the sync character was received. In other words, after the sync pattern is detected. the external logic must wait for two full Receive Clock cycles to activate the SYNC input. Once SYNC is forced Low, it should be kept Low until the CPU informs the external synchronization detect logic that synchronization has been lost or a new message is about to start. Character assembly begins on the rising edge of RxC that immediately precedes the falling edge of SYNC in the External Sync mode.

In the internal synchronization mode (Monosync and Bisync), these pins act as outputs that are active during the part of the receive clock (\overline{RxC}) cycle in which sync characters are recognized. The sync condition is not latched, so these outputs are active each time a sync pattern is recognized, regardless of character boundaries.

In the Z-80 SIO/2 bonding option, SYNCB is omitted.

TxCA. TxCB. Transmitter Clocks (inputs). In asynchronous modes, the Transmitter Clocks may be 1, 16, 32 or 64 times the data rate; however, the clock multiplier for the transmitter and the receiver must be the same. The Transmit Clock inputs are Schmitt-trigger buffered for relaxed rise and fall-time requirements (no noise level margin is specified). Transmitter Clocks may be driven by the Z-80 CTC Counter Timer Circuit for programmable baud rate generation.

In the Z-80 SIO/9 bonding option, \overline{TxCB} is bonded together with RxCB.

TxDA. TxDB. Transmit Data (outputs, active High). Serial data at TTL levels. TxD changes from the falling edge of \overline{TxC} .

W/RDYA, W/RDYB. Wait/Ready A, Wait/ Ready B (outputs, open drain when programmed for Wait function, driven High and Low when programmed for Ready function). These dual-purpose outputs may be programmed as Ready lines for a DMA controller or as Wait lines that synchronize the CPU to the SIO data rate. The reset state is open drain.

Programming

The system program first issues a series of commands that initialize the basic mode of operation and then other commands that qualify conditions within the selected mode. For example, the asynchronous mode. character length, clock rate, number of stop bits, even or odd parity might be set first; then the interrupt mode; and finally, receiver or transmitter enable.

Both channels contain registers that must be programmed via the system program prior to operation. The channel-select input (B/\overline{A}) and the control/data input (C/D) are the commandstructure addressing controls, and are normally controlled by the CPU address bus. Figures 15 and 16 illustrate the timing relationships for programming the write registers and transferring data and status.

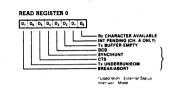
Read Registers. The SIO contains three read registers for Channel B and two read registers for Channel A (RR0-RR2 in Figure 13) that can be read to obtain the status information; RR2 contains the internally-modifiable interrupt vector and is only in the Channel B register set. The status information includes error conditions, interrupt vector and standard communications-interface signals.

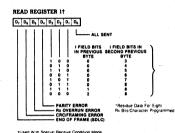
To read the contents of a selected read register other than RRO, the system program must first write the pointer byte to WR0 in exactly the same way as a write register operation. Then, by executing a read instruction, the contents of the addressed read register can be read by the CPU.

The status bits of RRO and RR1 are carefully grouped to simplify status monitoring. For example, when the interrupt vector indicates that a Special Receive Condition interrupt has occurred, all the appropriate error bits can be read from a single register (RR1).

Write Registers. The SIO contains eight write registers for Channel B and seven write registers for Channel A (WR0-WR7 in Figure 14) that are programmed separately to configure the functional personality of the channels; WR2 contains the interrupt vector for both channels and is only in the Channel B register set. With the exception of WR0, programming the write registers requires two bytes. The first byte is to WRO and contains three bits (D_0-D_2) that point to the selected register; the second byte is the actual control word that is written into the register to cor figure the SIO.

WR0 is a special case in that all of the basic commands can be written to it with a single byte. Reset (internal or external) initializes the pointer bits D_0 - D_2 to point to WR0. This implies that a channel reset must not be combined with the pointing to any register.







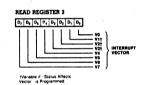


Figure 13. Read Register Bit Functions

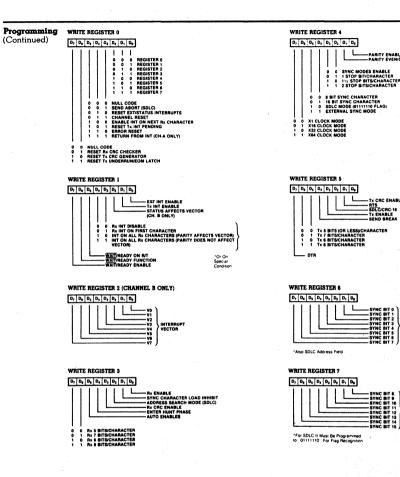


Figure 14. Write Register Bit Functions

TA CRC ENABLE

RTS SDLC/CRC-16

Tx ENABLE

ACTER

TS 806C THEORY OF OPERATION

- SECTION TITLE
 - 1.0 INTRODUCTION
 - 2.0 TS 806C TAPE CONTROLLER BOARD DESCRIPTION
 - 3.0 FUNCTION OF THE SYSTEM
 - 4.0 OPERATION OF THE SYSTEM
 - 5.0 INPUT/OUTPUT PORT ASSIGNMENT AND TIMING
 - 6.0 CONNECTOR DESCRIPTION
 - 7.0 SPECIFICATIONS



1.0 INTRODUCTION

The TS 806C is an intelligent tape-cartridge unit with its own controller designed to be used with the TS 806 service processor. The TS 806C has three ports: one of the ports (50 pin internal connector) is used to interface to the tape drive; one RS 422 port is used to communicate with the service processor; one RS 422 is used to pass the normal communications through to the work station. The service processor controls the communication links between the tape controller and the work station. The TS 806C receives commands from the TS 806 to perform tape read, write or erase functions; all other data is passed through to the work station.

The system block diagram is shown in Figure 1.

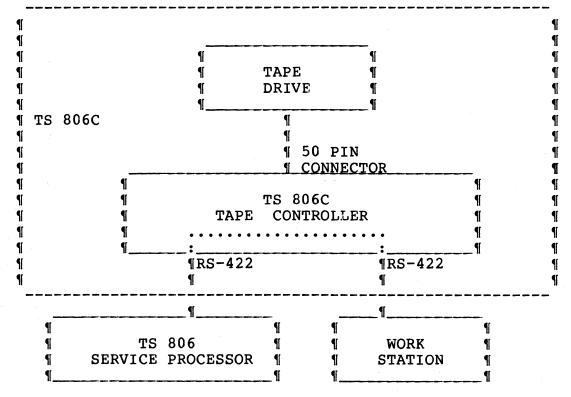


Figure 1 System Block Diagram

2.0 TS 806C TAPE CONTROLLER BOARD DESCRIPTION

The TS 806C controller board contains

Three Z80A family ICs

CPU, CTC and SIO

64 Kbytes of dynamic RAM

4 Kbytes of EPROM Two 8-bit output latches Associated control logic See Figure 2.

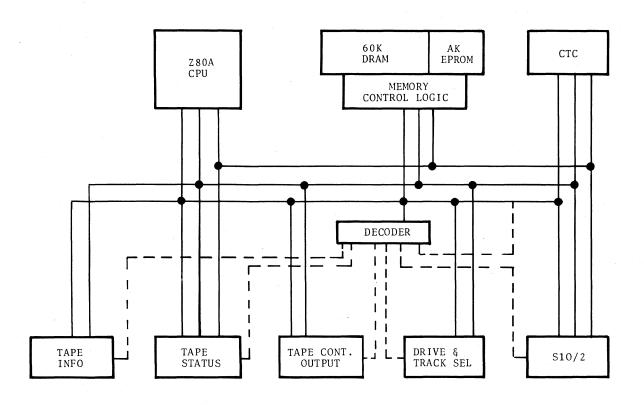


Figure 2 Block Diagram of TS 806C Controller Board

The Z80A ICs are driven by a 4 MHz system or tape clock. This clock is also used to generate the necessary timing for the memory control logic.

The Z80A CPU has a 16-bit address bus which can address up to 64 Kbytes for memory. The lower eight address lines are also used to address up 256 Input/Output devices. During instruction fetch cycles, the CPU sends out refresh addresses. The refresh enable address signals to the memory control which then refreshes 64 Kbytes of dynamic RAM.

A second bus is an eight-bit bidirectional data bus. The third bus is the control bus which includes the following signals: -M1, -MREQ, -IORQ, -RD, -WR, -RFSH, -WAIT, -INT and -RESET. The FS 806C has eight I/O devices: the SIO, CTC, LEDs, a four-bit lipswitch, two tri-state input latches, and two 8-bit output latches.

The Z80A CTC is a four-channel programmable clocking device which can be programmed as a timer or as a counter. The main function of this CTC is to provide transmit clocks for the RS 422 communication interface. The CTC is programmed to reset the Z80 SIO which effects both channels of the SIO.

The Z80A SIO is a dual-channel serial I/O controller. It is programmed in the SDLC mode and clocked at a data rate of 800 Kbits/sec.

The main memory in the TS 806C controller contains eight 64K x l dynamic memory devices (only 60 Kbytes are accessible). There are four kilobytes of EPROM.

3.0 FUNCTION OF THE SYSTEM

The TS 806C generates an 8-MHz clock and a 4-MHz clock. These are supplied to the CPU, SIO, CTC, DMA and memory control logic. After power is first turned on, the system is reset to the idle state and all devices are then programmed for operation.

There are two different types of memory devices in the TS 806C, dynamic RAM and EPROM. See Figure 3.

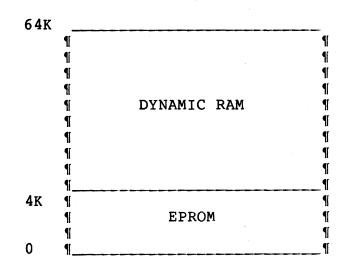


Figure 3 Memory Devices

The EPROM contains software for initialization of the controller, power-on diagnostics, and the download program.

The memory control logic supplies the necessary read/write signals, address latching signals, and refresh signal to the

dynamic RAM. Because of the speed difference between the access of the two types of memory, a one cycle "wait state" is inserted during read/write operations to the EPROM.

The daisy-chain style of communication makes it possible to reset the controller without effecting the link from the service processor to the work station.

4.0 OPERATION OF THE SYSTEM

The 16 MHz clock from the oscillator is counted down to 4 MHz by counter Al4 (74LS 161A) and becomes the system clock frequency. This 4 MHz signal is sent to all the Z80A family devices and the memory control logic. The level of the clock signal is critical to the Z80A devices. The high level of the clock must be between 4.4 V and 5.3 V and the low level must be between -0.3 V and 0.45 V (see Figure 4). A transistor (2N2907) is used to pull the clock's output high level to about five volts.

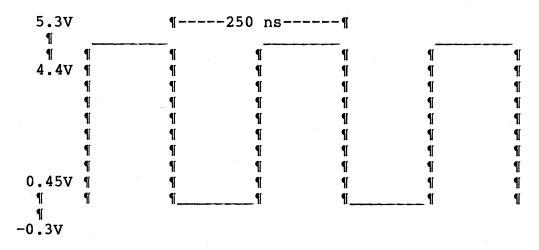


Figure 4 System Clock

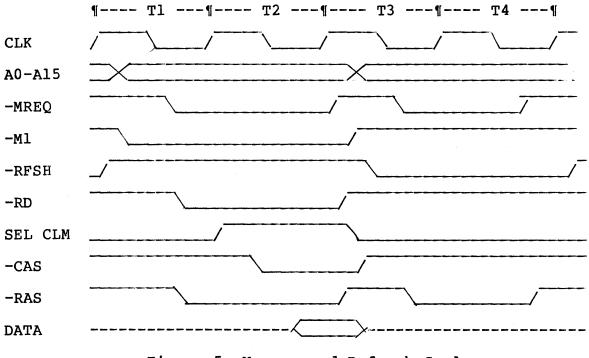
Each dynamic RAM used in the TS 806C is organized as 65536 x l bits; eight dynamic RAMs are used on the controller board. Multiplexed addressing and periodic refreshing are required in the operation of this type of memory. The memory and refresh cycle diagram is shown in Figure 5.

The -MREQ and -RFSH signals from the CPU are sent to the memory control logic. The memory control logic is comprised of the following devices:

Three "D" flip-flops 74S74 (A41, A46) Two nand gates 74S00 (A37) Three nor gates 74LS02 (A38) Two inverter gates 74LS04

Two three-input nor gates 74LS27 (A31)

Two multiplexers, All and Al2 (74S157), are used to multiplex the 16-bit address lines from the CPU into the 8-bit address lines of the dynamic memory. The row address and column address are latched internally by the falling edge of the -RAS and -CAS signals. Each memory cell in the memory device must be refreshed at least every two milliseconds. A serial resistor (33 ohm) is connected to the address line signal, the write signal, the -RAS signal, and the -CAS signal on the input to the dynamic memory. The purpose of these resistors is to reduce the signal undershoot caused from the capacitance of the devices.





The other type of memory used in this system is the EPROM. Since the EPROM has a slower access time compared to the Z80A memory access cycle, one "wait state" is automatically inserted by the "wait control circuit" when the EPROM is accessed. The memory latch timing diagram is shown in Figure 7. The "wait control circuit" is implemented with two "D" flip-flops (74LS74 A34), one three-input nor gate (74LS27 A31), and one inverter gate (74LS04 A30).

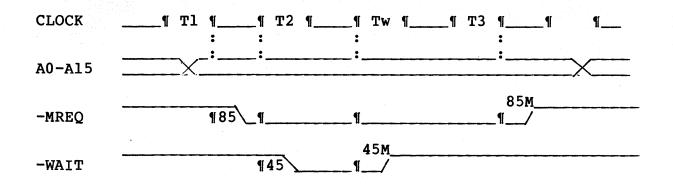


Figure 6 Wait Cycle Inserted Timing

The TS 806C controller board has eight logical I/O devices. They are the LED, CTC, SIO, four-bit dipswitch, two tri-state input buffers, and two output latches. During CPU I/O operations, the address bus is decoded by A26 (74LS138 IC) to select one of the I/O devices. The LED and tape select have the same port address but different times of use.

4.1 TAPE OPERATION

The data latches (A23 and A28) are used to latch the data word that controls the tape drive operation.

4.1.1 Output Latches (from the computer)

The control signals to the tape drive are programmed by two data latches. These latches are write-only ports.

4.1.1.1 Tape Output Latch 2

Tape Output Latch 2 controls the track select and, if in a multiple drive system, the drive select.

Data word for "Tape output latch 2"

¶ -TR2 ¶ -TR1 ¶ -SL4 ¶ -SL2 ¶ -SL1 ¶ X ¶ X ¶ X ¶ ---D07----D06----D05----D04----D03---D02---D01---D00--

Track select is controlled via D07 and D06 according to the following chart with low being true. All heads for a given track (read, write, and erase) are selected simultaneously through the use of these control bits. The last track selection is stored in the tape drive even after deselection by the controller.

| Track | No. | Track | Select | Signals |
|-------|-----|-------|---------|----------|
| | | m | R2 | TR] |
| 1 | | | KZ H | TRI L |

| 2 | \mathbf{L} | H |
|---|------------------|--------------|
| 3 | \mathbf{L}_{i} | \mathbf{L} |
| 4 | H | н |

Tape drive select is controlled via D05, D04, AND D03. Together these bits determine the tape drive address (low true). The address function processes and remains active during any other drive I/O. The following chart shows the address selection.

| Drive No. | Drive | Select | Signals |
|-----------|--------------|--------------|-------------------|
| | SL4 | SL2 | SLl |
| 1 | H | H | · L |
| 2 | Н | L | Н |
| 3 | Н | \mathbf{L} | \mathbf{L}^{-1} |
| 4 | \mathbf{L} | н | H |
| 5 | \mathbf{L} | H | \mathbf{L} |
| 6 | \mathbf{L} | \mathbf{L} | Н |
| 7 | \mathbf{L} | \mathbf{L} | \mathbf{L} |
| 8 | н | н | н |

Data bits D02, D01, and D00 are not used and do not effect the system.

4.1.1.2 Tape Output Latch 1

The Tape Output latch 1 is called the Tape Command/Control Output Latch (14H).

Data Word for Output Latch 1

¶ SLG ¶ X ¶ -RWD ¶ FB ¶ -WEN ¶ -HSP ¶ -FWD ¶ -REV ¶ ---D7----D6-----D5-----D4-----D3-----D1-----D0---

Reverse direction (-REV) is D00. Active low causes the tape to move in the reverse direction.

Forward direction (-FWD) is D01. Active low causes the tape to move in the forward direction.

High speed (-HSP) causes the tape to move at high speed in the direction selected.

Tape motion continues until command signals go false. This can be triggered by the end of the tape (when going forward) or the beginning of the tape (when in reverse). The tape will also stop if both directions are given at the same time, if the ready signal goes inactive, and if a higher priority rewind command is given.

When in high speed motion, the speed will drop to low speed if the upper loading point hole is sensed in the reverse direction or if the upper early warning hole is sensed in the forward direction.

Write enable (-WEN) enables writing and erasing functions for the selected track. The writing or erasing operations is only allowed if the tape cartridge is in the unprotected state. Write enable should be set prior to the tape going in motion and should not be reset until the drive has stopped. At least two milliseconds are needed between the reset of write enable and the selection of a new track. The write enable signal is reset by either reverse or high speed commands.

Feedback (FB) is for general use and can also be read through port OCH at D04.

Rewind (-RWD) causes the tape to be positioned at the beginning at high speed. The drive must be selected to start a rewind but may be deselected after the sequence has started.

Data bit D06 is not used.

Select gate (SLG) allows selection per the select and address codes. It is used to prevent unwanted selection during the change of a select operation.

4.1.2 Input Latches (to the controller)

The status of the tape controller is read through two input ports. These ports are read-only status line which are latched by the tape drive.

4.1.2.1 Tape Status Input Latch 1

Input Latch 1 is called the Tape Status Input Latch (08H).

Data word for the Tape Status Input Latch

¶ RDY ¶ LPS ¶ EWS ¶ FLG ¶ FUP ¶ WEN ¶ BSY ¶ SLD ¶ --D07---D06---D05---D04---D03---D02---D01---D00--

Ready (RDY) is true when the tape cartridge is installed, the sensor lamp is drawing current and five volts is applied to the tape drive.

Load point sensed (LPS) is set and latched when the upper load point hole (the warning for the start of the tape) is passed in the reverse direction. The signal is reset when the load point is passed with the tape going in the forward direction. When this signal is true, the high speed signal is disabled in the reverse direction. At this time, reverse tape motion is allowed to proceed until the beginning of tape hole is encountered. When this happens, the drive stops and only accepts forward commands.

Early warning sensed (EWS) is set and latched when the upper early warning hole (warning for end of the tape) is passed in the forward direction. This signal is reset when the hole is passed in the reverse direction. When this signal is true, the high speed signal is disabled and forward motion is allowed to proceed until the end of tape hole. When this happens, the tape stops and only accepts reverse direction commands.

Flag (FLG) is set when the automatic sequence to position the tape at the beginning has been executed or a rewind command has been completed. This signal is reset by a subsequent receipt of a forward command.

File unprotected (FUP) is true when a tape cartridge is installed and it is in the unprotected state (meaning the tape can be written on).

Write enable (WEN) is true when a write enable is latched within the tape drive.

Busy (BSY) is true when the drive is performing an automatic rewind sequence (a cartridge is initially installed), normal rewind, forward or reverse command. This signal goes true when the command is received and remain true until tape motion has stopped. The time for slow speed commands is 30 milliseconds and for high speed commands is 80 milliseconds.

Selected (SLD) is true when the tape drive has received its proper address.

4.1.2.2 Tape Information Latch 2

Input latch 2 is called the Tape Information Input Latch (OCH).

Data word for Tape Information Input Latch 2

| ¶ SIO-A ¶ DIAG | ¶ DAD | ¶ | FΒ | ¶ | Х | ¶ | X | ¶ | X | ¶ | X | 1 |
|----------------|-------|---|------|---|-------|---|-------|---|-------|---|-------|---|
| D07D06 | D05 | ; | -D04 | | -D03- | | -D02- | | -D01- | | -D00- | |

The SIO ready signal for channel A is represented here.

Diagnostic (DIAG) is reserved for starting the diagnostic routines. These routines can be initiated by power-up or reset.

Data Detected (DAD) is true when the data has been detected during a read from the tape cartridge at either low or high speed.

Feedback (FB) is the same as bit four in "Tape output status 2".

4.2 DATA COMMUNICATION WITH THE TAPE DRIVE

The preamble for tape data is 39 zeros (0) followed by a one (1) at the beginning of each data block. The postamble is the

reverse of this with a one (1) followed by 39 zeros (0) at the end of each data block. The preamble is stripped from the read data when data is being read in the forward direction. The postamble is stripped from the read data when the data is read in the reverse direction.

Example of a data block on the tape cartridge:

39 'ZEROS'¶ 1 'ONE'¶ ADDRESS ¶ DATA ¶ 1 'ONE'¶ 39 'ZEROS'¶

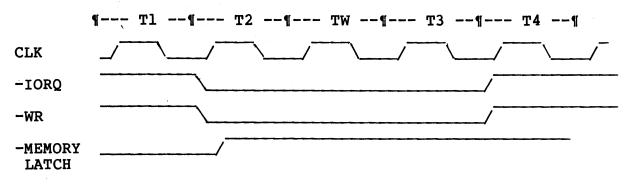
To create preambles and postambles, -DTR must be low on the SIO for the equivalent of a five-byte transmission time before and after the data is transmitted. During data transfer, -DTR is at a high level.

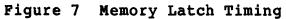
5.0 Input/Output Port Assignments and Timing

5.1 Input/Output Port Addresses

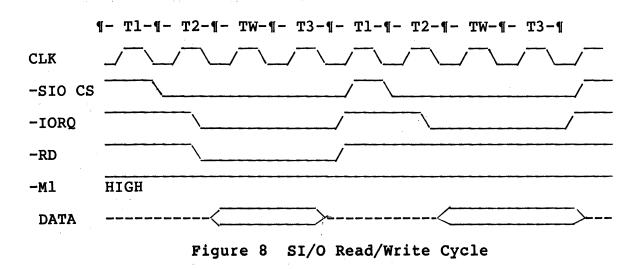
| . General Purpose Switch 00H> 03H |
|---|
| . CTC Channel 0 04H |
| . CTC Channel 1 05H |
| . CTC Channel 2 06H |
| . CTC Channel 3 07H |
| . Tape Input-Status 08H> 0BH |
| . Tape Input-Info 0CH> 0FH |
| . LED 10H> 13H |
| . Tape Select (Drive and Track) 10H> 13H |
| . Tape Command/Control 14H> 17H |
| . SIO Channel A (Data) 18H |
| . SIO Channel A (Command/Status) 19H |
| . SIO Channel B (Data) lAH |
| . SIO Channel B (Command/Status) 1BH |
| . Party Line Priority Controller 1CH> 1FH |

5.2 Memory Latch Timing

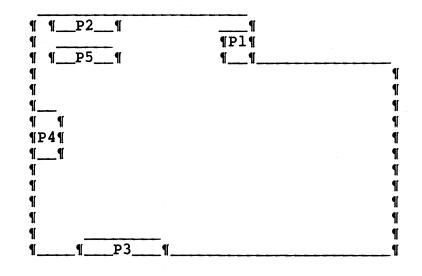




5.3 SI/O Read/Write Cycle



6.0 Connector Description





Legend

| Connector No. | Description | | | | |
|---------------|---|--|--|--|--|
| Pl | POWER CONNECTOR | | | | |
| P2 | RS422 CONNECTOR (TS806) | | | | |
| P5 | RS422 CONNECTOR (TS800) | | | | |
| P3 | TAPE INTERFACE CONNECTOR | | | | |
| P4 | RESET SWITCH CONNECTOR (OPTION) | | | | |
| | and the second secon | | | | |

6.1 Pl Power Connector (5-Pin Wafer)

| Pin No. | Description |
|---------|-------------|
| | |
| 1 | - 12 V |
| 2 | UNUSED |
| 3 | GROUND |
| 4 | + 5 V |
| 5 | + 12 V |

6.2 P2, P5 Communication Ports RS 422

(16 pin header on the PCB), (Rear panel 15 pin D-type conn.)

| D-conn. | : | Header | : | Description |
|---------|---|---------------|---|--------------|
| 1 | | 1 | | GND |
| 9 2 | | 2 3 | | -TxD +TxD |
| 10 3 | | 4 5 | | -RxD +RxD |
| 11 4 | | 6 7 | | +RTS -RTS |
| 12 5 | | 8 9 | | +CTS -CTS |
| 13 6 | | 10 11 | | -TxC +TxC |
| 14 7 | | 12 13 | | -RxC +RxC |
| 15 8 | | 14 15 | | Test GND |
| | | 16 | | NC |

6.3 P3 Tape Connector (50 Pin St. Header)

| Pin No. | Signal | 1/0 |
|---------|--------|-----|
| 2 | -SLD | I |
| 4 | -RDY | I |
| 6 | -WND | I |
| 8 | -FLG | I |
| 10 | -LPS | I |
| 12 | -FUP | |

| 14 | -BSY | I |
|----|-------|---|
| 16 | -EWS | I |
| 18 | -RWD | 0 |
| 20 | -REV | 0 |
| 22 | -FWD | Ō |
| 24 | -HSP | Ō |
| 26 | -WEN | Ō |
| 28 | -SL1 | Ō |
| 30 | -SL2 | Ō |
| 32 | -SL4 | 0 |
| 34 | -SLG | Ō |
| 36 | -RNZ | I |
| 38 | -RDS | I |
| 40 | -DAD | I |
| 42 | -WDE | 0 |
| 44 | -WN Z | 0 |
| 46 | -TR2 | 0 |
| 48 | -WDS | I |
| 50 | -TR1 | 0 |
| | | |

NOTE

All odd numbered pins are ground.

6.4 P4 Reset Connector (5-Pin Wafer) (Option)

| PIN | DESCRIPTION |
|------------------|--|
| 1 2 3 4 | NC (Option for lighted switch) Void (No pin) Reset Signal Ground NC (Not used) |

7.0 SPECIFICATIONS

- . Power Requirement ----- 1.08 Amp @ +5 V
- . Power Consumption ----- TYPICAL 5.4 Watts
- . System Clock ----- 4 MHZ
- . Memory
 - . 60 Kbytes of dynamic RAM
 - . 4 Kbytes of EPROM
- . Tape Communication
 - . SDLC Mode
 - . 192 Kbits/sec (supplied by tape drive)

. TS 806 Communication (RS 422)

- . SDLC / Asynchronous mode
- . Fix Rate of 800 Kbits/sec

Interrupt Priority

 1st Priority ----- SIO
 2nd Priority ----- CTC



SYSTEM REPAIR PRICE AND SPARE PARTS PRICE LISTS

This section contains the Repair Price List for Computers and the Systems Spare Parts Price List in effect at the printing date of this manual. Use these lists for estimating repairs: prices are subject to change without prior notice.



Repairs Price List for Computers

March 1, 1983

| DESCRIPTION | PRICE |
|---|--------|
| Logic Board TS 800 (Obsolete) | 135.00 |
| Logic Board TS 800A, 802, 802H | 150.00 |
| Logic Board TS 801 (Obsolete). | 175.00 |
| Logic Board TS 806, 806/20 | 250.00 |
| Logic Board TS 816, 816/40 | 350.00 |
| Logic Board TS 1602G, 1602GH | 400.00 |
| Graphics Board TS 1602G, 1602GH | 175.00 |
| Floppy Controller (Daughter Board) TS 802, 802H | 50.00 |
| Winchester Disk Controller (5" Drive & 40MB 8" Drive) | 175.00 |
| Tape Controller TS 806C | 95.00 |
| Interface Board TS 816U. | 50.00 |
| Keyboard TS 800, 800A, 802, 802H | 50.00 |
| Keyboard TS 1602G, 1602GH | 50.00 |
| Power Supply Module TS 800, 800A | 50.00 |
| Power Supply TS 802, 802H, 806, 806H, 816, 1602G | 110.00 |
| Video Module TS 800, 800A, 802, 802H, 1602G, 1602GH | 50.00 |
| Floppy Disk Drive 5" | 160.00 |
| Winchester Disk Drive 5" | 160.00 |
| Winchester Disk Drive 8" | 450.00 |
| Tape Drive | 300.00 |
| Picture Tube Broken P31 | 214.00 |
| Picture Tube Broken P39 | 230.00 |
| Top Case Broken TS 802, 800A, 1602G | 80.00 |
| Bottom Case Broken TS 802, 800A, 1602G | 100.00 |
| Top or Bottom Case Broken Computer Boxes TS 806, 816 | 80.00 |
| Front or Rear Panel Broken | 60.00 |
| Basic Repair charge (This additional amount charged when an entire system is returned for repair) | 70.00 |

TeleVideo will bill per above price schedule when no trouble is found in the module returned for repair.

Out of Warranty

Customer to return defective replaceable module freight prepaid to the factory, 1170 Morse Avenue, Sunnyvale, CA 94086. TeleVideo will send replacement repaired module, billing per above price schedule plus return freight.

Prices subject to change without notice.

TeleVideo Systems, Inc.

1170 Morse Avenue • Sunnyvale, CA 94086

Eastern Region – (212) 308-0705 • Northeast Region – (617) 369-9370 • Midwest Region – (312) 969-0112 South Central Region – (214) 258-6776 • Southwest Region – (714) 752-9488 • Northwest Region – (408) 745-7760 Southeast Region – (404) 447-1231 • European Sales – (31) 075-28-7461 TLX:844-19122 U.K./Scandinavian Sales – (44) 0908-668778 TLX:851-825151



Systems Spare Parts Price List

TeleVideo Systems, Inc.



TELEVIDEO COMPUTER SYSTEMS SPARE PARTS PRICE LIST

05-10-83

DESCRIPTION PART PRICE NUMBER MANUALS [class A] 2004200 20.00 Guide, Installation & User's TS 800A 20.00 Guide, Installation & User's TS 802 2003700 2003900 20.00 Guide, Installation & User's TS 802H 20.00 Guide, Installation & User's TS 803 2248100 20.00 Guide, Installation & User's TS 1602G 2133700 20.00 Guide, Installation & User's TS 1602GH 2133800 20.00 Guide, Installation & User's TS 1603 2248600 10.00 Guide, Installation & User's TS 806 2003000 20.00 Guide, Installation & User's TS 806/20 2226500 10.00 Guide, Installation & User's TS 806C 2004700 2002300 10.00 Guide, Installation & User's TS 806H 20.00 Guide, Installation & User's TS 806H/20 2232000 10.00 Guide, Installation & User's TS 816 2004100 20.00 Guide, Installation & User's TS 816/40 50.00 Guide, User's TELEPLAN 2226400 2150300 20.00 Manual, Operator's Tele3780 2200200 40.00 Manual, TeleDBMS 2219700 25.00 Manual, Mmmost 2003200 2003400 50.00 Manual, CP/M 50.00 Manual, CPM/86 2162400 2150400 40.00 Manual, TeleVideo - COBOL 50.00 Manual, Maintenance TS 800A, 802, 802H 2133900 50.00 Manual, Maintenance TS 806/20, 806, 806C & H 2230700 50.00 Manual, Maintenance TS 816 2131400 2259000 50.00 Manual, Maintenance TS 1602G/GH 50.00 Manual, Maintenance TS 803 2291000 KITS [class B] 2000700 300.12 Kit, Spare Parts, Logic Board 8 Bit Systems * 703.50 Kit, Spare Parts, Logic Board 16 Bit Systems * 2252100 2202900 244.80 Kit, Spare Parts, Logic Board WDC * 2203000 199.59 Kit, Spare Parts, Logic Board FDC * 401.40 Kit, Spare Parts, Graphics Board TS 1602G * 2252000 153.28 Kit, Spare Parts, Video Mod ts 800A * 2280700 2202800 90.88 Kit, Spare Parts, Mechanical TS 802/800A* 2228400 235.72 Kit, Spare Parts, Data Cables TS 816's* 223.08 Kit, Spare Parts, Data Cbls TS 806 & User Stat * 2270800

* Contents of Kits at end of List

05-10-83 DESCRIPTION PART PRICE NUMBER MAJOR ASSEMBLIES/PRINTED CIRCUIT BOARDS [class B] 2198900 1,474.14 PCB Asy Mod 8" Flpy TS 801 2018000 1,199.67 PCB Asy Logic Board TS 800A,802 2018001 1,199.67 PCB Asy Logic Board TS 802H 2226000 1,047.78 PCB Asy Logic Board TS 803,800 2020400 2,157.93 PCB Asy Logic Board TS 1602G 2022000 1,529.10 PCB Asy Logic Board TS 1603 2012000 1,251.12 PCB Asy Logic Board TS 806 TS 806/20 2017501 658.32 PCB Asy Logic Board TS 806C 2012500 1,765.83 PCB Asy Logic Board TS 816 2012501 1,403.05 PCB Asy Logic Board TS 816/40 360.00 PCB Asy Logic Board TS 816U 2013000 93.00 Video Module TS 800A,802 2195800 2226900 72.00 Video Module TS 803,1603 866.19 PCB Asy 5" Winchester Disk Controller (806, 806/20) 2013500 2013501 866.19 PCB Asy 8" Winchester Disk Controller (816/40) 210.75 PCB Asy Floppy Disk Cont TS 802 (Daughter Bd) 2017000 2019000 656.87 PCB Asy Graphics TS 1602G 654.00 Power Supply Switching 100W TS 806,806/20 2109200 827.34 Power Supply Switching 150W TS 802,802H,1602G 2109201 2129202 954.00 Power Suppy Switching 200W TS 816,806C 103.00 Power Supply 3A/5V 2191500 TS 800A (TS 800 OBS) 275.04 Power Supply 120W TS 803,1603 2227500 2227400 63.84 Power Supply Transformer TS 1603,803 288.00 Power Supply (OPC) TS 1602/SGH,802S 2304000 288.00 Power Supply (OPC) TS 803H,1603H 2304800 2294000 288.00 Power Supply (OPC) TS 806S/20 2299100 324.84 Power Supply (OPC) TS 816/40 STORAGE DEVICES [class C]

591.00 Floppy Drive 48 TPI D/S 5 1/4" 2099200 .5 MB Full Height 399.00 Floppy Drive 48 TPI D/S 5 1/4" 2221300 Half Height .5 MB 489.00 Floppy Drive 96 TPI D/S 5 1/4" 1.0 2252300 MB Half Height 5 1/4" 20 2198800 1,800.00 Disk Drive Winchester MB 2099400 4,140.00 Disk Drive Winchester 8" 23 MB 2220300 4,575.00 Disk Drive Winchester 115V 8" 40 MB 2220100 4,575.00 Disk Drive Winchester 230V 8" 40 MB 2099500 2,100.00 Tape Drive W/Codec Bd 14 MB

05-10-83

| 05-10-83 | | | | |
|--------------------|-------------------|---------------------------------------|---|---------|
| ****** | * * * * * * * * * | | ****** | * * * * |
| PART | PRI | CE DESCRIE | TION | |
| NUMBER | | | | |
| * * * * * * * * | | | *************************************** | * * * * |
| | M | ISC. MAJOR PARTS | [class D] | |
| 2040200 | 170 00 | | | |
| 2049300 | | Picture Tube Bla | | |
| 2173800 | | | ck/Green 12" P39 Graphics (1602G) | |
| 2218700 | | | ck/Green 14" P31P | |
| 2090200 | | | ard TS 800A/802 * | |
| 2183701 | | Detachable Keybo | | |
| 2183700 | | | ard TS 803/1603 * | |
| 2099800 | | Case Bottom | TS 806/806-20/RWP | |
| 2099900 | 70.20 | · · · · · · · · · · · · · · · · · · · | TS 806/806-20/RWP | |
| 2100600 | 110.00 | - | TS 802/1602G | |
| 2100700 2100800 | 85.00 | Case Bottom Bezel | TS 802/1602G | |
| 2141700 | | | TS 802/1602G | |
| 2141700 | 97.80 | | TS 800A | |
| 2141800 | | Case Top Bezel | TS 800A TS 800A | |
| 2103100 | | Case Top | TS 816/816-40 | |
| 2103100 | | Case Bottom | TS 816/816-40 | |
| 2188300 | | | Crt TS 803/1603 | |
| 2188500 | | | . TS 803/1603 | |
| 2188600 | | Case CRT | TS 803/1603 | |
| 2188800 | | Bezel | TS 803/1603 | |
| 2189100 | | Case Arm Top | TS 803/1603 | |
| 2189200 | | Case Arm Bottom | TS 803/1603 | |
| 2218800 | | Thumb Wheel Adj | TS 803/1603 | |
| 2291100 | | Power Cord | TS 803/1603 | |
| | | | 20 000/ 1000 | |

* ORDER APPROPRIATE LABELS FROM LABELS/LOGO'S (PAGE 7)

| 05-10-83 ******** | *************************************** | k * |
|----------------------|--|------------|
| PART NUMBER | PRICE DESCRIPTION | |
| ****** | *************************************** | < * |
| | FIRMWARE | |
| 8000002 | 16.92 IC 2332 Char Gen Upper Character Cell | |
| 8000003 | 16.92 IC 2332 Char Gen Lower Character Cell | |
| 800000 9 | 23.28 IC 8048 With Program ROM for Keyboard TS 800A/802 | |
| 8000016 | 15.0 IC charicter Gen. for Graphics Systems | |
| 8000018 | 34.50 IC 2716 System Program EPROM Kybd TS 800A/802 | |
| 8000045 | 37.50 IC System Program EPROM Z80 Portion TS 800A/802 | |
| 8000046 | 37.50 IC System Program EPROM Lower Terminal Firmware | |
| 8000047 | 37.50 IC System Program EPROM Higher Terminal Firmware 37.50 IC System Program EPROM 280 Portion TS 802H | |
| 8000050 8000024 | 55.80 IC System Program EPROM 450ns TS 806 | |
| 8000053 | 30.90 IC System Program EPROM 450ns TS 800 | |
| 8000106 | 30.90 IC System Program EPROM 450ns TS 816/40 | |
| 8000054 | 30.90 IC System Program EPROM 450ns TS 806C | |
| 8000093 | 24.00 IC System Program EPROM 450ns TS 1602G/GH | |
| 8000080 | 5.40 IC Memory Decode TS 1602G | |
| 8000 079 | 5.40 IC I/O Decode TS 1602G | |
| 8000027 | 55.80 IC Diagnostic EPROM TS 806 | |
| 8000052 | 55.80 IC Diagnostic EPROM TS 806C | |
| 8000096 | 55.80 IC Diagnostic EPROM TS 816 | |
| 8000107 | 55.80 IC Diagnostic EPROM TS 816/40 | |
| 8000035 8000036 | 23.70 IC L2-7 System Program ROM WDC (74S472 512 X 8) 23.70 IC MX-7 System Program ROM WDC (74S472 512 X 8) | |
| 8000037 | 23.70 IC FX-7 System Program ROM WDC (745472 512 X 8) 23.70 IC FX-7 System Program ROM WDC (745472 512 X 8) | |
| 8100024 | 100.00 Listing System Program TS 806 * | |
| 8100045 | 100.00 Listing 800A, 802, [Z80 Portion] * | |
| 8100046 | 500.00 Listing 800A, 802, 802H [6502 Portion] * | |
| 8100050 | 100.00 Listing TS 802H [Z80 Portion] * | |
| 8100053 | 100.00 Listing System Program EPROM TS 816 W/Code * | |
| 8100054 | 100.00 Listing System Program EPROM TS 806C * | |
| 8100093 | 100.00 Listing System Program EPROM TS 1602G/GH * | |
| ^ Require | non-disclosure agreements and letter of intended use. | |
| | CABLES/CONNECTORS/WIRE ASSEMBLIES | |
| 2005700 | 25.44 Cbl Asy, Keyboard TS 800A/802 | |
| 2006800 | 14.22 Cbl Asy, 20 Pin 12" TS 802H WDC To Winchester | |
| 2006900 | 18.60 Cbl Asy, 34 Pin 15" TS 802H Daughter Bd/Floppy WDC/Wind | ch |
| 2007001 | 32.28 Cbl Asy, 40 Pin 16" TS 802H Daughter Bd To WDC | |
| 2007300 | 29.16 Cbl Asy, 34 Pin 14" TS 802 Daughter Bd To Floppy | |
| 2006100 2006901 | 11.88 Cbl Asy, 40 Pin 2" TS 802/H Logic to Daughter Bd | . . |
| 2008901 | 18.96 Cbl Asy, 34 Pin TS 806 Logic To Floppy/WDC To Winch 20.28 Cbl Asy, 40 Pin 13" TS 806 Logic To WDC | 1 |
| 2006801 | 14.52 Cbl Asy, 20 Pin 15" TS 806-816/40 WDC To Winchester | |
| 2006501 | 11.34 Cbl Asy, 16 Pin 8" TS 806 RS 422 To Logic | |
| 2006600 | 36.00 Cbl Asy, 34 Pin 7" TS 806 Parallel Printer | |
| 2006400 | 24.12 Cbl Asy, 50 Pin 16" TS 816 WDC To Winchester | |
| | | |

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| 05-10-83 | و باو باو باو باو باو باو | 1 | مالد مالد عالم عالم عالم عالم عالم و | و علو علو علو علو علو علو | **** | |
|--------------------|---------------------------|-----------|--------------------------------------|---------------------------|---|-------------------|
| | | | | | | ~ ~ ~ ~ ~ ~ ~ ~ ~ |
| PART | PRIC | نظن | DESCRIP | TION | | • |
| NUMBER | ****** | | ****** | ****** | * | **** |
| | | | · | | 816U To Logic, Logic To T | |
| 2006201 2006601 | 20.10 | Cbl Asy, | | | 816 Parallel Printer | ape |
| | 30.34 | CDI ASY, | 34 P10 | | | |
| 2006500 | 32.70 | Cbl Asy, | 10 P1n | 3° TS | 806C RS 422 Internal | Dring |
| 2007100 | | Cbl Asy, | | | 806C Logic Board To Tape | |
| 2007101 | 42.00 | Cbl Asy, | 20 Pin | | 806 Internal For TS 806H | |
| 2007800 | /3.08 | Cbl Asy, | 34 P1n | 10" TS | 806 Internal For TS 806H | |
| 2128500 | 19.32 | CDI ASY, | 20 Pin | | 806H External to TS 806 | |
| 2006300 | 23.40 | CDI ASY, | 34 P1n | | 806H External to TS 806 | Law |
| 2007600 | 42.60 | CDI ASY, | 20 Pin | IZ" TS | 806H Internal for Winches | ter |
| 2007700 | | | | | 806H Internal for Winches | ter |
| 2135700 | | Cbl Asy, | | | 1602G Logic to Graphics | |
| 2007002 | | | | | 816/40 Logic to WDC | 01 617 |
| 2006204 | | | | | 816/40 Logic to Tape/Logic | C to 8160 |
| 2006200 | | | | | 816/40 Logic to WDC | |
| 2006404 | | | | | 816/40 WDC to Winchester | |
| 2224300 | | | | | 816/40 Parallel Printer | 1 6 / 4 0 |
| 2235400 | | Cbl Asy, | | | 816/40H External to TS 8 | |
| | | | | | | 16/40 |
| 2235500 | | | | | 816/40 Internal for TS 8 | |
| | 150.06 | CDI Asy, | 50/5/ P | in TS | 816/40 Internal for TS 8 | 16/4UH |
| 2161000 | | | | | 816/40H Internal for Winc | |
| 2235600 | | | | | 816/40H Internal for Winc | hester |
| 2008600 | | Harness A | | | | |
| 2007900 | | | | | Le 16" TS 806 | |
| 2136500 | | | | | Le 12" TS 806 | |
| 2008000 | | | | | Le 24" TS 816 | |
| 2008101 | | | | | ette Power-1 | |
| 2008201 | | | | | ette Power-2 | |
| 2176200 | | Harness A | | | | |
| 2008800 | | | | | Serv. Board TS 816 | |
| 2008900 | | | | | ette Power 1 & 2 | |
| 2136702 | | Harness A | | | | |
| 2192901 | | Harness A | | | Coster D Div 140 | |
| 2008400 | | | | | : Switch 3 Pin 14" | |
| 2008401 | | | | | : Switch 3 Pin 15" | State Anna State |
| 2008700 2097900 | | | | | ppy General 14" | • |
| 2141200 | | | | | PCB Mount (AMP) | |
| 2208400 | | | | | PCB Mount 6 Pin | |
| 2098000 | | RJ12 Modu | | | Banala DCD Namet DC400 | |
| 2163100 | | | | | Female PCB Mount RS422 | 400 |
| | | | | | Metal Female PCB Mount RS | 422 |
| 2097800 | | | | | Female PCB Mount RS232 | 222 |
| 2165300 2098100 | | | | | Metal Female PCB Mount RS | L3L |
| | | | | | - Straight | |
| 2216400 | | | | | - Straight | |
| 2098103 | | | | | - Straight | |
| 2098104 | | | | | - Straight | |
| 2098106 | | | | | - Straight | |
| 2098107 | | | | | - Straight | |
| 2098108 | 1.50 | connector | 50 Pin | неаде | r - Straight | |
| , | | | | | | |

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| *************************************** | ***** |
|---|-------|
| PART PRICE DESCRIPTION NUMBER | |
| *************************************** | ***** |
| 2174401 8.10 Connector 50 Pin Header - Angle | |
| 2098300 1.80 Connector 2 Position Jumper | |
| 2098703 1.74 Connector 2 Pin Right Angle Molex | |
| 2098800 .72 Connector 2 Pin Straight Wafer | |
| 2098801 3.72 Connector 3 Pin Straight Wafer | |
| 2098802 .72 Connector 5 Pin Straight Wafer | |
| 2098700 1.02 Plug 5 Pin Molex Right Angle Wafer | |
| 2001200 11.34 Jack Socket Connector Kit | |
| 2109000 19.87 Power Cord 3 Conductor 3 Prong 6 FT | |

CRYSTALS

| 2098602 | 6.60 | Crystal | 1.8432 | MHz | | | | |
|---------|-------|----------|-----------|-----------|-----------|-------|------|------|
| 2216500 | 18.00 | Crystal | 4 | MHz | | | | |
| 2098603 | 4.80 | Crystal | 8.0000 | MHz | | | | |
| 2098605 | 3.54 | Crystal | 13.6080 | MHz | | | | |
| 2048800 | 18.00 | Crystal | 15 | MHz | | | | |
| 2042800 | 27.00 | Crystal | 16 | MHz | | (MOT, | CTS, | HYT) |
| 2098604 | 4.50 | Crystal | 20.000 | MHz | | | | |
| 2035200 | 37.08 | Crystal | 23.814 | MHz Klll | 4A | (MOT, | CTS) | |
| 2099700 | 1.02 | Insulato | or Mounti | ing Pad F | or Crysta | 1 | | |
| | | | | | | | | |

FANS

| 2099000 | 35.28 E | an l | .15V/23 | OV AC | 36-4 | 47 CFM | (AIR | OVER) |
|---------|---------|-------|---------|-------|------|---------|------|-------|
| 2141500 | 52.50 E | Fan B | 30x 23 | OV AC | TS 8 | 302/802 | H | |
| 2245800 | 56.28 E | Fan B | Box 11 | 5V AC | TS] | L602G/G | H | |
| 2245700 | 67.20 E | Fan B | 30x 23 | OV AC | TS I | L602G/0 | H | |

BEZEL'S/CASE ASSEMBLIES

| 2142000 | 40.00 | Panel | Front | TS | 806 |
|---------|-------|---------|-----------------------------------|----|----------------|
| 2142100 | 40.00 | Panel | Front | TS | 806C |
| 2142200 | 40.00 | Panel | Front | TS | 806H |
| 2105401 | 40.00 | Panel | Back | TS | 806 |
| 2105500 | 40.00 | Panel | Back | TS | 816 |
| 2105600 | 40.00 | Panel | Front | TS | 816 |
| 2100102 | 10.00 | Shroud | Connector | TS | 800A |
| 2100500 | 10.00 | Shroud | Connector | TS | 802/1602G |
| 2219800 | 5.00 | Card Gu | uide | TS | 802 |
| 2100300 | 15.00 | Cover H | Fan | TS | 816 |
| 2105700 | 28.80 | Panel H | Floppy Cover | TS | 802H (Plastic) |
| | | | 그는 그는 것은 것은 집안에서 가지 않는 것이 많이 있었다. | | |

MISCELLANEOUS SPARES

| 2101600 | 64.80 | Chassis Mounting | ТS | 806/806H |
|---------|-------|--------------------------|------|-----------|
| 2101900 | 5.88 | Panel Cover Hard Disk | TS | 806 |
| 2103500 | 19.98 | Cover Top Power Supply | TS | 802 |
| 2152800 | 6.72 | Speaker (8 ohm) With Con | nnec | tor |
| 2150100 | 41.70 | Filter A.C. Line SAE HP: | 2-2 | |
| 2001300 | 2.88 | Bail Mount Enclosure | | |
| 2101900 | 5.88 | Panel Shield Winchester | Dis | sk TS 806 |
| 2204900 | 2.89 | Shield Board TS 816 | | |
| 2225200 | 31.80 | Bracket TS 816/40 WDC | | |

LABELS/LOGO'S

| 2105000 | 1.50 Label Logo Plastic "TeleVideo" Systems |
|-----------------|---|
| 2191300 | 2.40 Label Logo Plastic "TeleVideo" Printer |
| 2105104 | .72 Label Keyboard TS 800A |
| 2105105 | .72 Label Keyboard TS 802 |
| 210510 6 | .72 Label Keyboard TS 802H |
| 2208500 | .72 Label Keyboard TS 1602G |
| 2208600 | .72 Label Keyboard TS 1602GH |
| 2105300 | .96 Label Back Panel "RS 232" TS 806-806/20 |
| 2105301 | .96 Label Back Panel "Terminal" TS 806-806/20 |
| 2154402 | 1.50 Label Back of Unit TS 800A |
| 2142900 | .90 Label S1 TS 800A |
| 2142901 | .90 Label S2 TS 800A |

BOXES/PACKING MATERIAL

| 2208800 | 25.00 | Carton | Inner | Shipping | TS | 802/TS | 1602G |
|---------|-------|--------|--------|----------|-----|--------|-------|
| 2143600 | 16.86 | Carton | Outer | Shipping | TS | 802/TS | 1602G |
| 2143700 | 20.52 | Carton | Inner | Shipping | TS | 806 | |
| 2208900 | 31.98 | Carton | Outer | Shipping | TS | 806 | |
| 2143800 | 26.22 | Carton | Inner | Shipping | TS | 816 | |
| 2209000 | 31.98 | Carton | Outer | Shipping | TS | 816 | |
| 2185700 | 3.18 | Corner | Blocks | Shipping | Car | ton | |
| 2208200 | 16.00 | Carton | | Shipping | Haı | d Disk | |
| 2214401 | 12.00 | Formed | Foam | TS 806 | | | |
| 2237300 | 10.00 | Formed | Foam | TS 800A | | | |
| | | | | | | | |

ALL PRICES SUBJECT TO CHANGE WITHOUT NOTICE

MINIMUM ORDER \$500.00

05-10-83 PARTS KITS PRICE SPARE Kit, Spare Parts, Power - Video Module TS 800A \$ 153.28 2280700 2042200 2N3906 Vertical Amplifier 2046500 2N3904 Vertical Drive 2N4401 Horizontal Drive 2200000 2N5551 Reference Amplifier 2047100 2046700 KTC1627A 75Volt Regulator 2201400 DS135D/IN391 2047500 IN914 2200800 Yoke Deflection With Connector 2201200 Transformer Horizontal Drive Linearity Coil 5.40uh 2200900 2201300 Transformer Flyback (High Voltage) 2201000 Inductor 27uh 2126900 Voltage Regulator LAS 16CB 2A/13.8Volts Resistor CF 390 Ohms 1/2Watt 5% 2186200 2201600 IN759A Zener Diode 2199300 Capacitor 220uf 16Volt Electrolytic Capacitor .luf 600Volt Mylar 2197300 2047300 2SC2233 MJE13006 2280800 Diode 30S2 I.R. 3 Amp 2126600 Voltage Regulator LAS L1405 3 Amp 5 Volt 2000700 Kit, Spare Parts, Logic Bd 8 Bit Systems \$ 300.12 2042600 26LS32 RS422 interface 2042400 26LS31 RS422 interface 2050600 SIO/2 serial communications chip 2050800 CTC counter timer chip 2051000 CPU Z80A central processor unit DMA direct memory access chip 2051200 2051600 64K dynamic RAM (4 each) 2029200 75188N 2029400 75189N 2228400 Kit Data Cables TS 816's \$ 235.72 2006201 Logic to Tape/Logic to 816U 2006400 Logic to Winchester 816 2006801 816/40 WDC To Winchester 2006601 816 Parallel Printer to Logic 2007002 816/40 logic to WDC 816/40 Logic to Tape/Logic to 816U 2006204 2006200 816/40 Logic to WDC 2006404 816/40 WDC to Winchester 2224300 816/40 Parallel Printer 2203000 Kit, Spare Parts Logic bd FDC \$ 199.59 2040200 WD2143-01 four phase clock logic 2040400 WD1691 floppy support logic 2040600 FD1793-02 (93816 fair amd) floppy controller

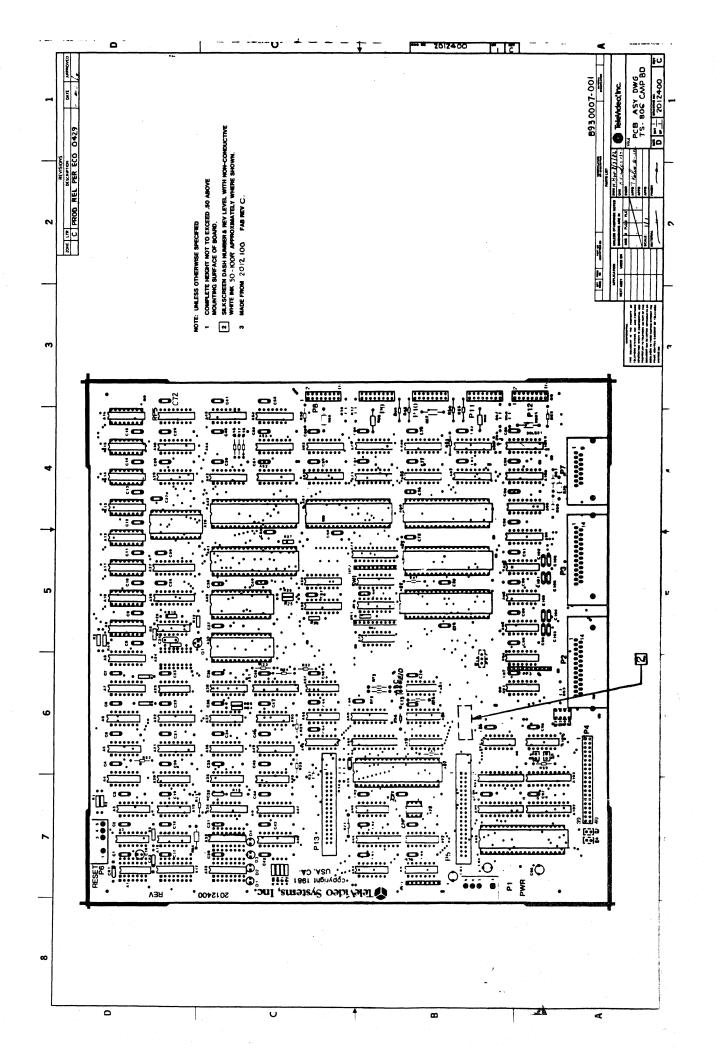
05-10-83 PARTS KITS PRICE SPARE Kit Data Cables User Stations and TS 806's \$ 223.08 2270800 2006800 802H WDC To Winchester 802H Daughter Board To WDC 2007001 2007300 802 Daughter Board To Floppy 802 Logic to Daughter Board 2006100 1602G Logic to Graphics Board 2135700 802H Daughter Board To Floppy/ WDC To Winchester 2006900 806 RS422 to Logic 2006501 806 Logic To Floppy/ WDC To Winchester 2006901 2007000 806 Logic To WDC 806 WDC To Winchester 2006801 2006600 806 Parallel Printer \$ 90.88 Kit, Spare Parts, Mechanical TS802/800A 2202800 cord for keyboard 2005700 3 amp 125V fuse (25 each) 2223700 2199400 keyswitch 2096800 10 position side dip switch 2223300 1 amp 250V fuse (25 each) 2182100 RS232 connector RJ-11 connector 2097900 2098000 RS422 connector 2100500 connector shroud \$ 244.80 2202900 Kit, Spare Parts Logic bd WDC 2056200 parallel converter 2056400 MFM converter 2056600 AM detector 2056800 CRC generator/checker parallel to serial converter 2057000 2057200 delay line 2252000 Kit, Graphics Board \$ 401.40 2057400 IC Gate Array Graphics/1603 IC 7220 Graphic Display Controller IC Dynamic RAM 4116 16K x 1 (120ns) [4 each] 2139800 2139200 2252100 Kit, Spare Parts, Logic Board 16 Bit Systems \$ 703.50 2051600 64K dynamic RAM (4 each) 2029200 75188N 2029400 75189N 2042600 26LS32 RS422 interface 26LS31 RS422 interface 2042400 2054000 IC 8284A Clock Generator 2054200 IC 8288 Bus Controller IC P8088 CPU IAPX 88/10 2054400 IC 8274 USART 2054600 2054800 IC 8254 Program Interval Timer 2055000 IC 8259A Priorty Interrupt Controller IC 8089 IOP 2055200

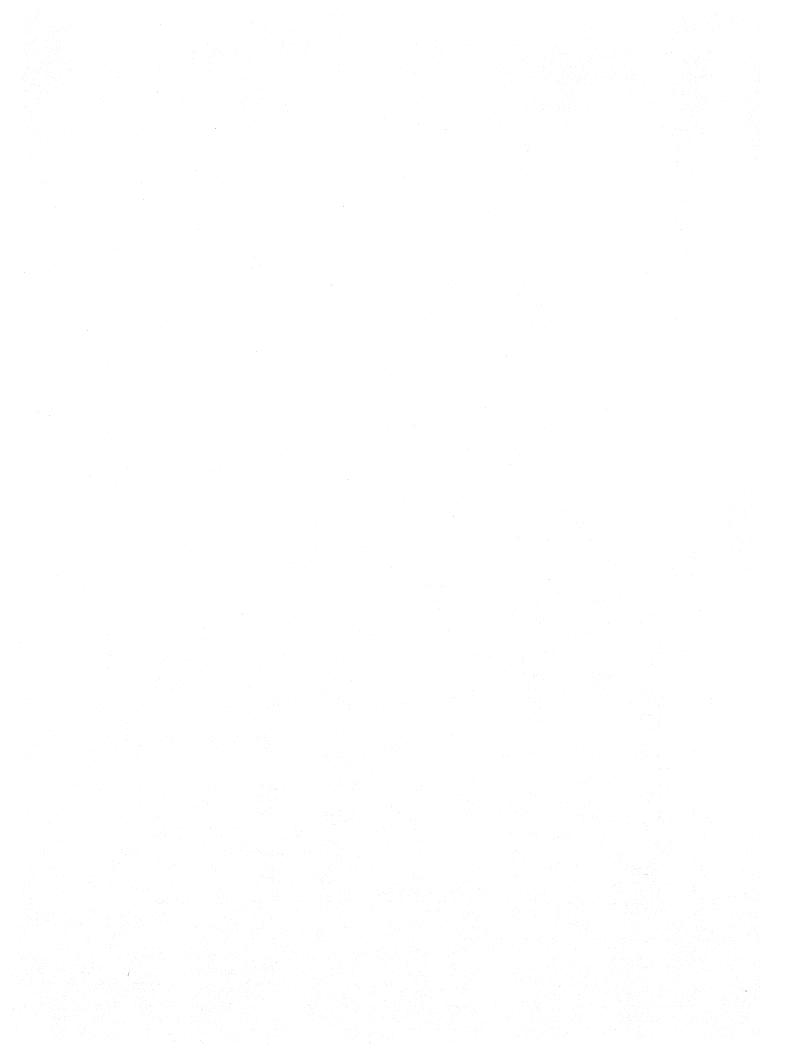


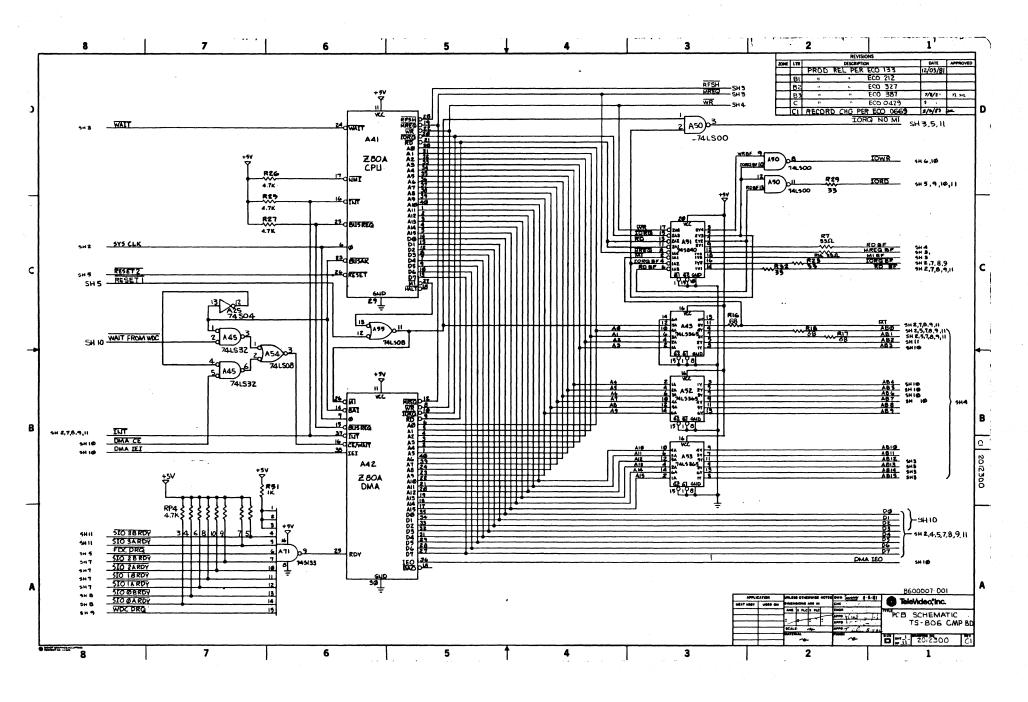
DRAWINGS

This section contains the latest board assembly drawings and logic diagrams. When ordering parts, use the component type or value shown in the diagrams to refer to the TeleVideo part number listed in the Spare Parts Price List. These prices are for estimating only.

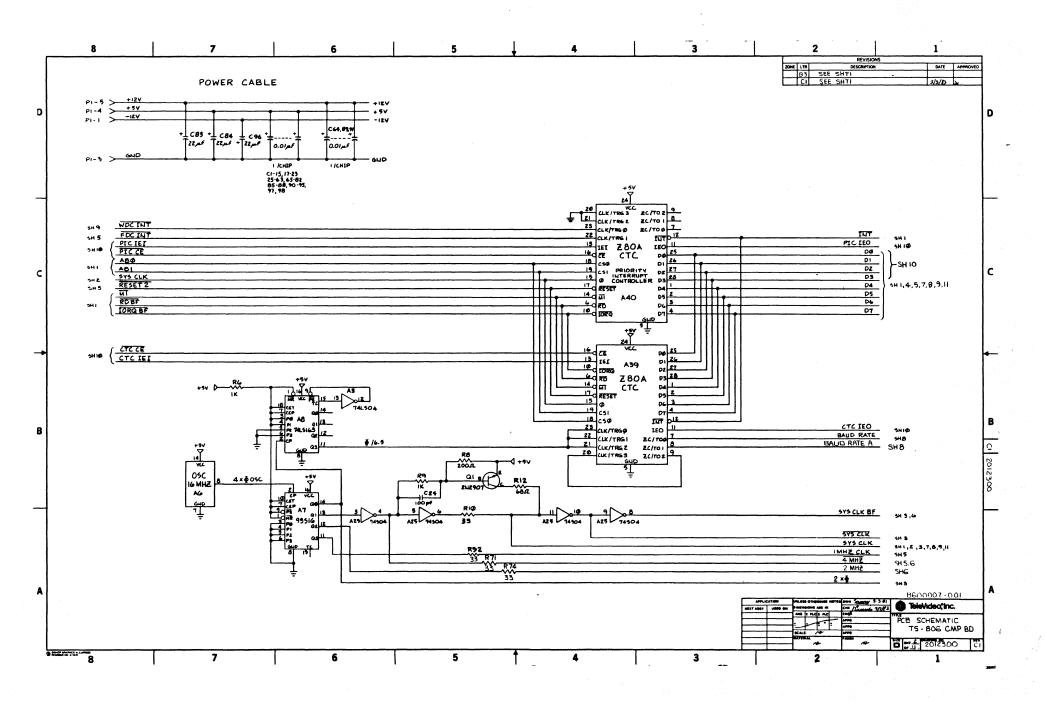




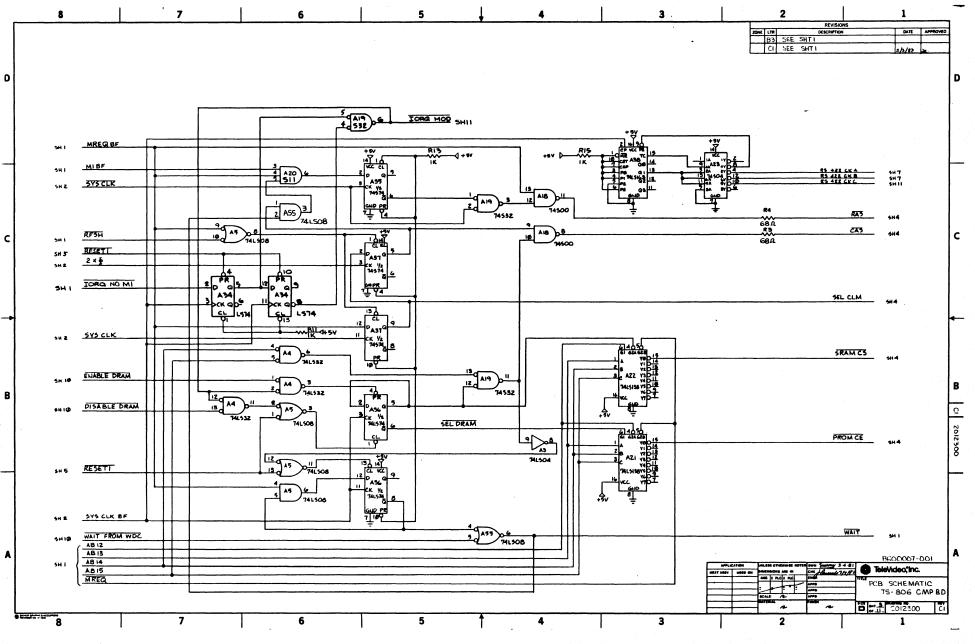


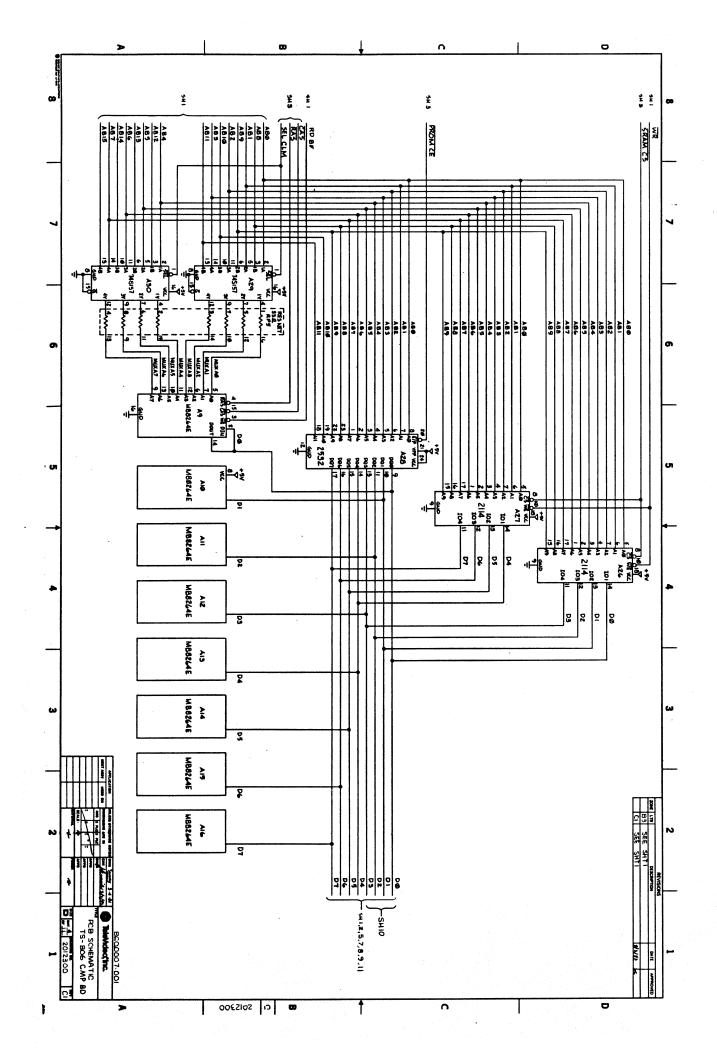




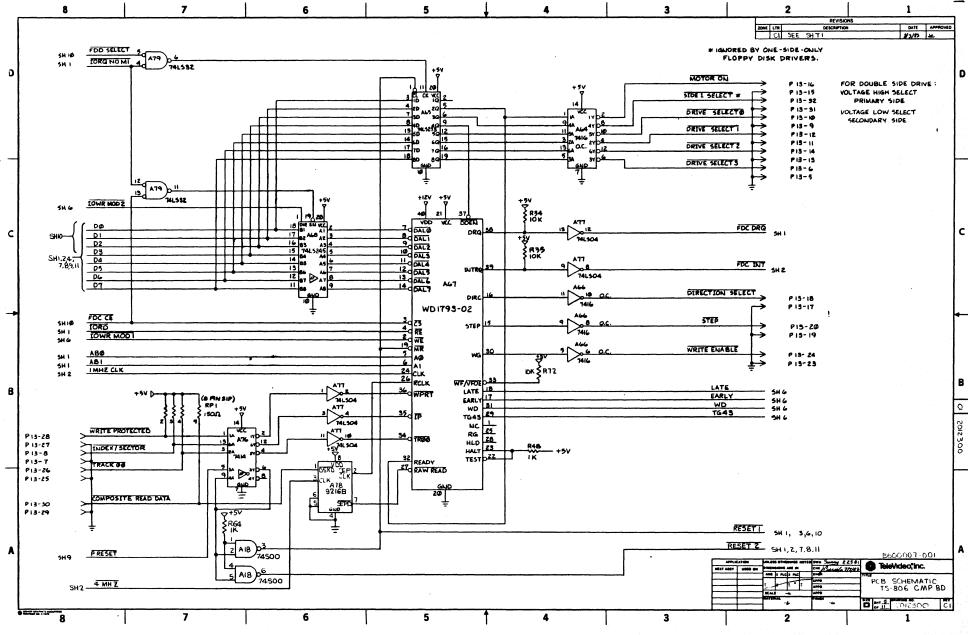




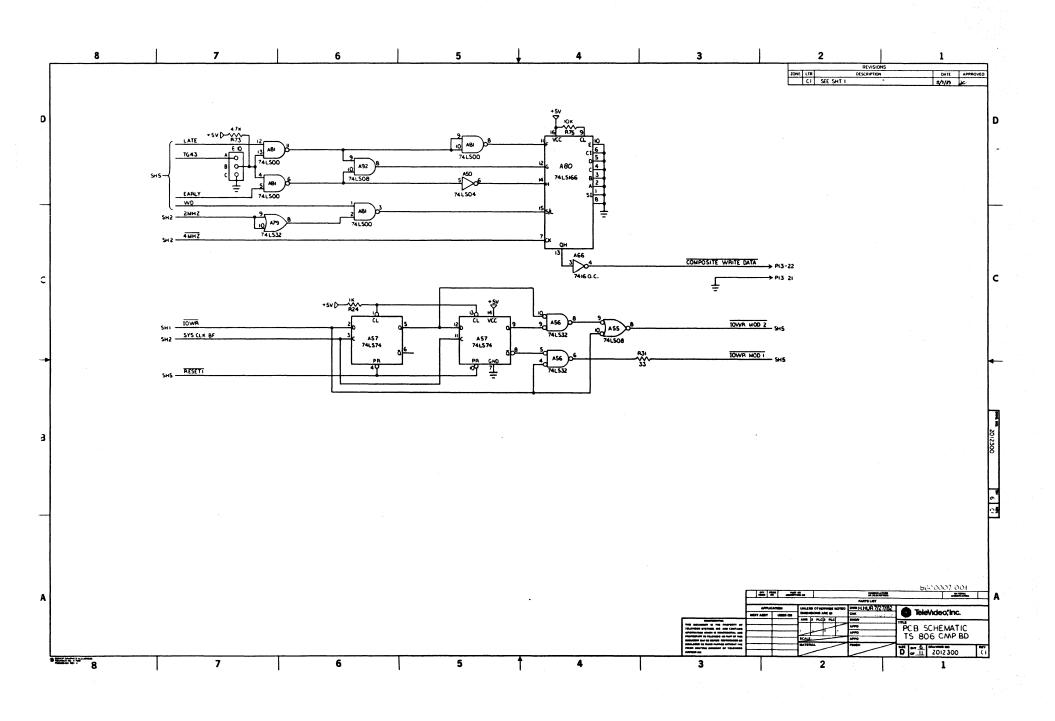




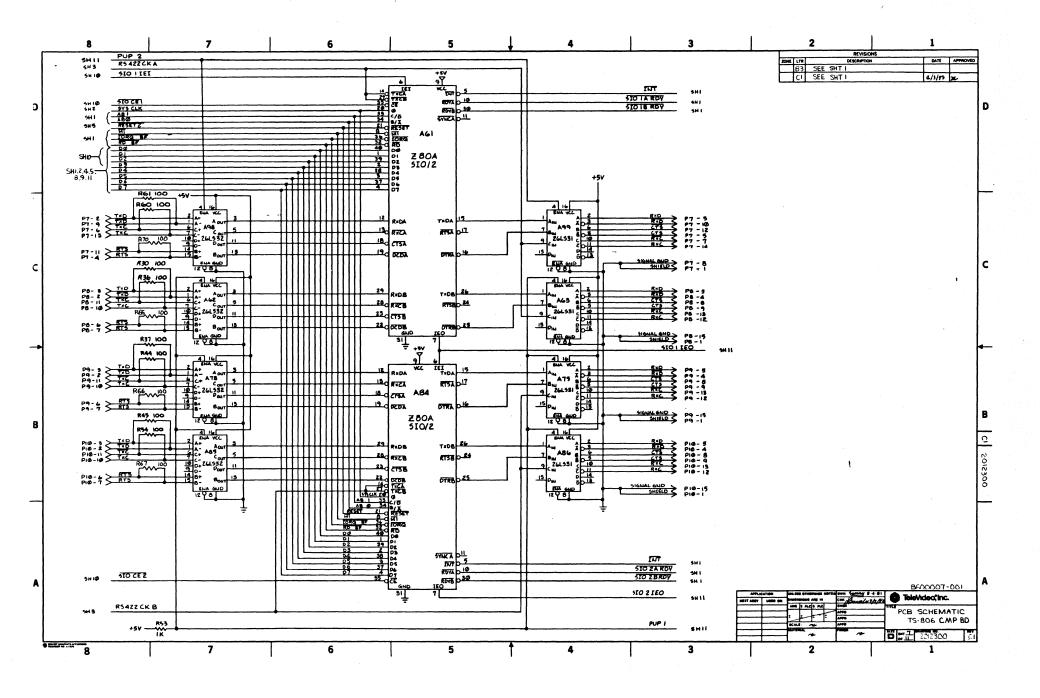






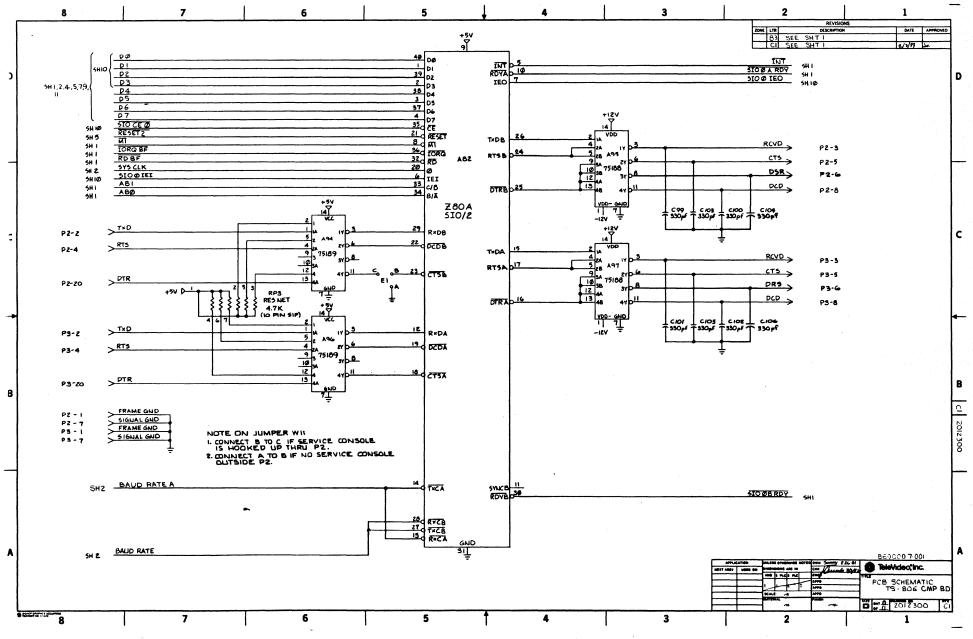




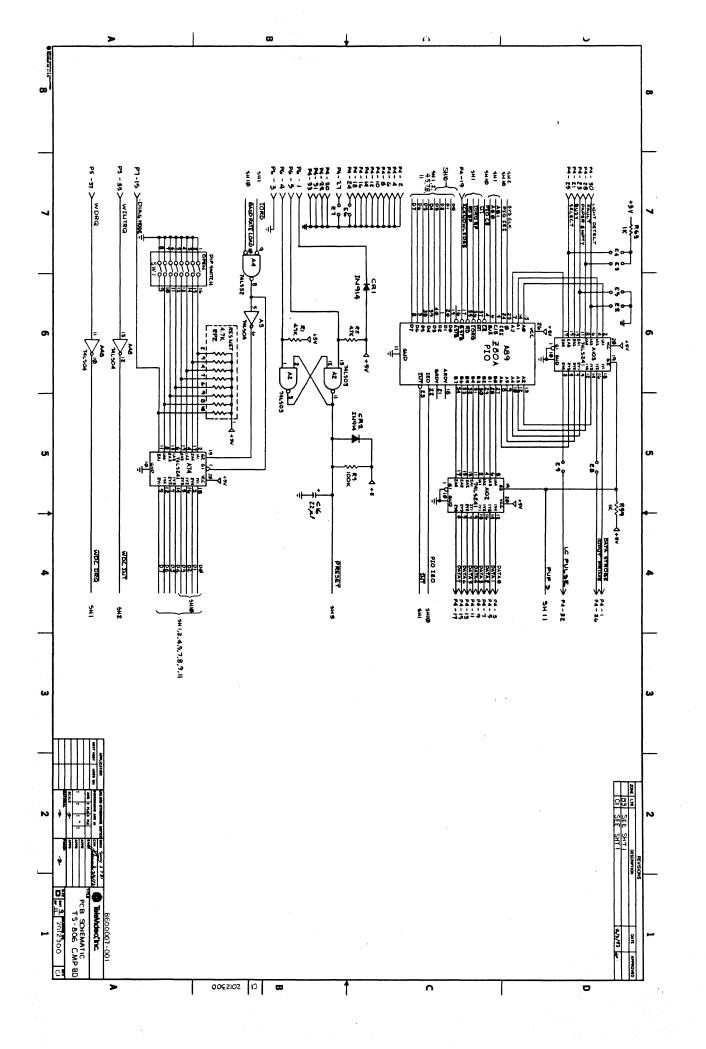


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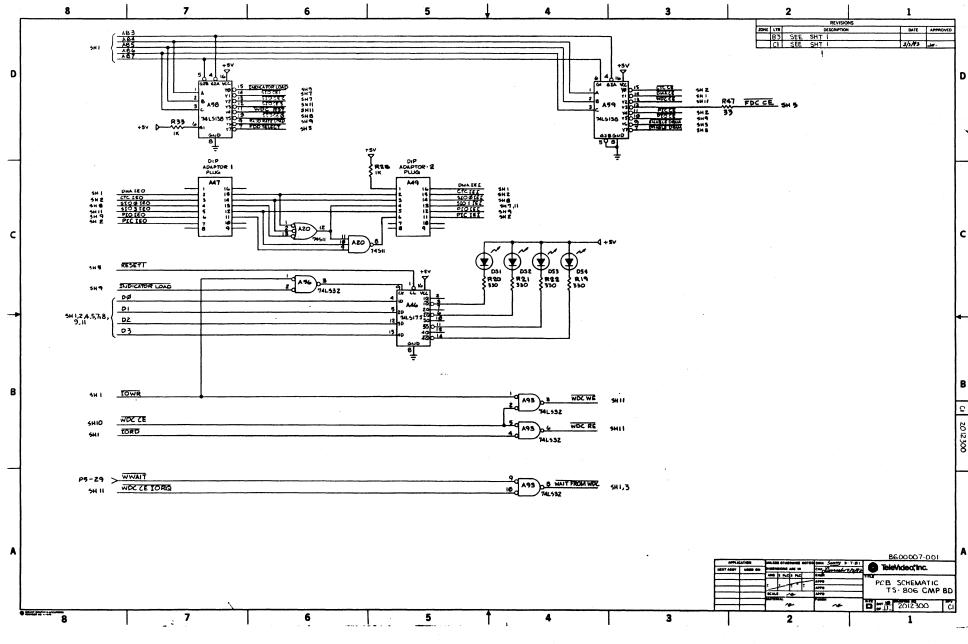




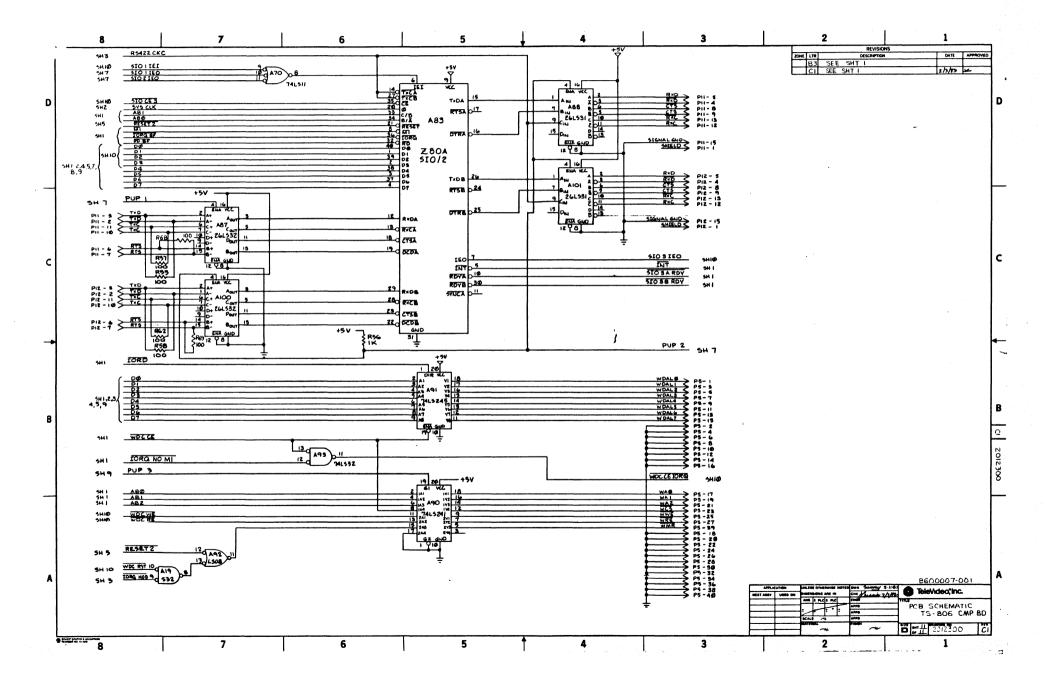




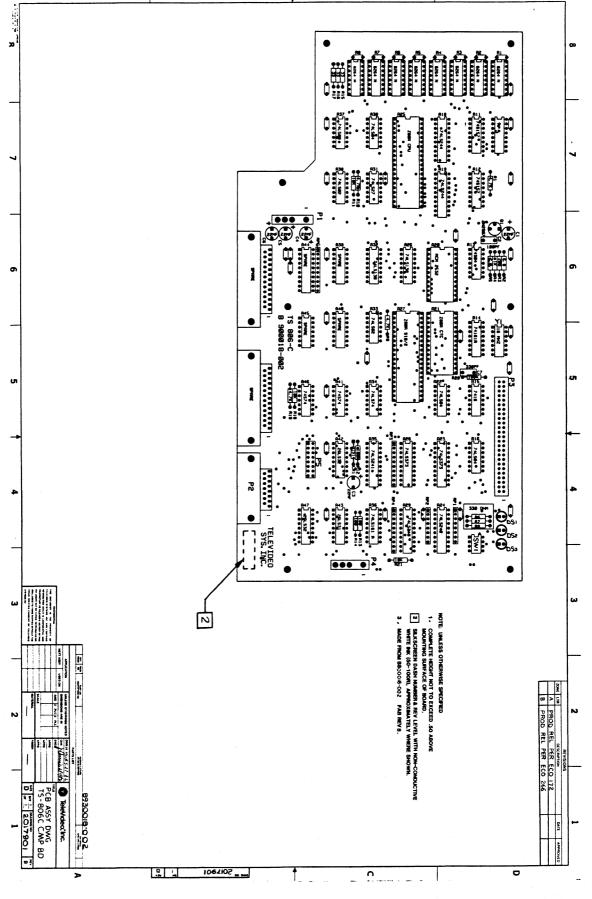




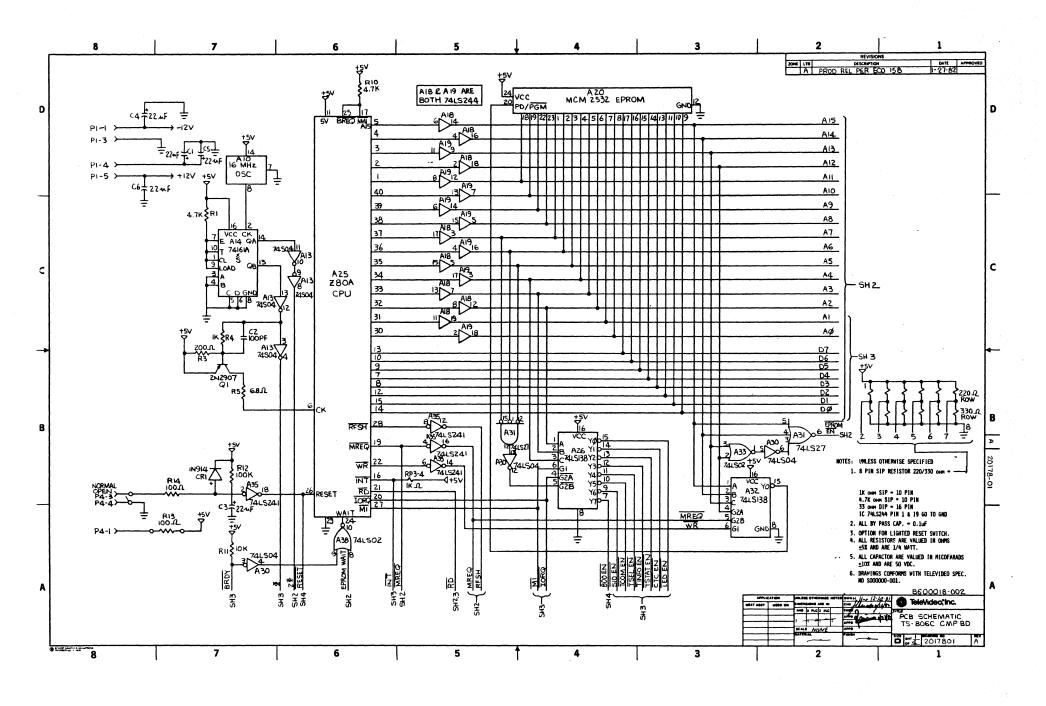


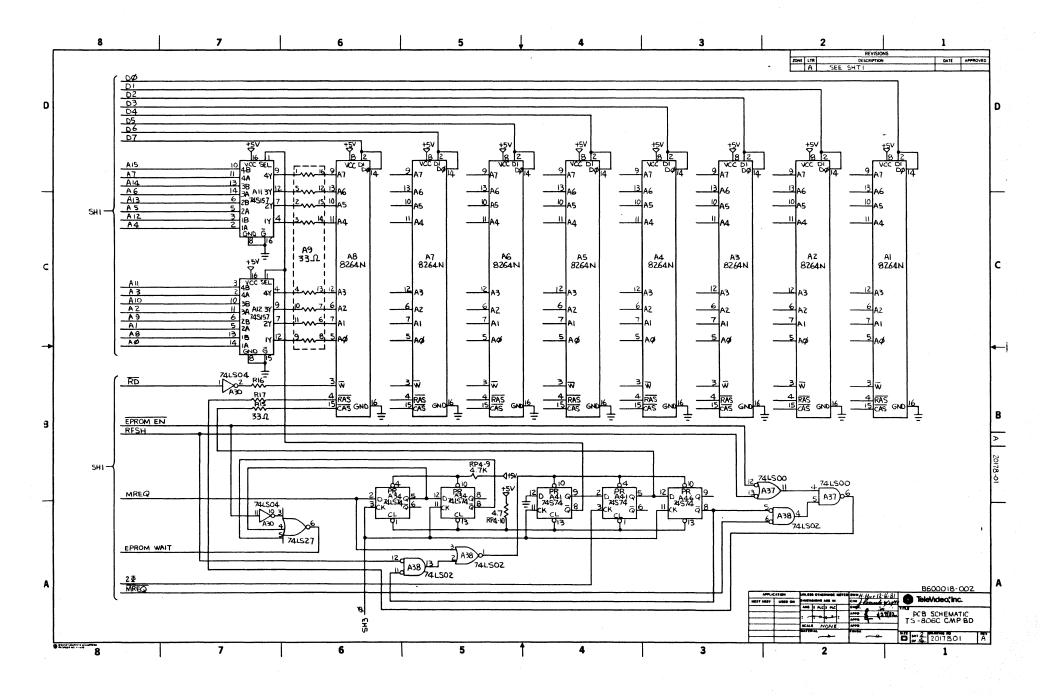




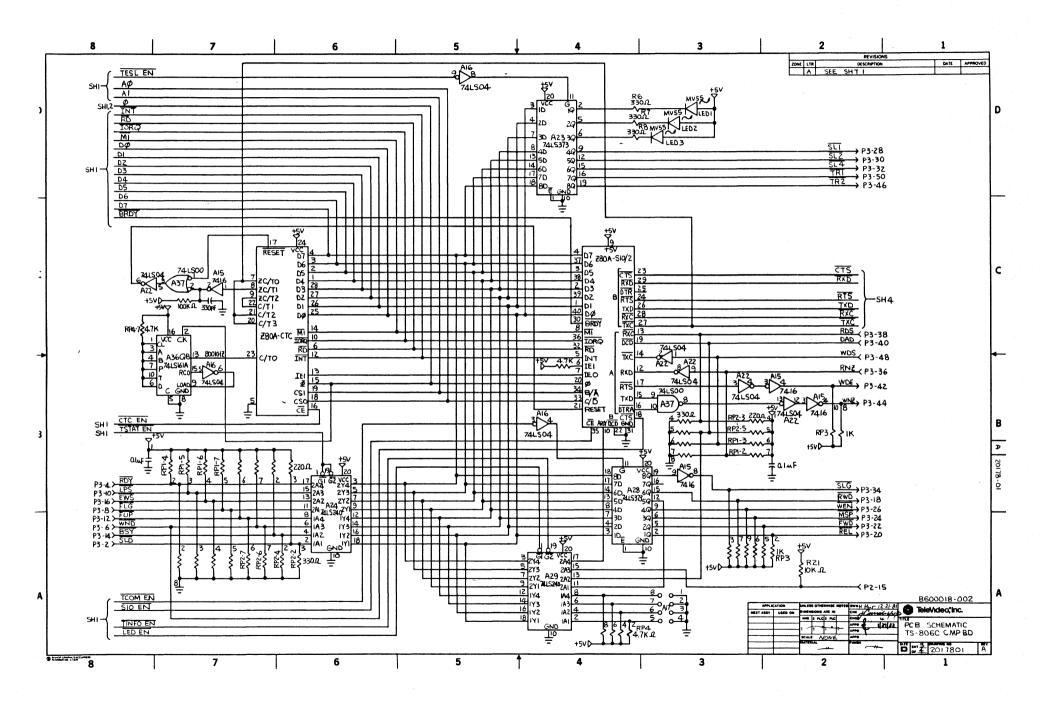




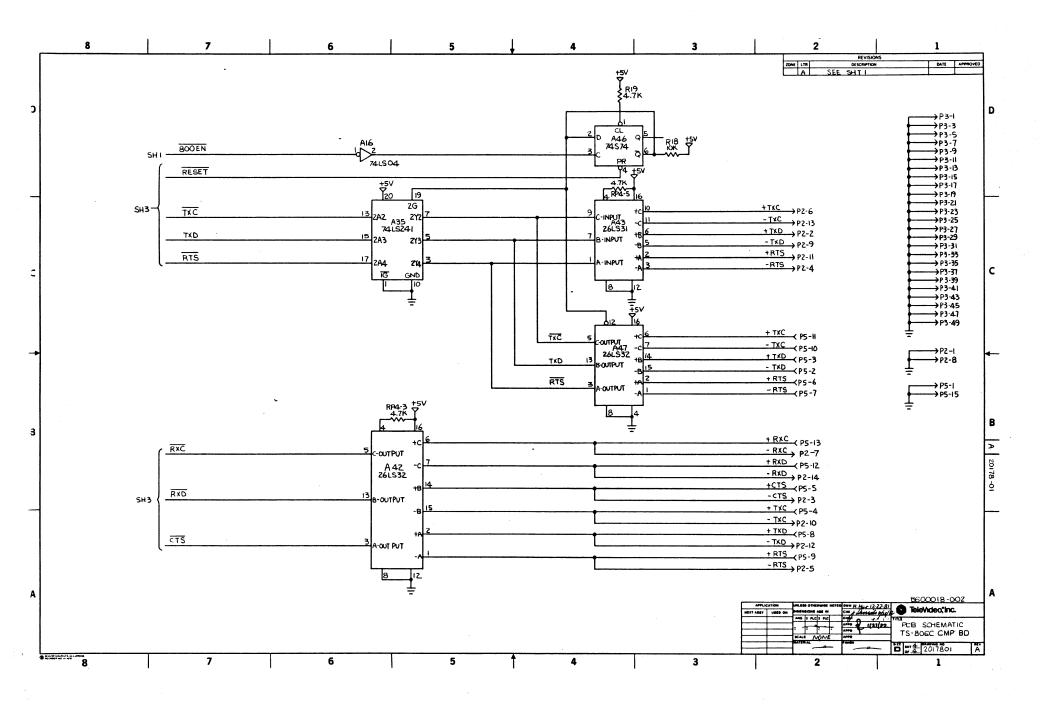




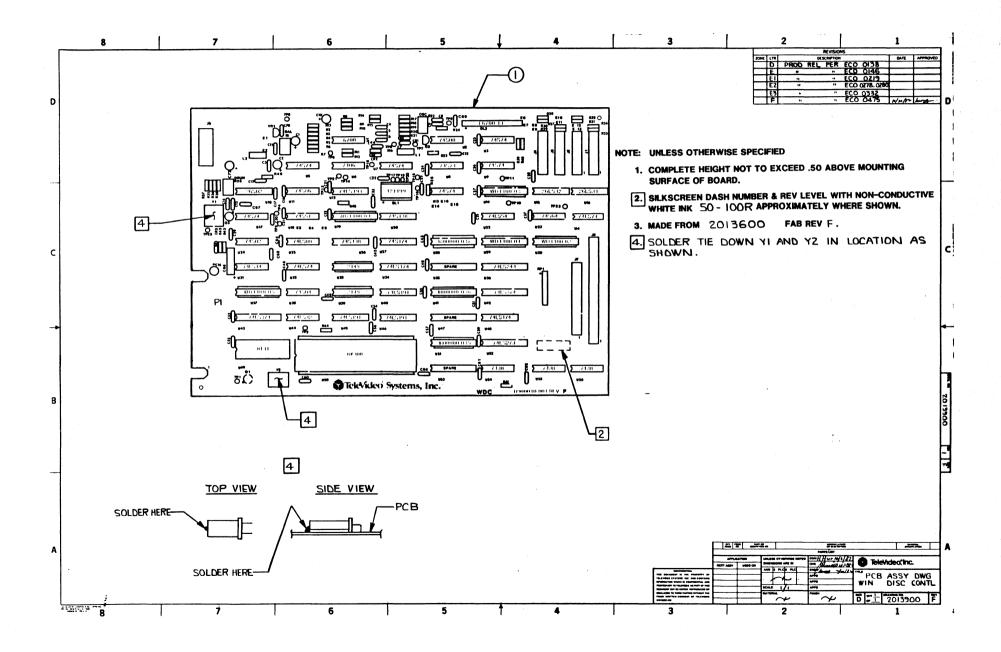








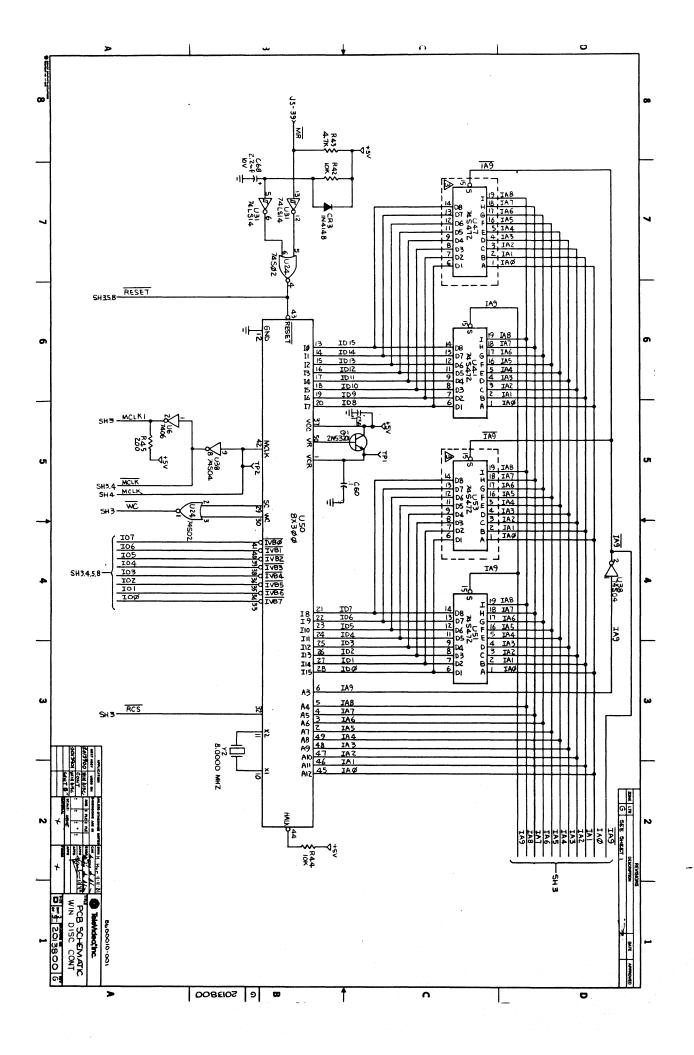


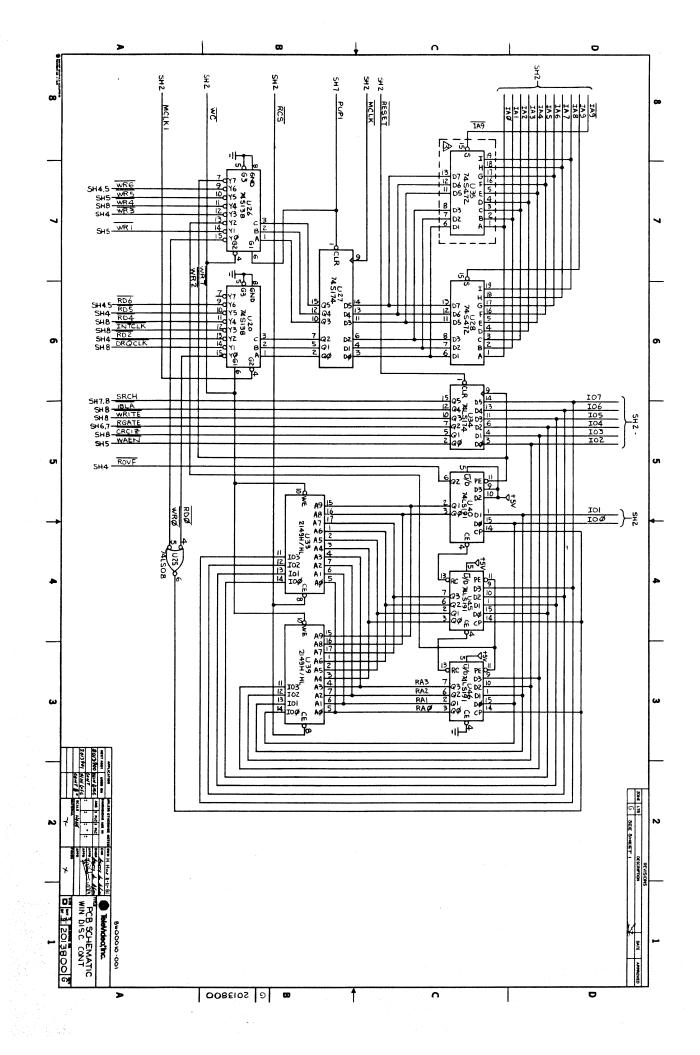


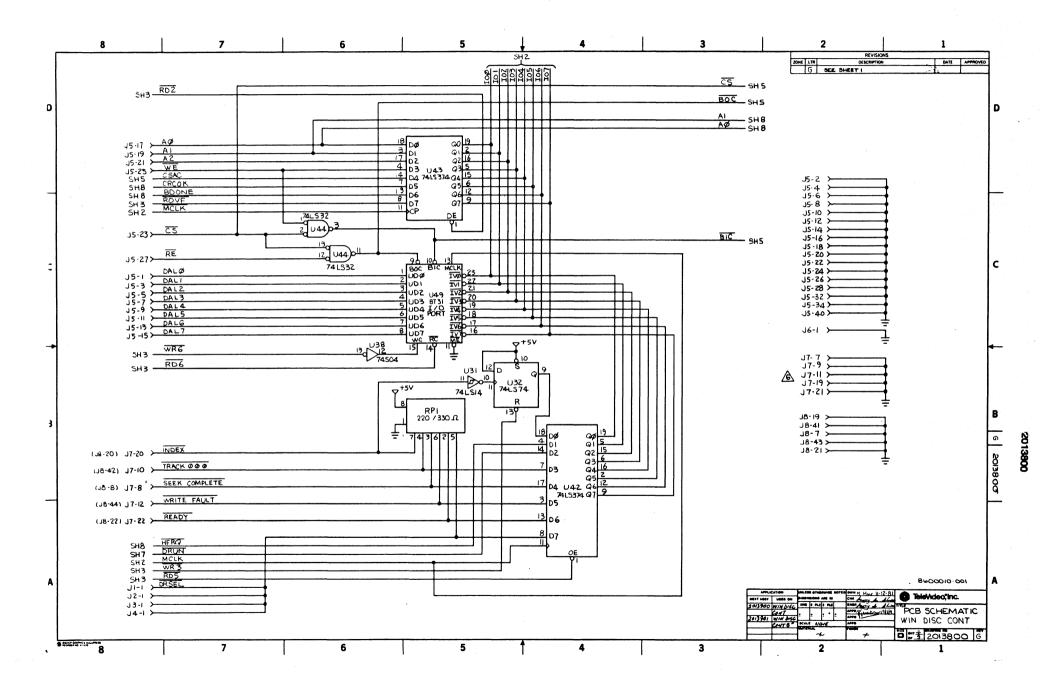


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|---|-----------------|--|--------------------|---|--|---|--------------------|---------------------------|--------------------------------|--------------------------|
| | • | | | | | | | ZONE | | DATE |
| | | | | | | | | A | PROD REL PER ECO | 125 4/30/82 |
| | | | | | | | •. • | | | |
| | MNEMONIC | MEANING | | | | MNEMONIC | 1. N | MEANING E | ECO (| 0146 |
| | | | | | | | | | ECO C | 0322 0415 ⊂ੇ੩ יי/א/אי |
| | AMDET | ADDRESS MARK DETECT | | | · · · · · | | | G | ECO (| 0594 5 1111/11 |
| | A2-AØ | TASK FILE ADDRESS SELECT | | | -SEEK COMPLETE SRCH | SEEK COMPLETE SEARCH | STATUS FROM DRIVE | | | |
| | BIC | BUS INPUT CONTROL | | | -STEP PULSE | STEP PULSE | TO DRIVE | | | |
| | BOC | BUS OUTPUT CONTROL | | | TIMCLK | TIMING CLOCK FOR SAIGOO TRACK ØØØ STATUS FROM DRIVE | | | | |
| | CLKS | CLOCK DATA | | | TRACK ØØØ | | | | | |
| | CRCIZ | CYCLIC REDUNDANCY CHECK WORD INITIALIZE | | | | WAEN | WAIT ENABLE | | | |
| | CRCOK | CRC OKAY | | | | WALT | | T READY SIGNAL | | |
| | | CARD SELECT | | | | WCLK | WRITE CLOC | | | |
| | CSAC | CARD SELECT ACCESS CONTRO | | | | WE | WRITE ENAB | | | |
| | DAL7-DALØ | DATA ACCESS LINE | | | WGI | WRITE GATE | | | | |
| | DIRECTION IN | DIRECTION CONTROL | | | -WRITE FAULT | | STATUS FROM DRI | IE. | | |
| | DHOLD | DATA HOLD | | | | WR7-WRØ | WRITE CONT | | | |
| | DHOLD | DELAYED DATA | | | | IBLA | BYTE LOOK | | | |
| | DRQ | DATA REQUEST | | | | ZXDR | | NAHEAD REFERENCE CLOCK | | |
| | DROCLK | DATA REQUEST CLOCK | | | | LADA . | E X DAIA | NEFENENCE CLOCK | • | |
| | DRSEL | DRIVE SELECT | | | | | | | | |
| | DRS4-DRSI | DRIVE SELECT BITS 4-1 | | | | | | | | |
| | DRUN | DATA RUN | | | | | | | | |
| | HERQ | HIGH FREQUENCY | | | | • | . • | | | |
| | HSAC | HOST SELECT ACCESS CONTROL | | | | NOTES: UNLESS OTHERWISE SPECIFIED | | | | |
| | HS2-HSØ | HOST SELECT ALLESS CONTROL HEAD SELECT BITS 2-0 | | | | I. RESISTOR VALUES ARE IN OHMS ±5% , 1/4 W | | | | |
| | IA9-IAØ | INSTRUCTION ADDRESS LINES | | | | | | | | |
| | IDIS-IDØ | INSTRUCTION DATA BITS 15-6 | | | JUMPER ES TO E4 TO QUALIFY WAIT BY CS JUMPER E3 TO E4 IF BIC -BOC TO QUALIFY WAIT | | | | | |
| | -INDEX | INDEX PULSE FROM DRIVE | | | | • | | | - | |
| | INTCLK | INDER PULSE FROM DRIVE INTERRUPT CLOCK INTERRUPT REQUEST | | | | A NOT USED, USED ONLY WHEN FOR EXPANDING MEMORY TO IK. | | | | |
| | INTRO | | | | | 4. ALL CAPACITOR ARE VALUED IN JA 50V AND ARE 10%. | | | | |
| | 107-10Ø | I/O LINES 7-0 | | | | | | | | |
| | IVBØ - IVB7 | INPUT VECTOR BUS BITS OF | 7 | | | A NOT USED F | OR 51/4" DRIVE. | | | |
| | MCLK | MASTER CLOCK | • | | | | UK 374 UKIVE | | | • |
| | MFMW | MODIFIED FREQUENCY MODUL | ATION WRITE STREAM | • | | A JT IS FOR | 5 1/4" DRIVE, J8 1 | S FOR &" DRIVE. | | |
| | MR | MASTER RESET | | • | | | | | | |
| | 0SC | OSCILLATOR OUTPUT | | | | | | | | |
| | RA9-RAØ | RAM ADDRESS BITS 9-0 | | | | | | | | |
| | RCLK | READ CLOCK | | | | | | | | |
| | RCS | RAM CHIP SELECT | | | | | | | | |
| | RDAT | READ DATA | | | | | | | | |
| | RD6-RD4 RD2 RDØ | READ CONTROL LINES | | | | | | | | |
| | RE | READ ENABLE | | | | | | | | |
| | -READY | READY STATUS FROM DRIVE | | | | | | | | |
| | RESET | RESET SIGNAL | | | | | | | | |
| | RGATE | READ GATE | | | | | | | | |
| | ROVE | RAM OVERFLOW | | | | | | | | \$600010-001 |
| | RWC | | | | | | - | APPLICATION 401.000 | orientes sores own H. Her II-I | 4 41 4 |
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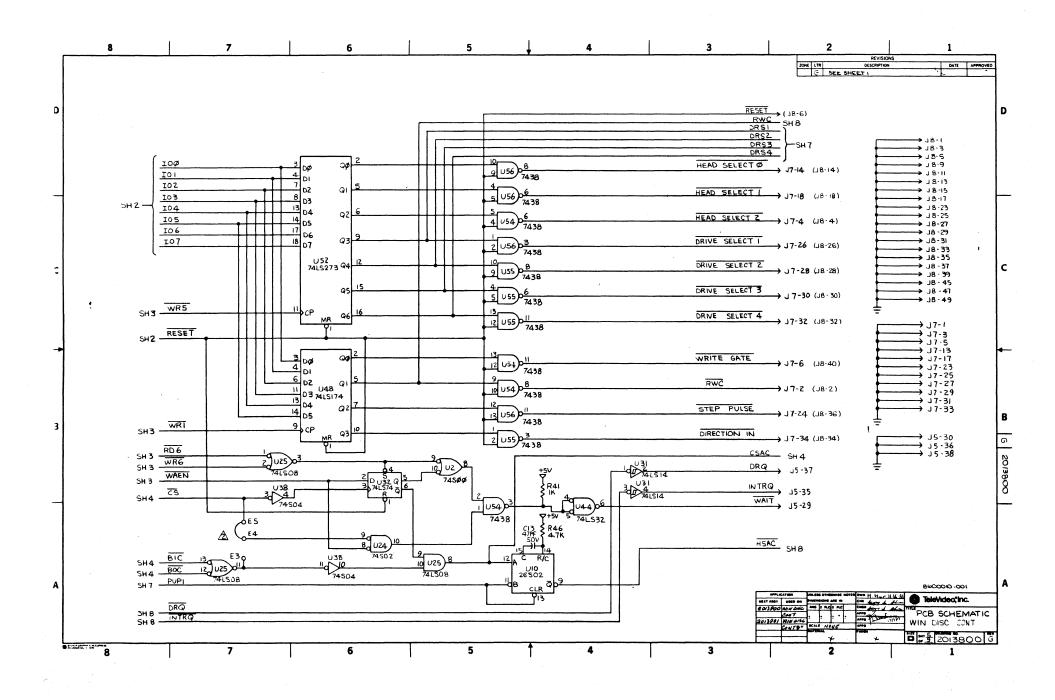


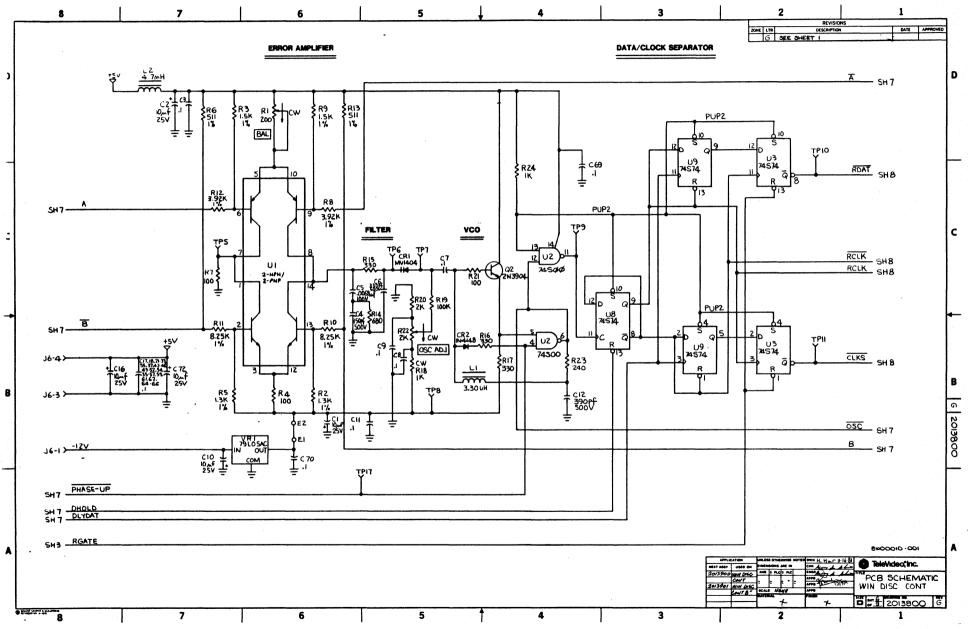






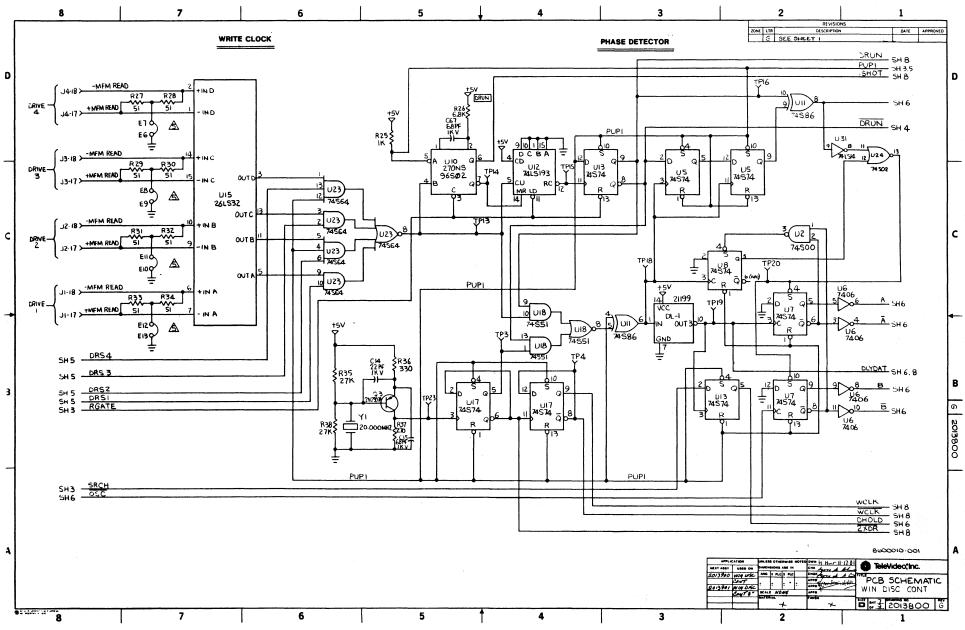




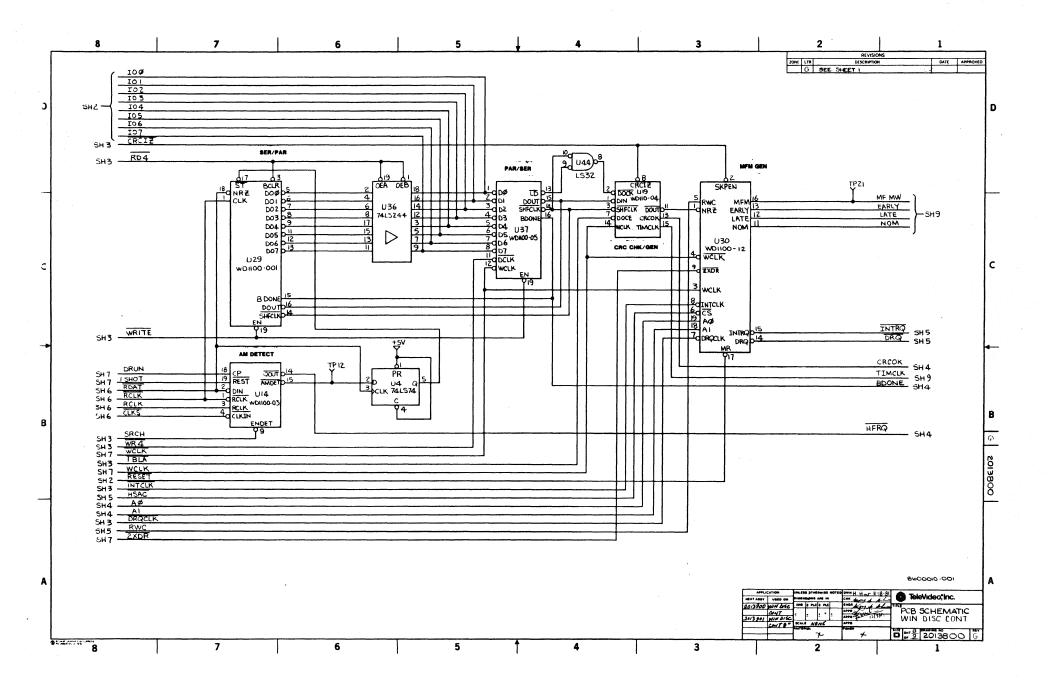


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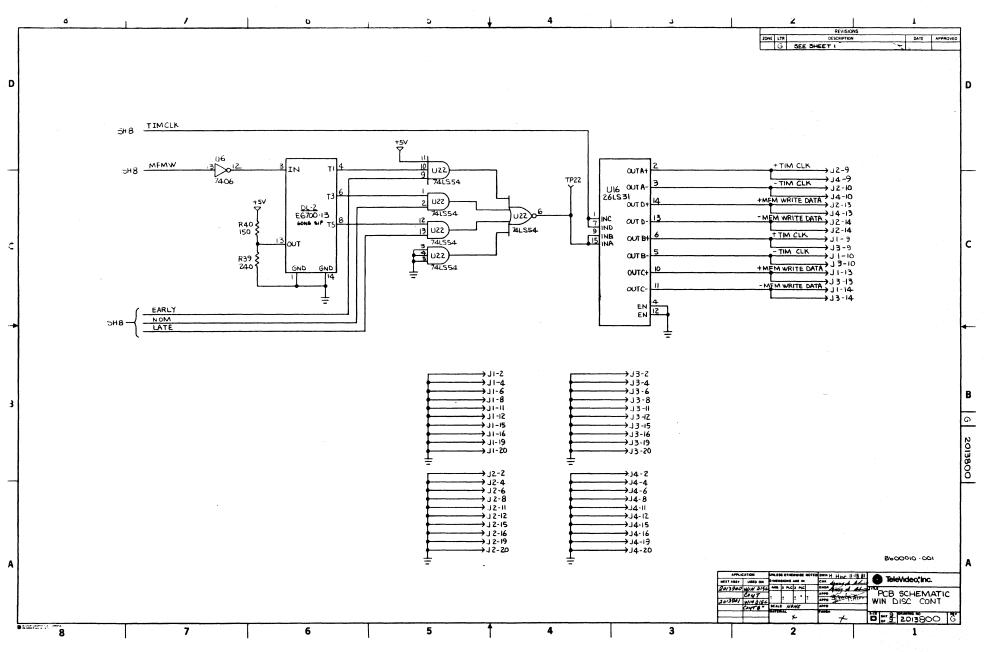




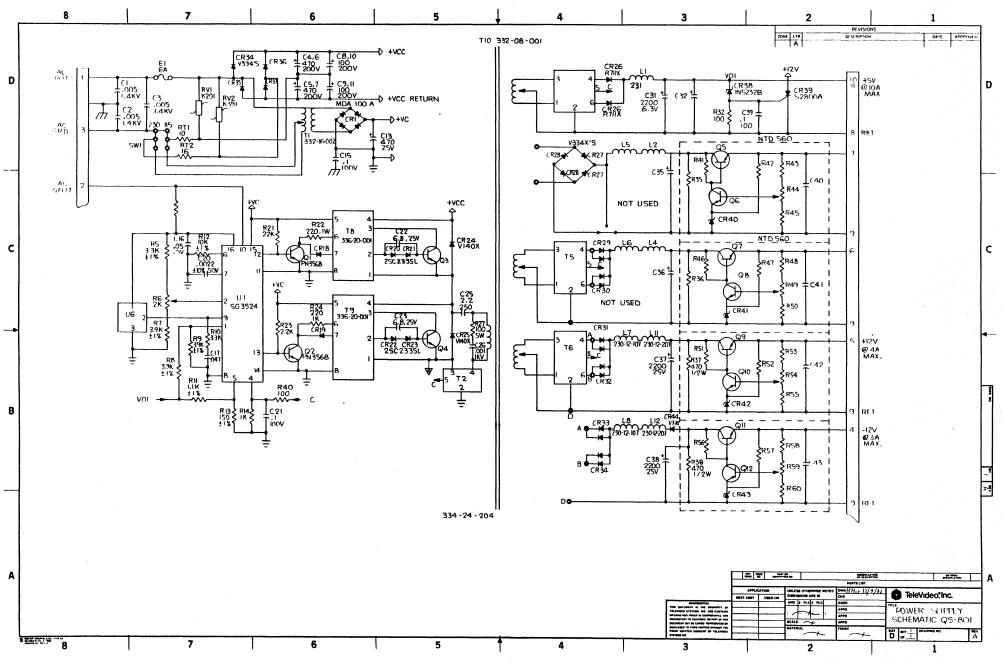












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